

Vocabulary
(signature)

$(+, *)$

$(+, *, <)$

Structures

$(\mathbb{N}, +, *)$

$(\mathbb{R}, +, *)$

$(\mathbb{N}, +, *, <)$

$(\mathbb{R}, +, *, <)$

Syntax

$\exists x \exists y$

$(x = y + y)$

$\exists x \forall y$

$y > x$

$y = x$

φ

Semantics
(satisfaction)

φ is true

in $(\mathbb{N}, +, *, <)$

false in

$(\mathbb{R}, +, *, <)$

Structures over a vocabulary σ

- a universe (set) U

- an interpretation of

each function symbol f_i of arity n_i

as a function $U^{n_i} \rightarrow U$

• each constant as an element of U

• each predicate symbol P_i as

a subset of U^{m_i}

Finite strings as structures

Vocabulary

$(\underbrace{P_a, P_b}_{\text{arity}}, \underbrace{<}_{\text{arity 2}})$

alphabet Σ
 $\Sigma = \{a, b\}$

$s = abaaab$

$\mathcal{M}_s = (\{0, 1, 2, 3, 4\}, \underbrace{P_a = \{0, 2, 3\}, P_b = \{1, 4\}}_{\text{arity}}, \underbrace{<}_{\text{arity 2}})$

First Order Logic (FO)

Terms

each variable x is a term

$$FV = \{x\}$$

• each constant symbol c_i is a term

$$FV = \emptyset$$

• If f is a funct. symb of arity n and t_1, \dots, t_n are terms, then $f(t_1, \dots, t_n)$ is a term,

$$\text{and } FV = \bigcup_i FV(t_i)$$

If no function symbols, atomic flos. all terms are variable and constants

Formulas

• if t_1, \dots, t_n are terms and P a pred symb of arity n , then

$P(t_1, \dots, t_n)$ is a flos

$$FV(P(t_1, \dots, t_n)) = \bigcup_i FV(t_i)$$

• If t_1 and t_2 are terms then $t_1 = t_2$ is a flos

$$\text{and } FV = FV(t_1) \cup FV(t_2)$$

- If $\varphi, \varphi_1, \varphi_2$ are formulas
then $\varphi_1 \vee \varphi_2, \varphi_1 \wedge \varphi_2, \neg \varphi$

are formulas

$$FV(\varphi_1 \vee \varphi_2) = FV(\varphi_1) \cup FV(\varphi_2)$$

$$FV(\neg \varphi) = FV(\varphi)$$

- If φ is a f.l.c., then $\exists x \varphi, \forall x \varphi$ are f.l.c.s

$$\text{and } FV(\exists x \varphi) = FV(\forall x \varphi) = FV(\varphi) - \{x\}$$

Bound = not free

$$FV(\varphi/x, x_n) \text{ if } FV(\varphi) = \{x, x_n\}$$

f.l.c. w/o free variable = sentence

$$\psi = \exists x \underbrace{\forall y (x < y \vee x = y)}_{\varphi(x)}$$

$$(N, <) \models \psi \Leftrightarrow (N, <) \models \varphi(0)$$

for each $n \in N$ $0 < n$ or $0 = n$ is true

atomic, $\forall, \wedge, \neg, \exists$
 ~~$\forall, \wedge, \neg, \exists$~~
 \wedge, \neg, \exists

$$\forall x \varphi(x) \Leftrightarrow \neg \exists x \neg \varphi(x)$$

$$G_f \cong U^{n+1}$$

$$G_f = \left\{ \begin{array}{l} (a_0, a_1, \dots, a_n) \\ a_0 = f(a_1, \dots, a_n) \end{array} \right\}$$