# Database Design

□Goal: specification of database schema □Methodology:

- Use E-R model to get a high-level graphical view of essential components of the model and how they are related
- Convert E-R diagram to DDL
- **□**E-R Model is viewed as a set of
  - Entities
  - Relationships among entities

#### Entities

□ *Entity*: an object that is involved in the enterprise

- Ex: Mike, DBS
- □ *Entity Type*: set of similar objects
  - Ex: students, courses, professors
- □ *Attribute*: describes one aspect of an entity type
  - Ex: name, maximum enrollment, cno, etc.

# Entity Type

□ Entity type described by set of attributes

• Person: *Id*, *Name*, *Address*, *Hobbies* 

*Domain*: possible values of an attribute

- Value can be a set (in contrast to relational model)
  - (123456789, Mike, 123 College, {stamps, coins})
- □ *Key*: minimum set of attributes that uniquely identifies an entity
- □ *Entity Schema*: entity type name, attributes (and associated domain), key constraints

# Entity Type cont'd

#### Graphical Representation in E-R diagram:



# Relationships

*Relationship*: relates two or more entities

John *majors in* Computer Science

*Relationship Type*: set of similar relationships

 Student (entity type) related to Department (entity type) by MajorsIn (relationship type).

Distinction:

- *relation* (relational model) set of tuples
- *relationship* (E-R Model) describes relationship between entities Both entity types and relationship types (E-R model) may be represented as relations (in the relational model)

#### Attributes and Roles

□ *Attribute* of a relationship type describes the relationship

- e.g., Mike majors in CS *since* 2004
  - Mike and CS are related
  - 2004 describes relationship value of SINCE attribute of MajorsIn relationship type

□ *Role* of a relationship type names one of the related entities

- e.g., Mike is value of *Student* role, CS value of *Department* role of MajorsIn relationship type
- (Mike, CS; 2004) describes a relationship

#### **Graphical Representation**

Roles are edges labeled with role names (omitted if role name = name of entity set). Most attributes have been omitted.



### Single-role Key Constraint

- □ If, for a particular participant entity type, each entity participates in *at most* one relationship, corresponding role is a key of relationship type
  - E.g., *Professor* role is unique in WorksIn
- □ Representation in E-R diagram: arrow



# Entity Type Hierarchies

□One entity type might be subtype of another

- Freshman is a subtype of Student
- □ A relationship exists between a Freshman entity and the corresponding Student entity
  - e.g., Freshman John is related to Student John
- □ This relationship is called *IsA* 
  - Freshman IsA Student
  - The two entities related by IsA are always descriptions of the same real-world object

## Properties of IsA

□ *Inheritance* - Attributes of supertype apply to subtype.

- E.g., *GPA* attribute of Student applies to Freshman
- Subtype *inherits* all attributes of supertype.
- Key of supertype is key of subtype
- □ *Transitivity* Hierarchy of IsA
  - Student is subtype of Person, Freshman is subtype of Student, so Freshman is also a subtype of Student

# Advantages of IsA

- □Can create a more concise and readable E-R diagram
  - Attributes common to different entity sets need not be repeated
  - They can be grouped in one place as attributes of supertype
  - Attributes of (sibling) subtypes can be different

#### IsA Hierarchy - Example



#### Participation Constraint

□ If every entity participates in *at least* one relationship, a *participation constraint* holds:

• e.g., every professor works in *at least* one department



### Participation *and* Key Constraint

- □ If every entity participates in *exactly* one relationship, both a participation and a key constraint hold:
  - e.g., every professor works in *exactly one* department



#### Representation of Entity Types in the Relational Model

An entity type corresponds to a relation
Relation's attributes = entity type's attributes

*Problem*: entity type can have set valued attributes, e.g.,

Person: *Id, Name, Address, Hobbies* 

- *Solution*: Use several rows to represent a single entity
  - (123456789, Mike, 123 College St, stamps)
  - (123456789, Mike, 123 College St, coins)
- Problems with this solution:
  - Redundancy
  - Key of entity type (Id) not key of relation
  - Hence, the resulting relation must be further transformed (we'll see how later)

#### Representation of Relationship Types in the Relational Model

- □ Typically, a relationship becomes a relation in the relational model
- □ Attributes of the corresponding relation are
  - Attributes of relationship type
  - For each role, the primary key of the entity type associated with that role

**D** *Example*:



- Professor (<u>Id</u>, DeptId, Name)
- Teaching (<u>CrsCode</u>, <u>SecNo</u>, Id, RoomNo, <u>TAs</u>)

Representation of Relationship Types in the Relational Model

□Key of corresponding table = key of relation

- Except when there are set valued attributes
- Example: Teaching (CrsCode, SectNo, Id, RoomNo, TAs)
  - Key of relationship type = (*CrsCode, SectNo*)
  - Key of relation = (*CrsCode, SectNo, TAs*)



CrsCode	SectNo	Id	RoomNo	TAs
CSC343	1	1234	BA1180	Joe
CSC343	2	1234	GB119	Mary

### Representation in SQL

- Each role of relationship type produces a foreign key in corresponding relation
  - Foreign key references table corresponding to entity type from which role values are drawn

### Example 1



CREATE TABLE WorksIn (

Since DATE, -- attribute Status CHAR (10), -- attribute ProfId INTEGER, -- role (key of Professor) DeptId CHAR (4), -- role (key of Department) PRIMARY KEY (ProfId), -- since a professor works in at most one department FOREIGN KEY (ProfId) REFERENCES Professor (Id), FOREIGN KEY (DeptId) REFERENCES Department )



#### Representing Participation Constraints in the Relational Model



□ *Inclusion dependency*: Every professor works in *at least* one dep't.

- in the relational model: (easy)
  - Professor (Id) references WorksIn (ProfId)
- in SQL:
  - Simple case: *If ProfId is a key in* **WorksIn** (i.e., every professor works in *exactly one* department) then it is easy:
    - FOREIGN KEY *Id* REFERENCES WorksIn (*ProfId*)
  - General case *ProfId* is *not* a key in **WorksIn**, so can't use foreign key constraint (not so easy):

CREATE ASSERTION ProfsInDepts CHECK ( NOT EXISTS ( SELECT \* FROM Professor P WHERE NOT EXISTS ( SELECT \* FROM WorksIn W WHERE P.Id = W.ProfId ) ) )

#### Representing Participation Constraint in the Relational Model

Example (can't use foreign key in Professor if ProfId is not a candidate key in WorksIn)



# Representing Participation *and* Key Constraint in SQL

If both participation and key constraints apply, use foreign key constraint in entity table (but beware: if candidate key in entity table is not primary, presence of nulls violates participation constraint).

CREATE TABLE Professor ( *Id* INTEGER,

PRIMARY KEY (*Id*), -- *Id can't be null* FOREIGN KEY (*Id*) REFERENCES WorksIn (*ProfId*) --all professors participate



#### Entity or Attribute?

□ Sometimes information can be represented as either an entity or an attribute.



### Entity or Relationship?

