

Database Design

- ❑ Goal: specification of database schema
- ❑ Methodology:
 - Use E-R model to get a high-level graphical view of essential components of the model and how they are related
 - Convert E-R diagram to DDL
- ❑ E-R Model is viewed as a set of
 - Entities
 - Relationships among entities

Entities

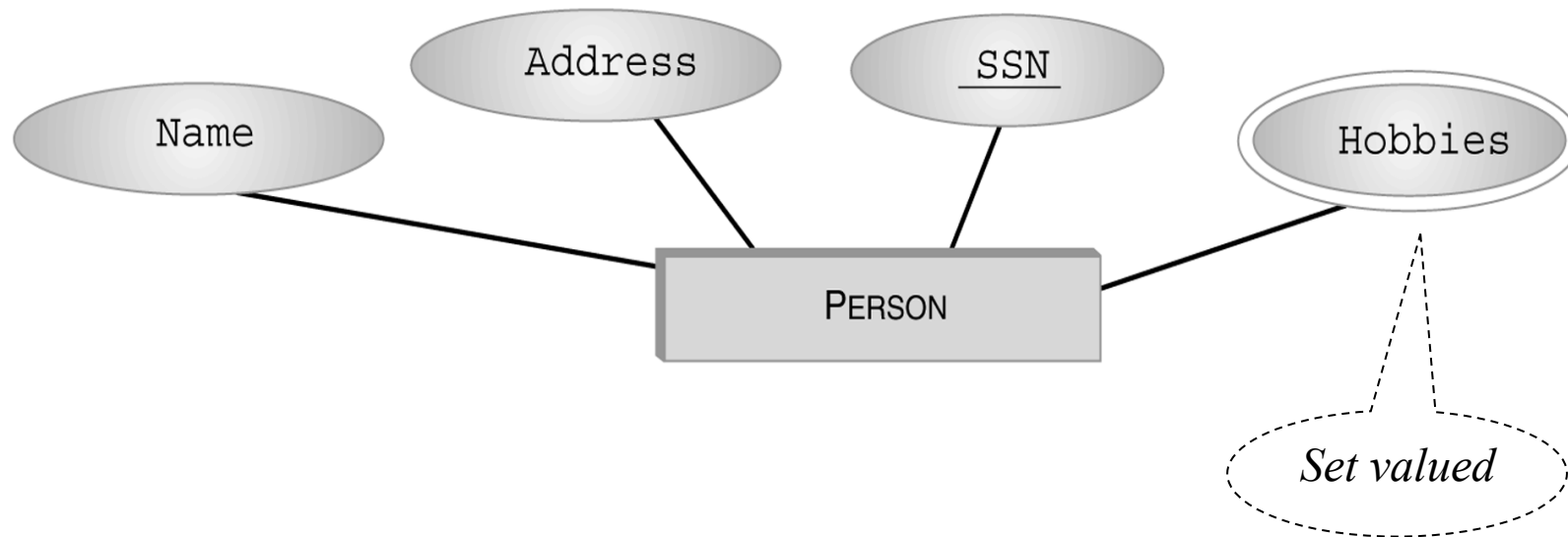
- ❑ *Entity*: an object that is involved in the enterprise
 - Ex: Mike, DBS
- ❑ *Entity Type*: set of similar objects
 - Ex: students, courses, professors
- ❑ *Attribute*: describes one aspect of an entity type
 - Ex: *name, maximum enrollment, cno, etc.*

Entity Type

- ❑ Entity type described by set of attributes
 - Person: *Id, Name, Address, Hobbies*
- ❑ *Domain*: possible values of an attribute
 - Value can be a set (in contrast to relational model)
 - (123456789, Mike, 123 College, {stamps, coins})
- ❑ *Key*: minimum set of attributes that uniquely identifies an entity
- ❑ *Entity Schema*: entity type name, attributes (and associated domain), key constraints

Entity Type cont' d

□ Graphical Representation in E-R diagram:



Relationships

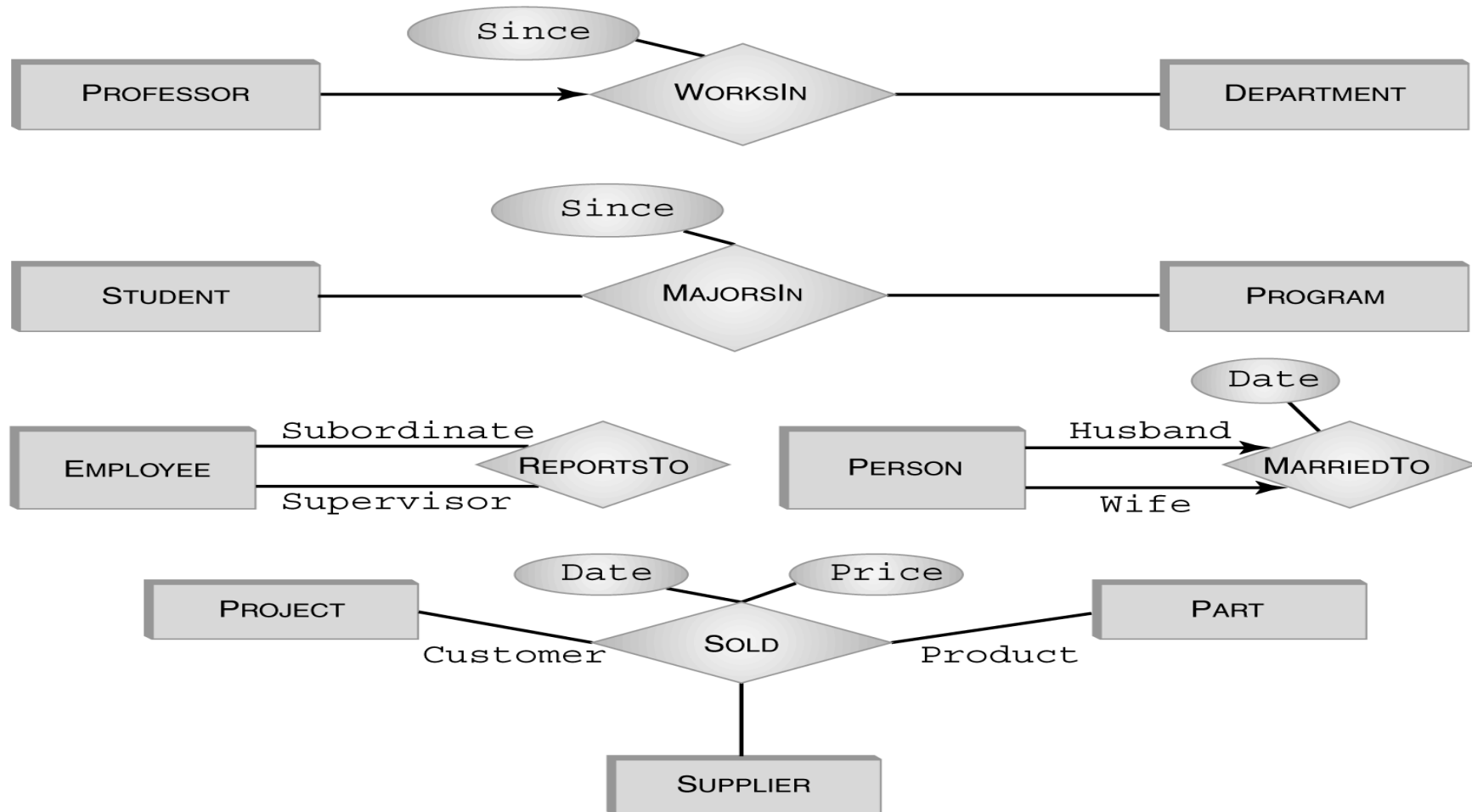
- ❑ *Relationship*: relates two or more entities
 - John *majors in* Computer Science
- ❑ *Relationship Type*: set of similar relationships
 - Student (entity type) related to Department (entity type) by MajorsIn (relationship type).
- ❑ Distinction:
 - *relation* (relational model) - set of tuples
 - *relationship* (E-R Model) – describes relationship between entities Both entity types and relationship types (E-R model) may be represented as relations (in the relational model)

Attributes and Roles

- ❑ *Attribute* of a relationship type describes the relationship
 - e.g., Mike majors in CS *since* 2004
 - Mike and CS are related
 - 2004 describes relationship - value of SINCE attribute of MajorsIn relationship type
- ❑ *Role* of a relationship type names one of the related entities
 - e.g., Mike is value of *Student* role, CS value of *Department* role of MajorsIn relationship type
 - (Mike, CS; 2004) describes a relationship

Graphical Representation

- ❑ Roles are edges labeled with role names (omitted if role name = name of entity set). Most attributes have been omitted.



Single-role Key Constraint

- ❑ If, for a particular participant entity type, each entity participates in *at most* one relationship, corresponding role is a key of relationship type
 - E.g., *Professor* role is unique in *WorksIn*
- ❑ Representation in E-R diagram: arrow



Entity Type Hierarchies

- ❑ One entity type might be subtype of another
 - Freshman is a subtype of Student
- ❑ A relationship exists between a Freshman entity and the corresponding Student entity
 - e.g., Freshman John is related to Student John
- ❑ This relationship is called *IsA*
 - Freshman IsA Student
 - The two entities related by IsA are always descriptions of the same real-world object

Properties of IsA

□ *Inheritance* - Attributes of supertype apply to subtype.

- E.g., *GPA* attribute of *Student* applies to *Freshman*
- Subtype *inherits* all attributes of supertype.
- Key of supertype is key of subtype

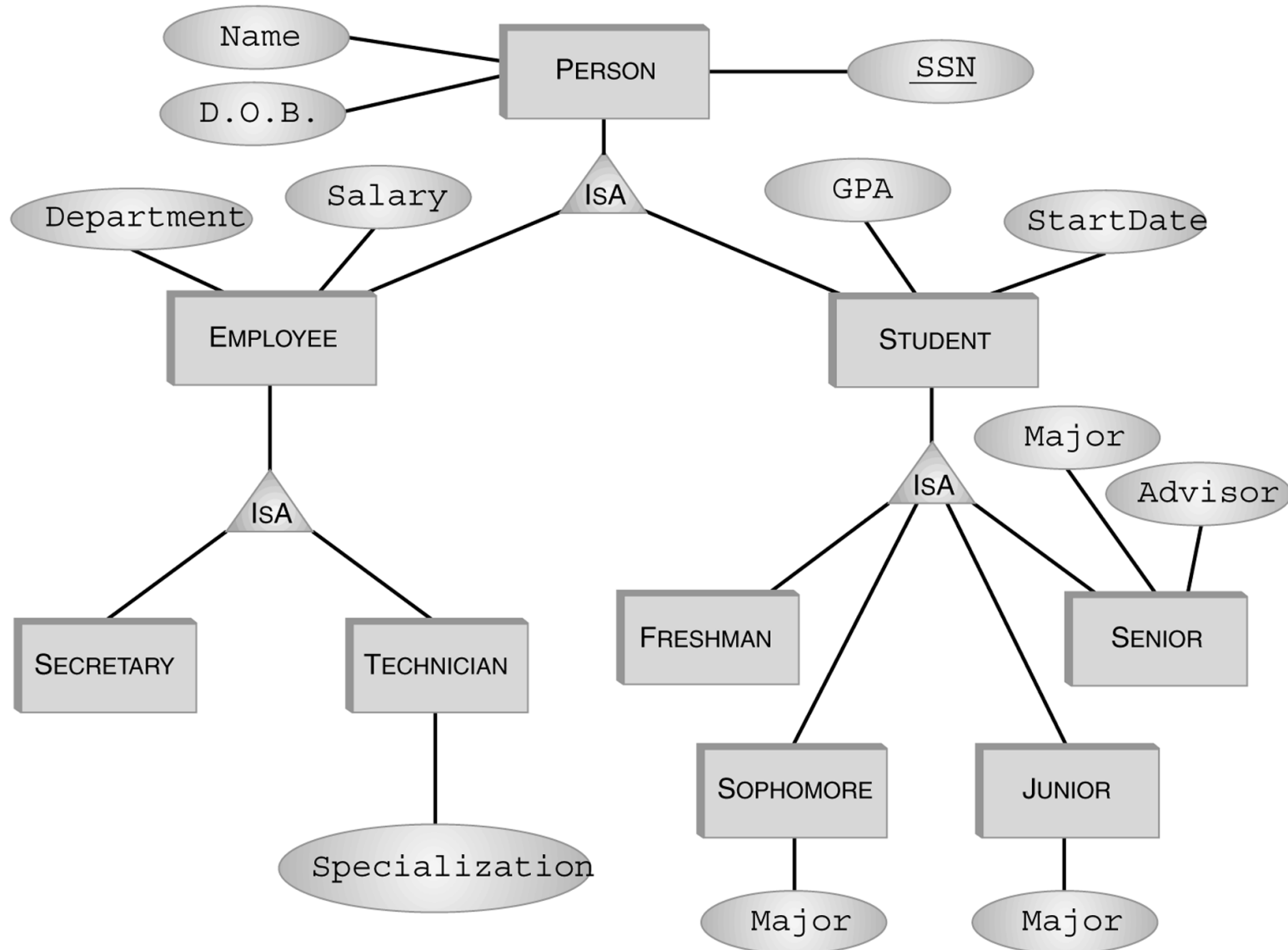
□ *Transitivity* - Hierarchy of IsA

- *Student* is subtype of *Person*, *Freshman* is subtype of *Student*, so *Freshman* is also a subtype of *Student*

Advantages of IsA

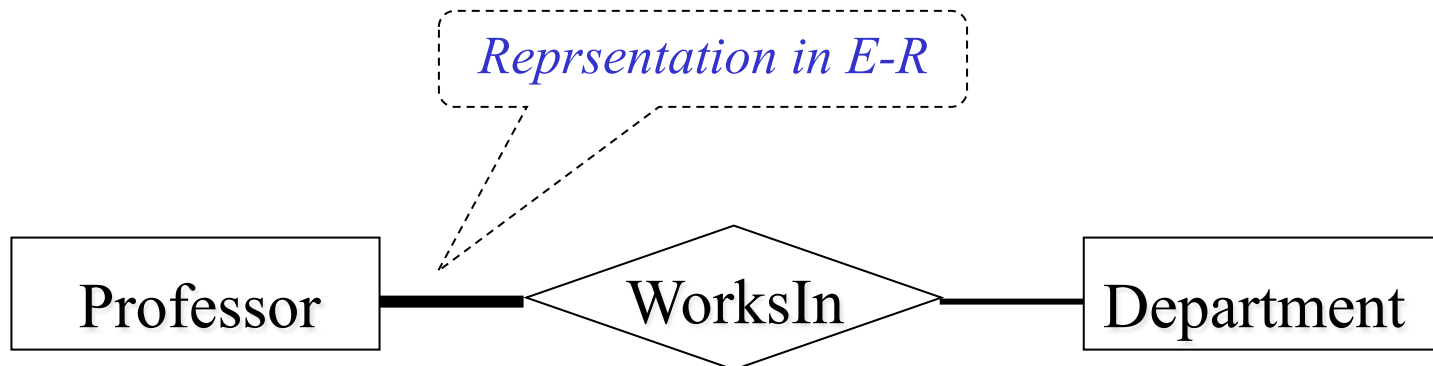
- ❑ Can create a more concise and readable E-R diagram
 - Attributes common to different entity sets need not be repeated
 - They can be grouped in one place as attributes of supertype
 - Attributes of (sibling) subtypes can be different

IsA Hierarchy - Example



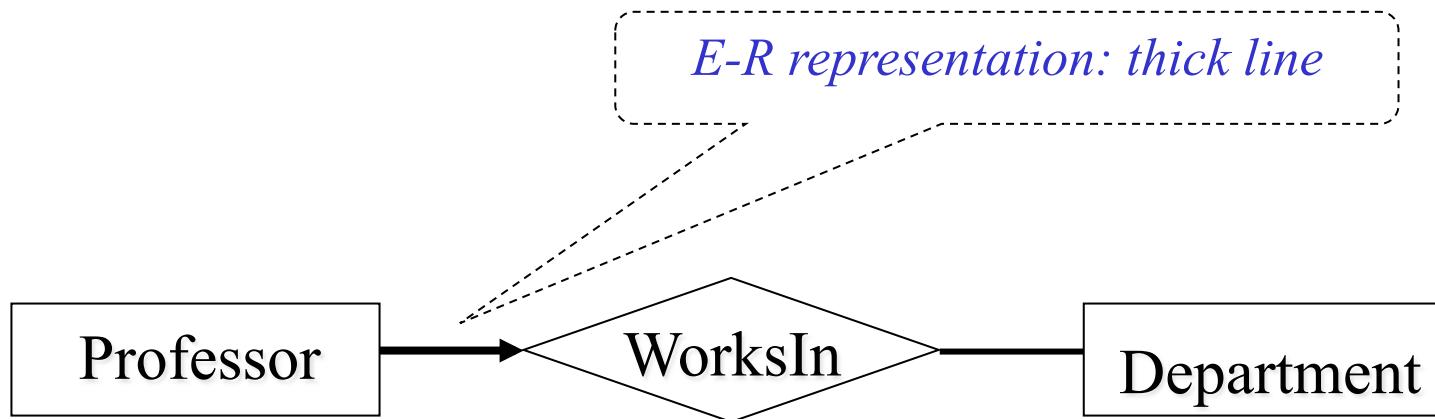
Participation Constraint

- If every entity participates in *at least* one relationship, a *participation constraint* holds:
 - e.g., every professor works in *at least* one department



Participation *and* Key Constraint

- If every entity participates in *exactly* one relationship, both a participation and a key constraint hold:
 - e.g., every professor works in *exactly one* department

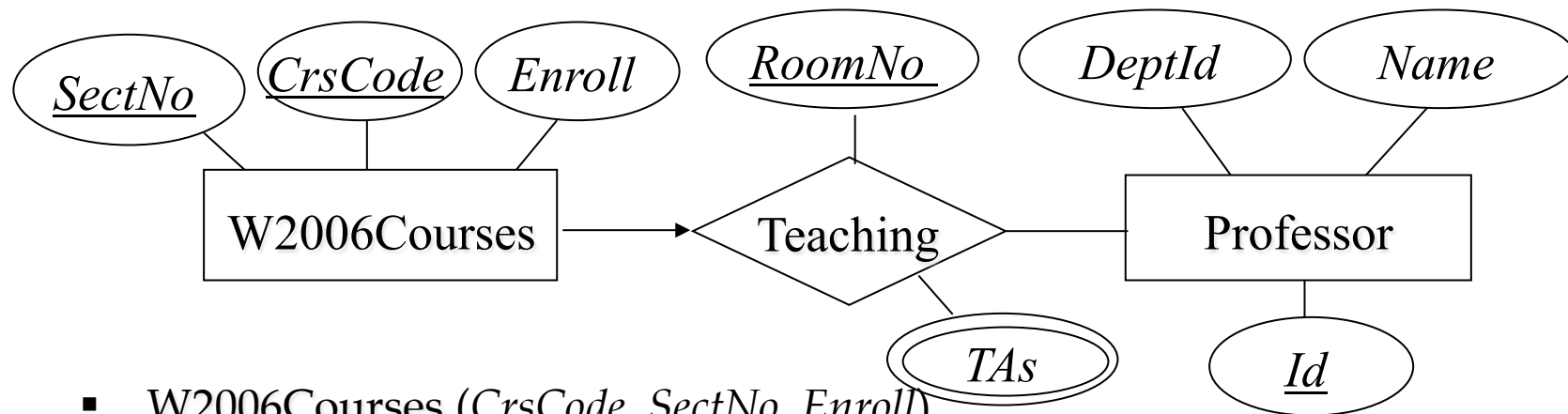


Representation of Entity Types in the Relational Model

- ❑ An entity type corresponds to a relation
- ❑ Relation's attributes = entity type's attributes
 - *Problem:* entity type can have set valued attributes, e.g.,
 Person: *Id, Name, Address, Hobbies*
 - *Solution:* Use several rows to represent a single entity
 - (123456789, Mike, 123 College St, stamps)
 - (123456789, Mike, 123 College St, coins)
 - Problems with this solution:
 - Redundancy
 - Key of entity type (Id) not key of relation
 - Hence, the resulting relation must be further transformed (we'll see how later)

Representation of Relationship Types in the Relational Model

- ❑ Typically, a relationship becomes a relation in the relational model
- ❑ Attributes of the corresponding relation are
 - Attributes of relationship type
 - For each role, the primary key of the entity type associated with that role
- ❑ *Example:*

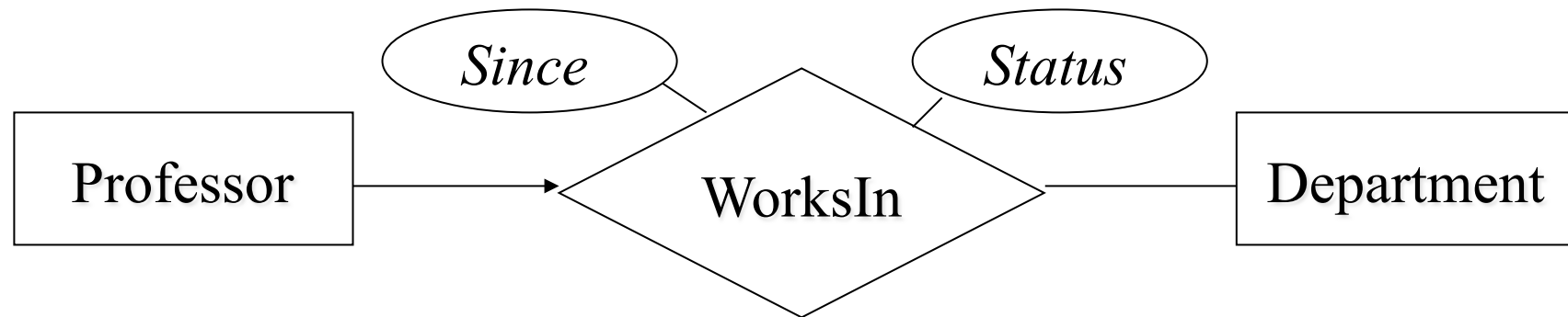


- W2006Courses (CrsCode, SectNo, *Enroll*)
- Professor (Id, *DeptId*, *Name*)
- Teaching (CrsCode, SecNo, Id, RoomNo, TAs)

Representation in SQL

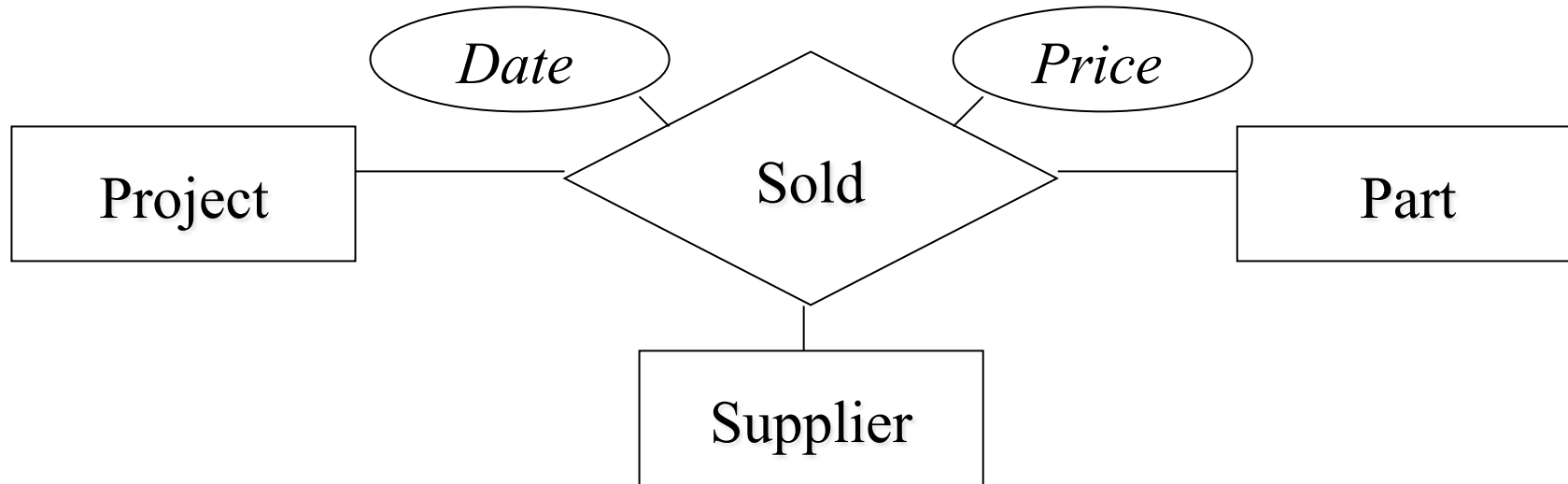
- ❑ Each role of relationship type produces a foreign key in corresponding relation
 - Foreign key references table corresponding to entity type from which role values are drawn

Example 1



```
CREATE TABLE WorksIn (  
  Since DATE,           -- attribute  
  Status CHAR (10),     -- attribute  
  ProfId INTEGER,       -- role (key of Professor)  
  DeptId CHAR (4),      -- role (key of Department)  
  PRIMARY KEY (ProfId), -- since a professor works in at most one department  
  FOREIGN KEY (ProfId) REFERENCES Professor (Id),  
  FOREIGN KEY (DeptId) REFERENCES Department )
```

Example 2



```
CREATE TABLE Sold (  
    Price INTEGER,           -- attribute  
    Date DATE,              -- attribute  
    ProjId INTEGER,         -- role  
    SupplierId INTEGER,     -- role  
    PartNumber INTEGER,     -- role  
    PRIMARY KEY (ProjId, SupplierId, PartNumber, Date),  
    FOREIGN KEY (ProjId) REFERENCES Project,  
    FOREIGN KEY (SupplierId) REFERENCES Supplier (Id),  
    FOREIGN KEY (PartNumber) REFERENCES Part (Number) )
```

Representing Participation Constraints in the Relational Model

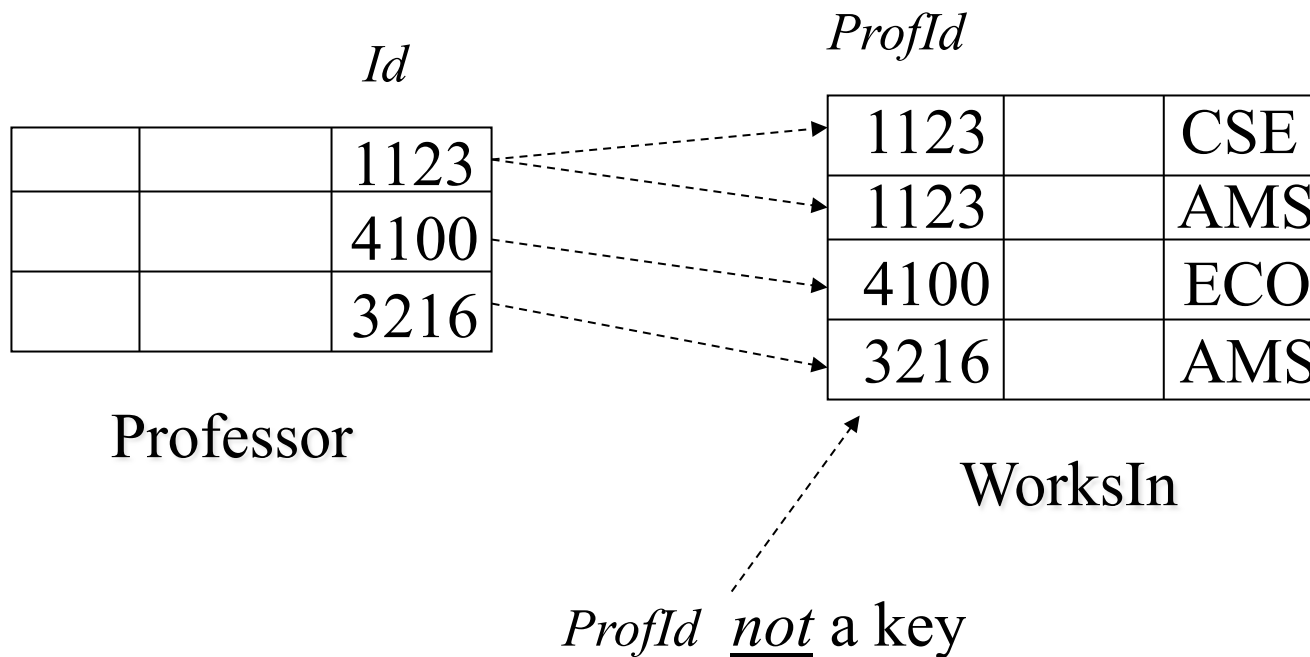


- ❑ *Inclusion dependency*: Every professor works in *at least* one dep't.
 - in the relational model: (easy)
 - Professor (*Id*) references WorksIn (*ProfId*)
 - in SQL:
 - Simple case: *If ProfId is a key in WorksIn* (i.e., every professor works in *exactly one* department) then it is easy:
 - FOREIGN KEY *Id* REFERENCES WorksIn (*ProfId*)
 - General case – *ProfId is not a key in WorksIn*, so can't use foreign key constraint (not so easy):

```
CREATE ASSERTION ProfsInDepts
CHECK ( NOT EXISTS (
  SELECT * FROM Professor P
  WHERE NOT EXISTS (
    SELECT * FROM WorksIn W
    WHERE P.Id = W.ProfId ) ) )
```

Representing Participation Constraint in the Relational Model

- ❑ Example (cannot use foreign key in Professor if ProfId is not a key in WorksIn)



Representing Participation *and* Key Constraint in SQL

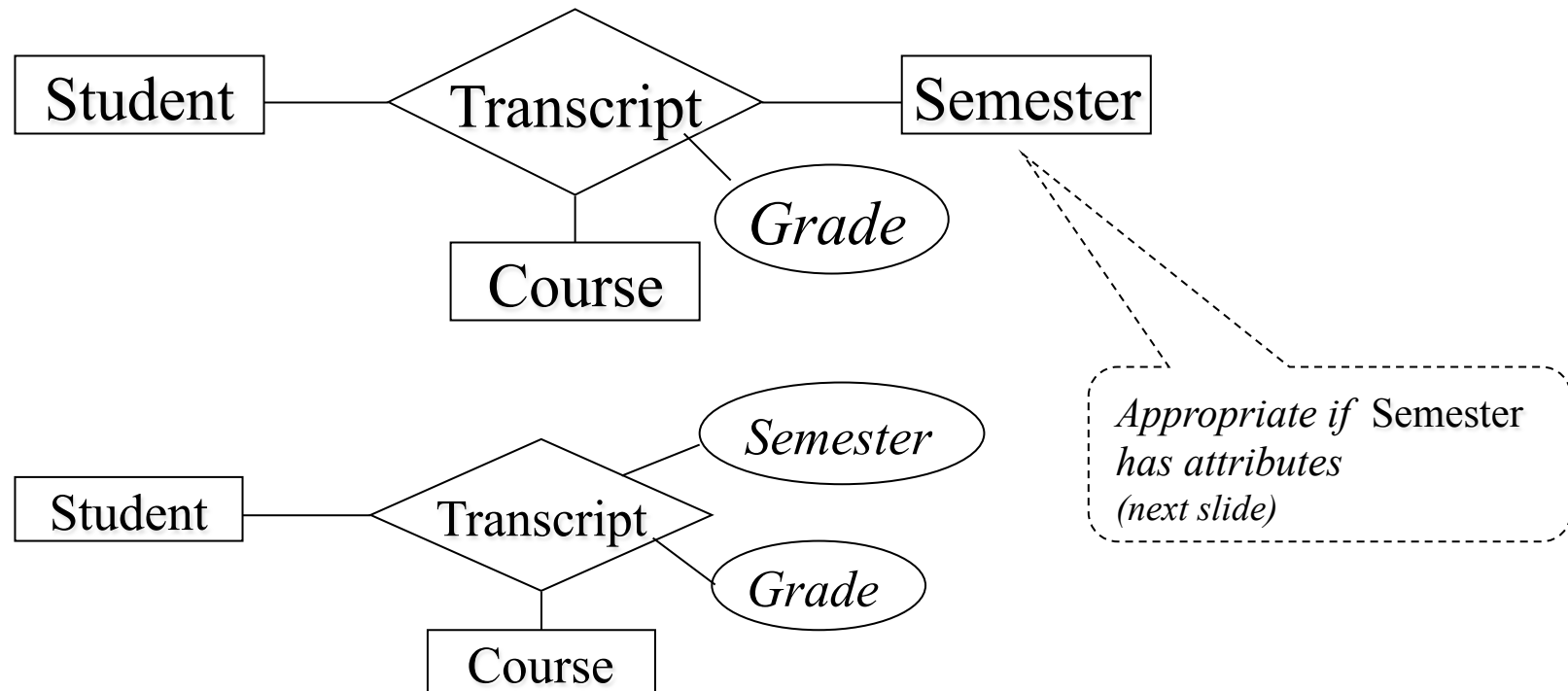
- ❑ If both participation and key constraints apply, use foreign key constraint in entity table (but beware: if a key in entity table is not primary, presence of nulls violates participation constraint).

```
CREATE TABLE Professor (  
    Id INTEGER,  
    .....  
    PRIMARY KEY (Id),      -- Id can't be null  
    FOREIGN KEY (Id) REFERENCES WorksIn (ProfId)  
                           --all professors participate  
)
```



Entity or Attribute?

- Sometimes information can be represented as either an entity or an attribute.



Entity or Relationship?

