

# SPAs for performance modelling: Lecture 9 — Hybrid Approximation

Jane Hillston

LFCS, School of Informatics  
The University of Edinburgh  
Scotland

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*of* EDINBURGH

# Outline

- 1 Adequacy of Fluid Approximation
- 2 Hybrid approximation
  - Example
  - Illustrations of the different transitions

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## 2 Hybrid approximation

- Example
- Illustrations of the different transitions

# Introduction

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Generally we are willing to trade some **accuracy** for **efficiency** or even **tractability**.

But we should remain aware that there will be cases for which the approach is inappropriate because too much error is introduced.

# Adequacy of fluid approximation

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It does not tell us how large an  $N$  is **big enough** to count as infinity.

Moreover the existing error bounds by Darling and Norris are extremely loose, and so do not help us to predict how big  $N$  should be.

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Bradley and Hayden from Imperial College have generalised this approach to construct sets of ODEs to also approximate **higher moments**.

This offers more information about the distribution of the population count, rather than simply its expectation.

# Example model

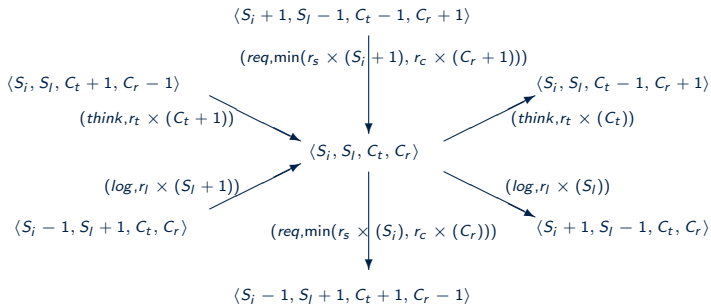
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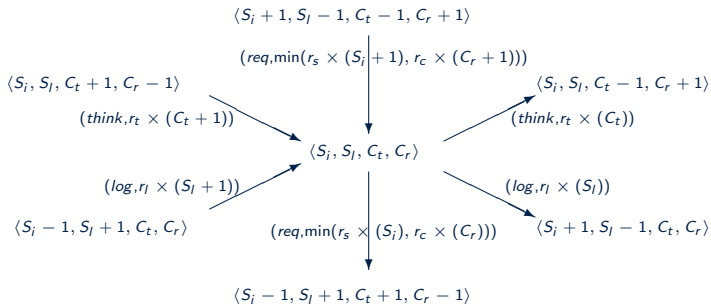
For some parametrizations of this model, the model's behaviour can accurately be characterized by the fluid flow approximations of its moments. However, for others, the moments are not sufficient to capture the model's behaviour, highlighting the danger of relying only on the results of fluid flow analysis.

# Transitions into and out of a **typical** state





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This view of the system is the basis of generating the ODEs for the moments of the system.

## Chapman-Kolmogorov equations

$$\begin{aligned}
 \frac{d p_{\langle S_i, S_l, C_t, C_r \rangle}(t)}{dt} = & \\
 & + (C_t + 1) \times r_t \times p_{\langle S_i, S_l, C_t+1, C_r-1 \rangle}(t) \\
 & + (S_l + 1) \times r_l \times p_{\langle S_i-1, S_l+1, C_t, C_r \rangle}(t) \\
 & + \min( (S_i + 1) \times r_s, (C_r + 1) \times r_c ) \times p_{\langle S_i+1, S_l-1, C_t-1, C_r+1 \rangle}(t) \\
 & - \min( S_i \times r_s, C_r \times r_c ) \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t) \\
 & - S_l \times r_l \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t) \\
 & - C_t \times r_t \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t).
 \end{aligned}$$

One variable/equation for every state of the system.

# First moment approximation

$$\begin{aligned}
 \frac{d \mathbb{E}[C_r](t)}{dt} &= \sum_{\langle S_i, S_l, C_t, C_l \rangle \in \mathbb{D}} \frac{C_r \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t)}{d t} \\
 &= + \sum_{\langle S_i, S_l, C_t, C_l \rangle \in \mathbb{D}} C_t \times r_t \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t) \\
 &\quad - \sum_{\langle S_i, S_l, C_t, C_l \rangle \in \mathbb{D}} \min(S_i \times r_s, C_r \times r_c) \times p_{\langle S_i, S_l, C_t, C_r \rangle}(t) \\
 &= + r_t \times \mathbb{E}[C_t](t) - \mathbb{E}[\min(S_i \times r_s, C_r \times r_c)](t).
 \end{aligned}$$

One variable/equation for each component of the state vector.

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 &= + r_t \times \mathbb{E}[C_t](t) - \mathbb{E}[\min(S_i \times r_s, C_r \times r_c)](t).
 \end{aligned}$$

One variable/equation for each component of the state vector.

Note  $\mathbb{E}[\min(S_i \times r_s, C_r \times r_c)]$ .

# First moment approximation

Approximating  $\mathbb{E}[\min(S_i \times r_s, C_r \times r_c)]$  with  
 $\min(\mathbb{E}[S_i \times r_s], \mathbb{E}[C_r \times r_c]) = \min(r_s \times \mathbb{E}[S_i], r_c \times \mathbb{E}[C_r])$ .

$$\begin{aligned} \frac{d \mathbb{E}' C_t(t)}{dt} &= -r_t \times \mathbb{E}' C_t(t) + \min(r_c \times \mathbb{E}' C_r(t), r_s \times \mathbb{E}' S_i(t)) \\ \frac{d \mathbb{E}' C_r(t)}{dt} &= -\min(r_c \times \mathbb{E}' C_r(t), r_s \times \mathbb{E}' S_i(t)) + r_t \times \mathbb{E}' C_t(t) \\ \frac{d \mathbb{E}' S_i(t)}{dt} &= -\min(r_c \times \mathbb{E}' C_r(t), r_s \times \mathbb{E}' S_i(t)) + r_l \times \mathbb{E}' S_l(t) \\ \frac{d \mathbb{E}' S_l(t)}{dt} &= +\min(r_c \times \mathbb{E}' C_r(t), r_s \times \mathbb{E}' S_i(t)) - r_l \times \mathbb{E}' S_l(t) \end{aligned}$$

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# Parameterisation 1

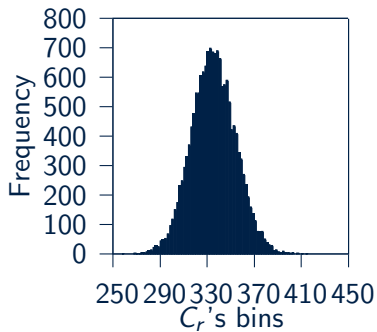
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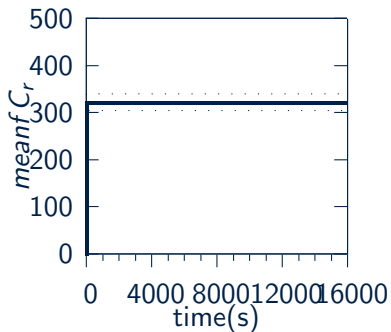
Parameter	Value	Description
$r_s$	500	On average, it takes 1/500th of an hour for a server to initiate a communication link with a client.
$r_l$	120	On average, it takes 1/120th of an hour for a server to process a request.
$c_r$	2	On average, it takes 1/2 of an hour for a client to initiate a communication link with a server.
$c_t$	0.06	On average, it takes 1/0.06th of a hours for a client to think.
$n_s$	10	Total population of servers.
$n_c$	10000	Total population of clients.

# Distribution of $C_r$ found by SSA





# Mean and std deviation found via fluid approximation



## Some numerical results

	$n_s$	3	4	5	6	7	8	9	10
$\mathbb{E}[C_r]$	F.F.A.	5645	4193	2741	1290.3	322	322	322	322
	M.C.	5644	4192	2740	1290	490	384	349	335
	Err.(%)	0.01	0.01	0.03	0.04	34	16	7.7	3.8
$\sigma[C_r]$	F.F.A.	60.62	70	78.26	85.73	17.66	17.66	17.66	17.66
	M.C.	60.45	69.45	78.79	86.5	36.90	23.49	19.84	18.72
	Err.(%)	0.26	0.25	0.67	0.9	52.3	25.20	11.44	5.6

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 S_{idle} &\stackrel{\text{def}}{=} (\text{req}, r_s).S_{logging} + (\text{brk}, r_b).S_{broken} \\
 S_{logging} &\stackrel{\text{def}}{=} (\text{log}, r_l).S_{idle} \\
 S_{broken} &\stackrel{\text{def}}{=} (\text{fix}, r_f).S_{idle} \\
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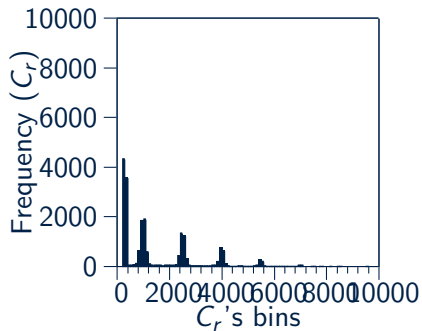
# More numerical results

	$n_s$	3	4	5	6	7	8	9	10	12	14	16	18
$\mathbb{E}[C_i]$	F.F.A.	7119	6159	5199	4239	3279	2319	1359	399	243	243	243	243
	M. C.	7156	6177	5236	4295	3387	2599	1975	1460	843	533	378	309
	Err.(%)	0.6	0.2	0.7	1.3	3.18	10.77	31.1	72.6	71.1	54.4	35.7	21
$\sigma[C_i]$	F.F.A.	1240	1432	1601	1753	1894	2025	2148	959	15.42	15.42	15.42	15.42
	M.C.	1245	1420	1609	1758	1808	1792	1656	1456	1048	713	470	314
	Err.(%)	0.42	0.79	0.52	0.25	4.7	13	29	34.1	98	97	96	95

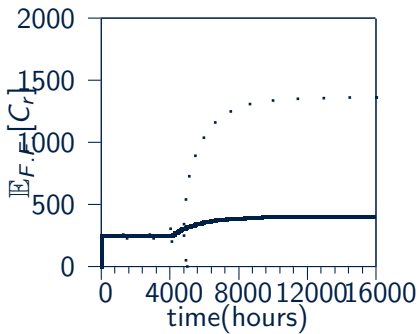
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- Example
- Illustrations of the different transitions

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However, we have also seen that there are cases where this technique should not be used because it will lead to inaccurate estimates of the performance of the system, and then simulation becomes the best way to tackle the system.

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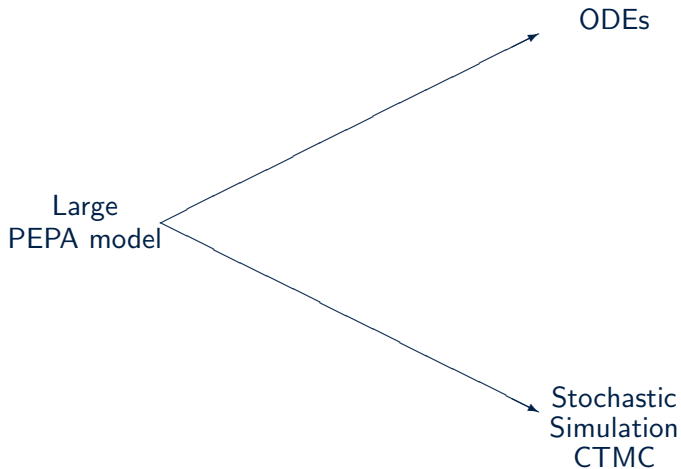
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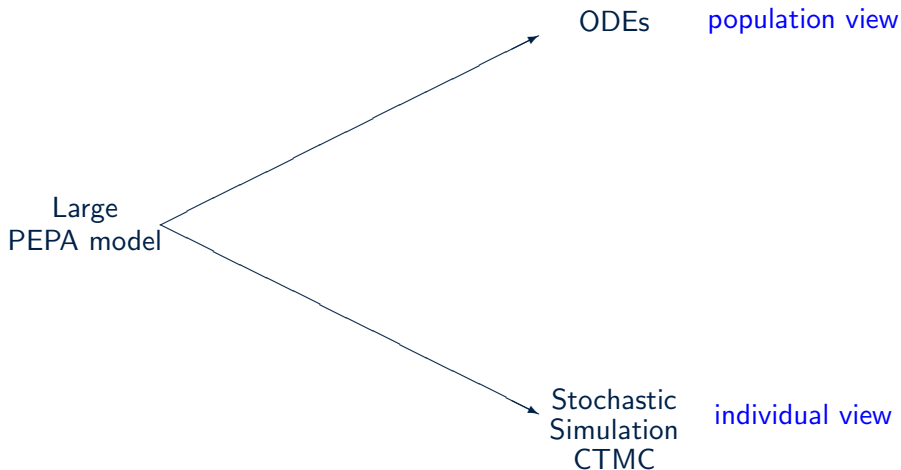
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- Typically problems arise when there is a mix of some large populations and some small, or some fast actions and some slow.
- So it is natural to consider if there might be a way to combine the approaches.
- In particular we aim to resort to the less efficient discrete approach for those parts of the model where it is strictly necessary.

# Motivation: Alternative Representations

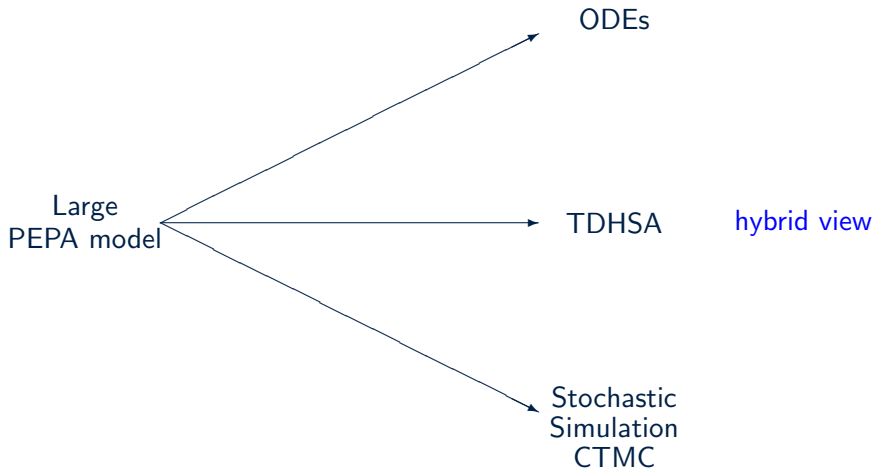




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In the **hybrid approximation** we choose to approximate **some** populations as **continuous** whilst keeping the others **discrete**.

The result is a **set of discrete states** each of which has an associated set of **ODEs**.

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But these are little difficult to work with directly, so we will use a form of automata, called **Transition Driven Stochastic Hybrid Automata (TDSHAs)** as an intermediary.

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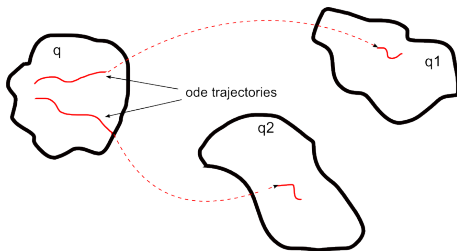
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PEPA             $\longrightarrow$             TDSHA             $\longrightarrow$             PDMP

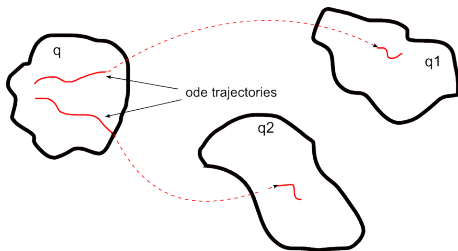
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- initial state:  $(q, (x_1, \dots, x_n))$

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  - source mode, target mode, event name
  - guard: activation condition over variables
  - reset: function determining new values of variables
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- continuous transitions (flows)
  - source mode
  - vector specifying variables involved
  - Lipschitz continuous function

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- instantaneous behaviour: fire when guard becomes true
- stochastic behaviour: fire according to rate
- product of TDSHAs
  - pairs of modes and union of variables
  - combining transitions  
(with conditions on resets and initial values)

# TDSHA synchronised product

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  - $a \notin L$ :  $(q_1, q_2)$  has every transition from  $q_1$  and from  $q_2$
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- stochastic transitions:
  - $a \notin L$ :  $(q_1, q_2)$  has every transition from  $q_1$  and from  $q_2$
  - $a \in L$ :  $(q_1, q_2)$  has every transition that both  $q_1$  and  $q_2$  have with  $a$ , new rate is PEPA cooperation rate and conjunction of resets is taken

# Overview of the mapping

- PEPA has two-level syntax
  - sequential components:  $S ::= (a, r).S \mid S + S$
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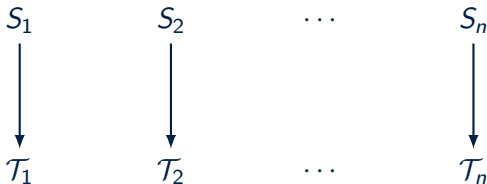
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$$\begin{array}{rcccccccc}
 P & \stackrel{\text{def}}{=} & S_1 & \underset{L_2}{\boxtimes} & S_2 & \underset{L_3}{\boxtimes} & \dots & \underset{L_n}{\boxtimes} & S_n \\
 \downarrow & & & & & & & & \\
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Any derivative that enables a **continuous action** will be treated as a **continuous variable** in the state vector.

# A client/server system with breakdowns and repairs

$$S_w \stackrel{\text{def}}{=} (\text{request}, r_{\text{reply}}).S_l + (\text{break}, r_{\text{break}}).S_b$$

$$S_l \stackrel{\text{def}}{=} (\text{log}, r_{\text{log}}).S_w$$

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 \end{aligned}$$

## State Representation

$$\omega = (\omega_{S_w}, \omega_{S_l}, \omega_{S_b}, \omega_{U_r}, \omega_{U_t})$$

Initial state is  $(1, 0, 0, N, 0)$

# In the discrete case

$$\begin{aligned}
 S_w &\stackrel{\text{def}}{=} (request, r_{reply}).S_l + (break, r_{break}).S_b \\
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## Request action

$$(1, 0, 0, N, 0) \xrightarrow{request, \min(1 \times r_{reply}, N \times r_{req})} (0, 1, 0, N - 1, 1)$$

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## Request action

$$\mathbf{x}(t) \xrightarrow{request, \min(1 \times r_{reply}, N \times r_{req})} \mathbf{x}(t) + (-1, +1, 0, -1, +1)$$

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# In the continuous case

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$$S_l \stackrel{\text{def}}{=} (log, r_{log}) \cdot S_w$$

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## ODE for $S_w$

$$\frac{dx_{S_w}(t)}{dt} = -\min(x_{S_w}(t)r_{reply}, x_{U_r}(t)r_{req}) - x_{S_w}(t)r_{break} + x_{S_l}(t)r_{log} + x_{S_r}(t)r_{re}$$



# Hybrid interpretation

$$\begin{aligned}S_w &\stackrel{\text{def}}{=} (\text{request}, r_{\text{reply}}).S_l + (\text{break}, r_{\text{break}}).S_b \\S_l &\stackrel{\text{def}}{=} (\text{log}, r_{\text{log}}).S_w \\S_b &\stackrel{\text{def}}{=} (\text{repair}, r_{\text{repair}}).S_w \\U_r &\stackrel{\text{def}}{=} (\text{request}, r_{\text{req}}).U_t \\U_t &\stackrel{\text{def}}{=} (\text{think}, r_{\text{think}}).U_r \\Sys &\stackrel{\text{def}}{=} S_w \boxtimes_{\{\text{request}\}} U_r[N]\end{aligned}$$

We may assume that the activities *break* and *repair* occur at a much lower frequency and a much lower rate than the other activities in the model.

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So for our hybrid approximation we treat these activities as **discrete** and the other activities as **continuous**.

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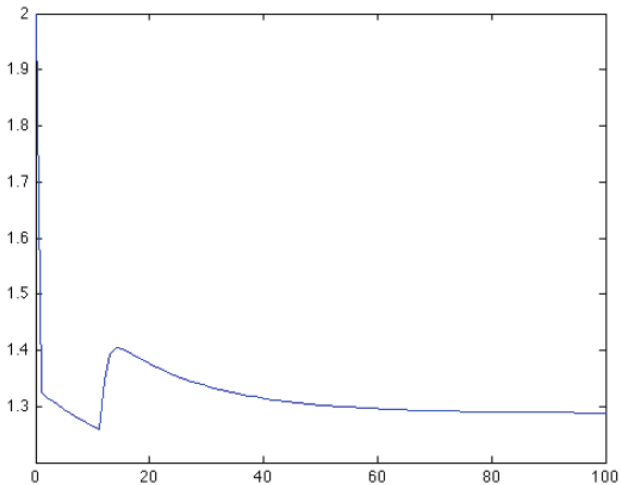
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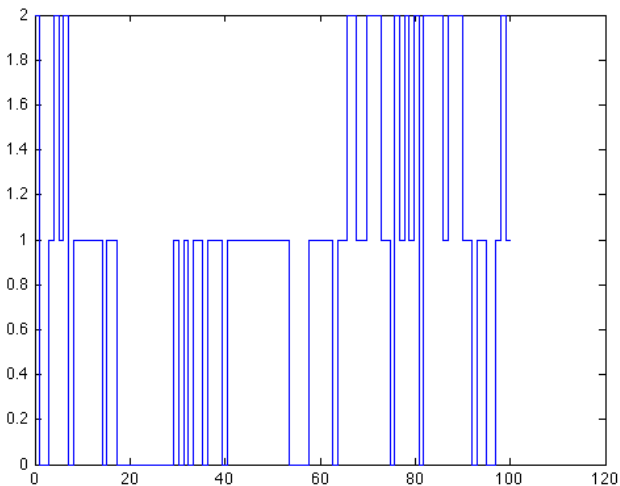
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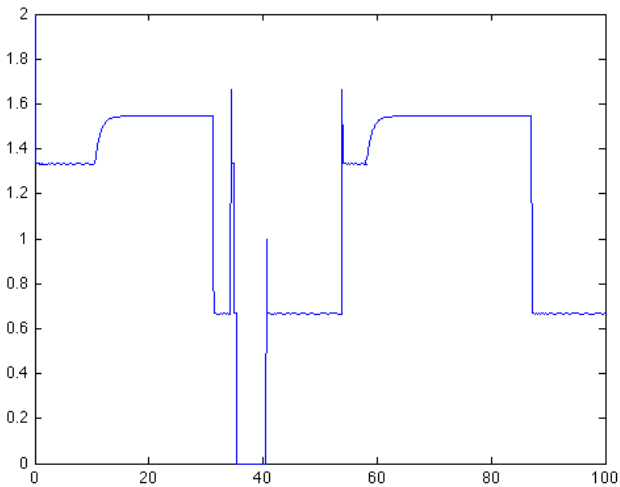
# Fluid Dynamics: Working servers vs. time



# Stochastic Dynamics: Working servers vs. time



# Hybrid Dynamics: Working servers vs. time



# Numerical Evaluation: set up

$$S_w \stackrel{\text{def}}{=} (\text{request}, \text{scale} \times 1000).S_l + (\text{break}, r_{\text{break}}).S_b$$

$$S_l \stackrel{\text{def}}{=} (\text{log}, \text{scale} \times 2000).S_w$$

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- $N_S \in \{2, 6\}$
- $N_C \in \{10, 100, 300\}$

# Numerical Evaluation: set up

For each model configuration we calculated the **steady-state probability of having 0 or 1 broken servers**.

**Errors** were computed with respect to the numerical solution of the Markov chain.



# Numerical Evaluation: results

$N_c$	$N_s$	$scale$	$\bar{X}^{S_b} = 0$	$\bar{X}^{S_b} = 1$	$H$	$S$	$S/H$
10	2	0.1	1.82%	3.58%	267	3	1.2E-2
100	2	0.1	0.67%	1.35%	1099	38	3.5E-2
300	2	0.1	0.70%	3.42%	529	69	1.3E-1
10	6	0.1	6.44%	0.52%	352	3	7.0E-3
100	6	0.1	2.35%	0.89%	566	18	3.1E-2
300	6	0.1	2.82%	1.53%	317	25	8.0E-2
10	2	10.0	0.54%	0.96%	547	253	4.6E-1
100	2	10.0	0.08%	0.21%	827	2618	3.2E+0
300	2	10.0	0.80%	3.20%	252	5092	2.0E+1
10	6	10.0	2.49%	2.64%	485	154	3.1E-1
100	6	10.0	3.86%	1.39%	623	1298	2.1E+0
300	6	10.0	1.30%	1.14%	876	5112	5.8E+0
10	2	100.0	0.13%	0.35%	204	3186	1.6E+1
100	2	100.0	0.35%	1.24%	589	20344	3.4E+1
300	2	100.0	0.01%	0.06%	438	51682	1.2E+2
10	6	100.0	2.19%	0.96%	217	1100	5.1E+0
100	6	100.0	2.14%	1.81%	301	13207	4.4E+1
300	6	100.0	0.09%	3.98%	592	39956	6.7E+1

# Ongoing issues

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- We currently assume that the modeller is responsible for partition action types and derivatives.
- There is an issue of how to make transitions from continuous state to discrete states in the general case: we have a solution but it may not be the best one.

## Illustrative example

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The example presented in the following slides is constructed to illustrate each of the different cases.

It illustrates some of the problems that can occur and our current solution to these problems.

# Clients and servers example

- clients

$$\text{Cr} \stackrel{\text{def}}{=} (\text{request}, r_{rq}).\text{Ct}$$
$$\text{Ct} \stackrel{\text{def}}{=} (\text{think}, r_{th}).\text{Cr}$$

# Clients and servers example

## ■ clients

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## ■ servers

$$\text{Sr} \stackrel{\text{def}}{=} (\text{request}, r_{rp}).\text{Sl} + (\text{break}, r_{bk}).\text{Sb}$$
$$\text{Sl} \stackrel{\text{def}}{=} (\text{log}, r_{lg}).\text{Sr} + (\text{remove}, r_{rm}).\text{Sm}$$
$$\text{Sm} \stackrel{\text{def}}{=} (\text{maint}, r_{mn}).\text{Sr} + (\text{replace}, r_{rc}).\text{Sr}$$
$$\text{Sb} \stackrel{\text{def}}{=} (\text{fix}, r_{fx}).\text{St}$$
$$\text{St} \stackrel{\text{def}}{=} (\text{test}, r_{ts}).\text{St} + (\text{compl}, r_{cm}).\text{Sr}$$

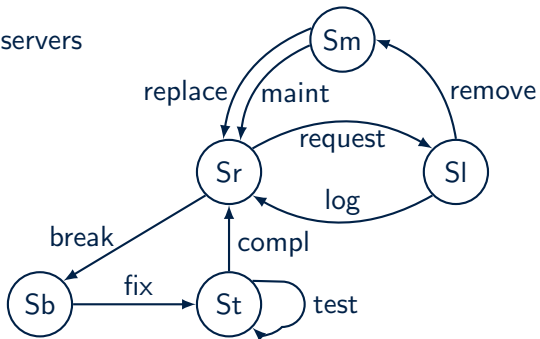


# Clients and servers example

## ■ clients



## ■ servers

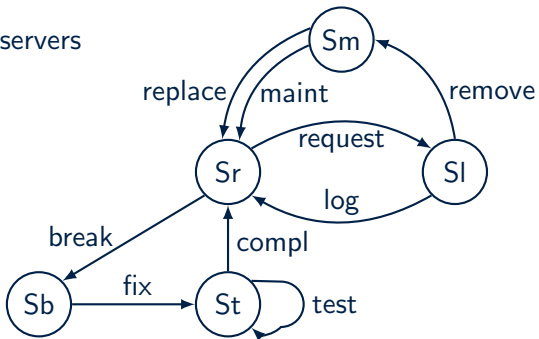


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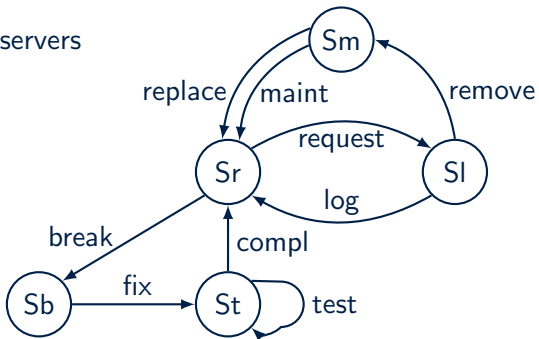


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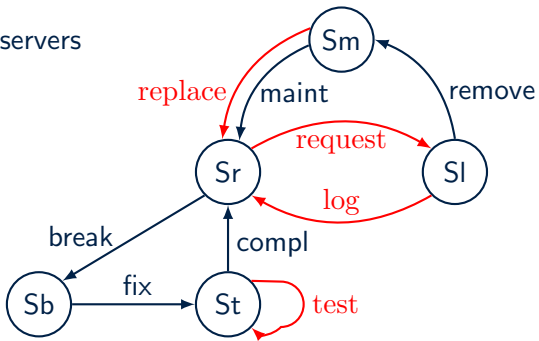


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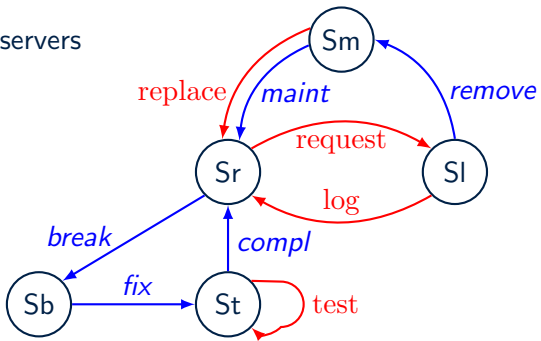


# Clients and servers example

■ clients



■ servers

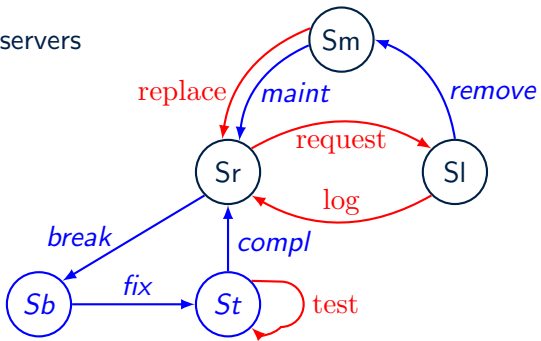


# Clients and servers example

## ■ clients



## ■ servers

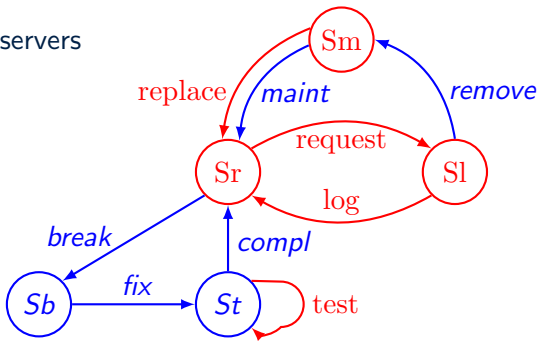


# Clients and servers example

## ■ clients



## ■ servers



# Clients and servers example

## ■ clients

$$\text{Cr} \stackrel{\text{def}}{=} (\text{request}, r_{rq}).\text{Ct}$$
$$\text{Ct} \stackrel{\text{def}}{=} (\text{think}, r_{th}).\text{Cr}$$

## ■ servers

$$\text{Sr} \stackrel{\text{def}}{=} (\text{request}, r_{rp}).\text{Sl} + (\text{break}, r_{bk}).\text{Sb}$$
$$\text{Sl} \stackrel{\text{def}}{=} (\text{log}, r_{lg}).\text{Sr} + (\text{remove}, r_{rm}).\text{Sm}$$
$$\text{Sm} \stackrel{\text{def}}{=} (\text{maint}, r_{mn}).\text{Sr} + (\text{replace}, r_{rc}).\text{Sr}$$
$$\text{Sb} \stackrel{\text{def}}{=} (\text{fix}, r_{fx}).\text{St}$$
$$\text{St} \stackrel{\text{def}}{=} (\text{test}, r_{ts}).\text{St} + (\text{compl}, r_{cm}).\text{Sr}$$



# Mapping to TDSHA

- continuous sequential components: *Cr, Ct, Sr, Sl, Sm*
- integral sequential components: *Sb, St*

# Mapping to TDSHA

- continuous sequential components:  $Cr, Ct, Sr, Sl, Sm$
- integral sequential components:  $Sb, St$
- population vector:  $(\#Cr, \#Ct, \#Sr, \#Sl, \#Sm, \#Sb, \#St)$

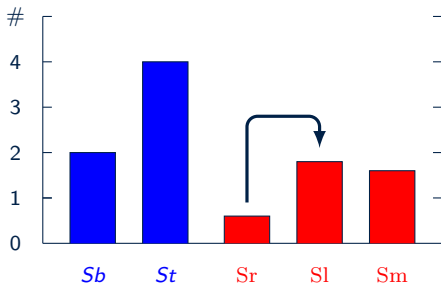
# Mapping to TDSHA

- continuous sequential components:  $\text{Cr}, \text{Ct}, \text{Sr}, \text{Sl}, \text{Sm}$
- integral sequential components:  $\text{Sb}, \text{St}$
- population vector:  $(\#\text{Cr}, \#\text{Ct}, \#\text{Sr}, \#\text{Sl}, \#\text{Sm}, \#\text{Sb}, \#\text{St})$
- PEPA is conservative: both  $N_C = \#\text{Cr} + \#\text{Ct}$  and  $N_S = \#\text{Sr} + \#\text{Sl} + \#\text{Sm} + \#\text{Sb} + \#\text{St}$  are invariant
- TDSHA
  - modes:  $(\#\text{Sb}, \#\text{St}) \in \{0, \dots, N_S\} \times \{0, \dots, N_S\}$
  - variables:  $(X_{\text{Cr}}, X_{\text{Ct}}, X_{\text{Sr}}, X_{\text{Sl}}, X_{\text{Sm}})$
  - initial state:  $((\#\text{Sb}, \#\text{St}), (\#\text{Cr}, \#\text{Ct}, \#\text{Sr}, \#\text{Sl}, \#\text{St}))$
  - continuous and stochastic transitions

# Continuous transitions between continuous components

■  $Sr \xrightarrow{(request, r_{rp} \cdot \#Sr)}_* Sl$

■ continuous transition: flow is determined by ODEs

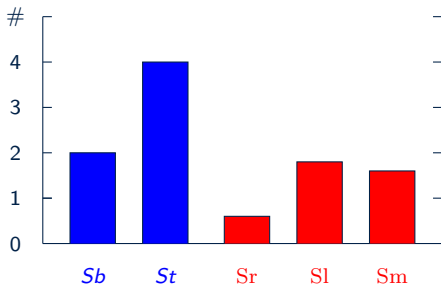


■  $((\#Sb, \#St), (0, 0, -1, 1, 0), r_{rp} \cdot \#Sr, request)$

# Continuous transition at a discrete component

■  $St \xrightarrow{(\text{test}, r_{ts} \cdot \#St)}_* St$

- continuous transition: no flow because single component

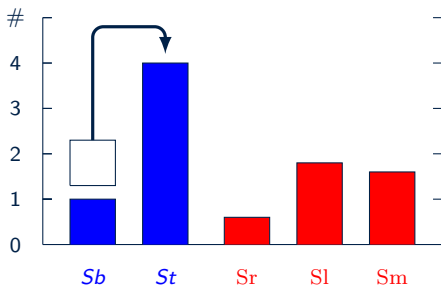


■  $((\#Sb, \#St), (0, 0, 0, 0, 0), r_{ts} \cdot \#St, \text{request})$

# Discrete transitions between discrete components

■  $Sb \xrightarrow{(fix, r_{fx} \cdot \#Sb)} \star St$

- stochastic transition: unit quantity is shifted

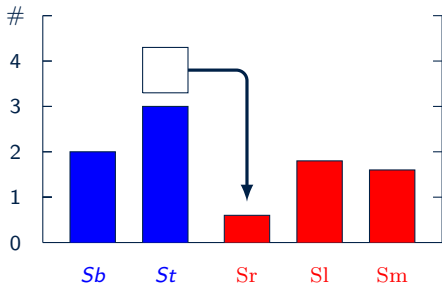


■  $((\#Sb, \#St), (\#Sb - 1, \#St + 1), true, true, r_{fx} \cdot \#Sb, fix)$

# Discrete transition from discrete to continuous component

■  $St \xrightarrow{(compl, r_{cm} \cdot \#St)} \star Sr$

- stochastic transition: unit quantity is shifted

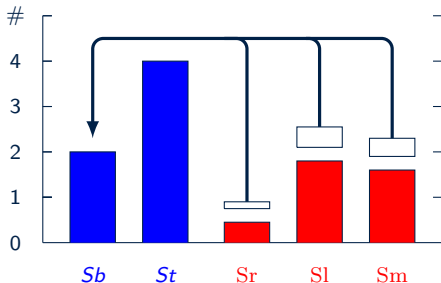


- $((\#Sb, \#St), (\#Sb, \#St - 1), true, R, r_{cm} \cdot \#St, compl)$  with  
 $R = (X'_{Sr} = X_{Sr} + 1)$

# Discrete transition from continuous to discrete component

■  $S_r \xrightarrow{(break, r_{bk} \cdot \#S_r)} \star S_b$

- stochastic transition: unit quantity is shifted proportionally



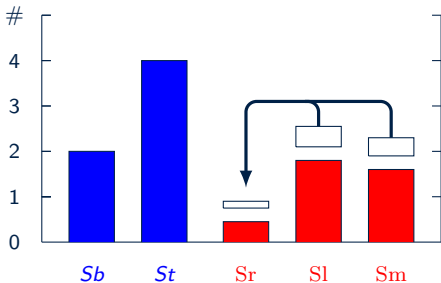
- $((\#S_b, \#S_t), (\#S_b + 1, \#S_t), true, R, r_{bk} \cdot \#S_r, break)$  with  
 $R = (X'_{S_r} = X_{S_r} - z_r) \wedge (X'_{S_l} = X_{S_l} - z_l) \wedge (X'_{S_m} = X_{S_m} - z_m)$   
 and  $z_r + z_l + z_m = 1$



# Discrete transition between continuous components

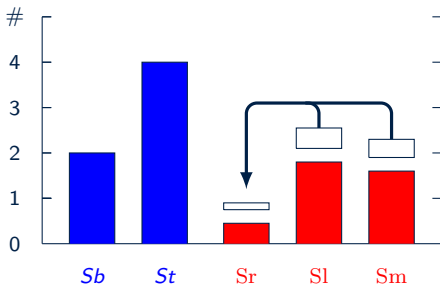
■  $S_m \xrightarrow{(maint, r_{mn} \cdot \#S_m)} S_r$

- stochastic transition: unit quantity is shifted proportionally

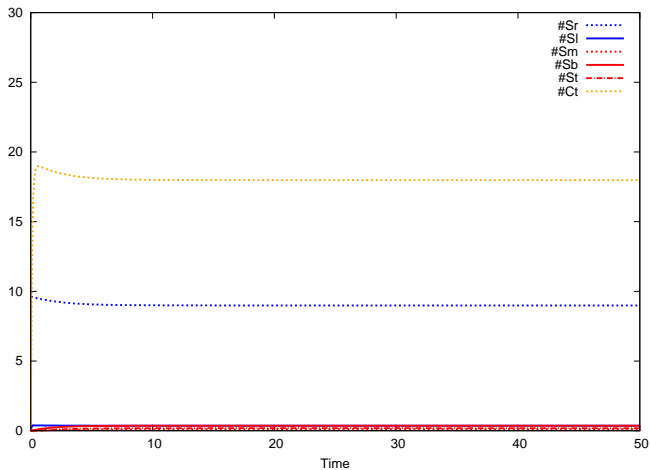


# Discrete transition between continuous components

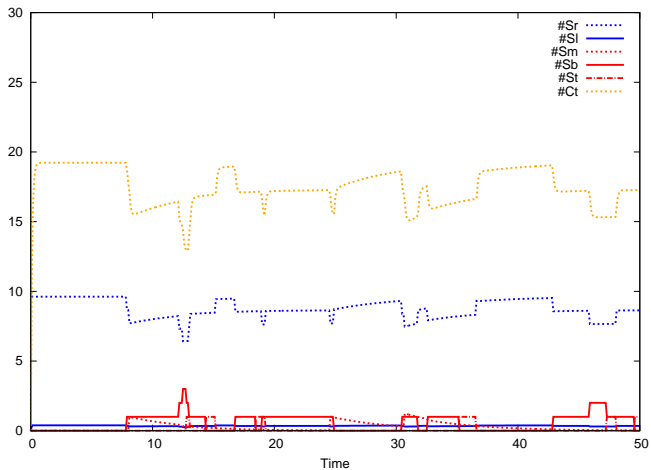
- $((\#S_b, \#S_t), (\#S_b, \#S_t), \text{true}, R, r_{mn} \cdot \#S_m, \text{maint})$  where  
 $R = (X'_{S_r} = X_{S_r} - z_r + 1) \wedge (X'_{S_l} = X_{S_l} - z_l) \wedge (X'_{S_m} = X_{S_m} - z_m)$   
 and  $z_r + z_l + z_m = 1$



# Continuous deterministic simulation



# Hybrid simulation



# Conclusions

- The hybrid semantics for PEPA is a bridge between the fully discrete approach and the deterministic approach of fluid approximation.

# Conclusions

- The hybrid semantics for PEPA is a bridge between the fully discrete approach and the deterministic approach of fluid approximation.
- The numerical results suggest that hybrid simulation may yield accurate results faster than full stochastic simulation