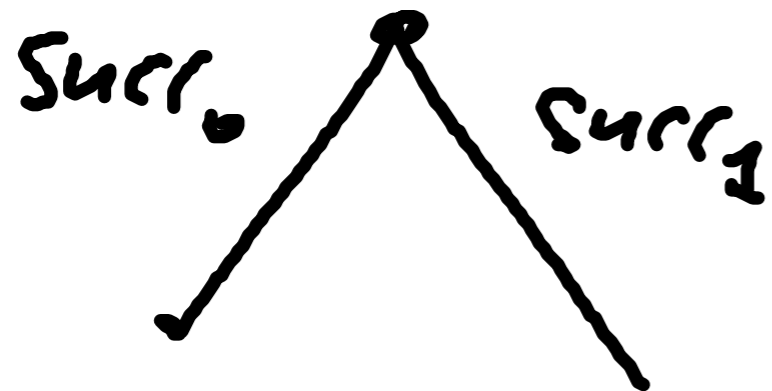


MSO on trees

$T = (D, \lambda)$, D - prefix-closed subset of $\{0, 1\}^*$
 $\lambda: D \rightarrow \Sigma$

$M_T = (D, \text{succ}_0, \text{succ}_1, P_a)$



$P_a = \{s \mid \lambda(s) = a\}$

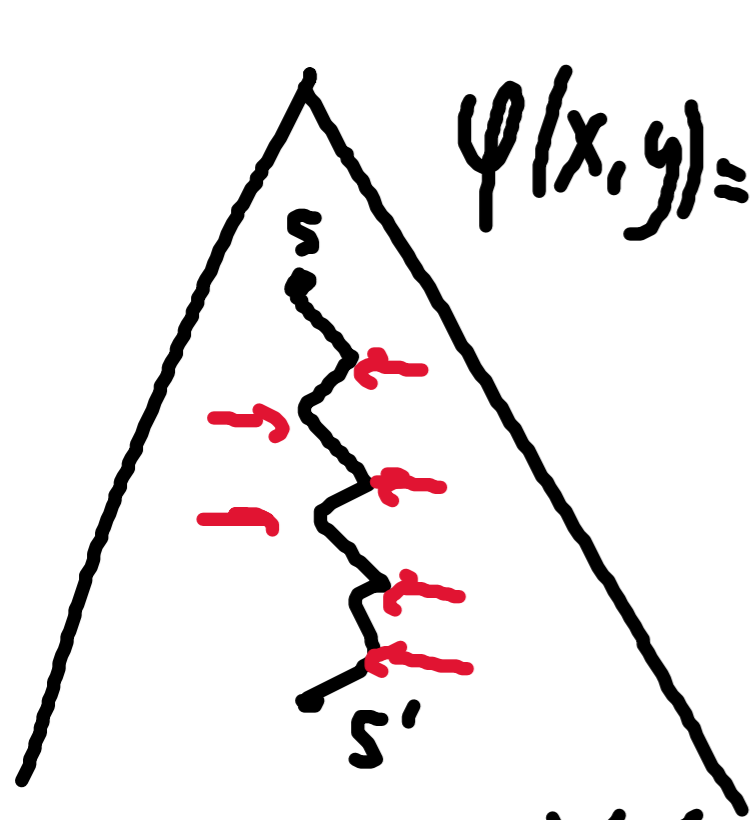
$\text{succ}_0 = \{(s, s0) \mid s, s0 \in D\}$

$\text{succ}_1 = \{(s, s1) \mid s, s1 \in D\}$

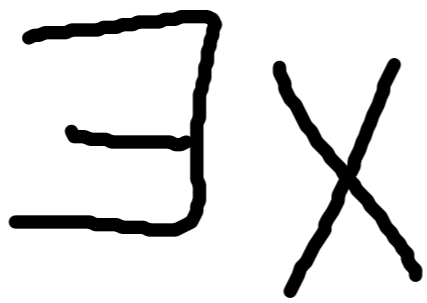
\preceq - descendant relation

$S \preceq S'$

S is a prefix of S'



$\varphi(x, y) =$



(X is a path from x to y)

FO formula

$\Pi(X, x, y)$

$\forall z, u (succ(u, z) \wedge z \neq x \wedge X(z) \rightarrow X(u))$

$X(x) \wedge X(y)$

$\forall z (succ(y, z) \rightarrow \neg X(z))$

$\forall z, u, v (X(z) \wedge succ_0(z, u) \wedge succ_1(z, v) \rightarrow$

$(X(u) \leftrightarrow X(v))$

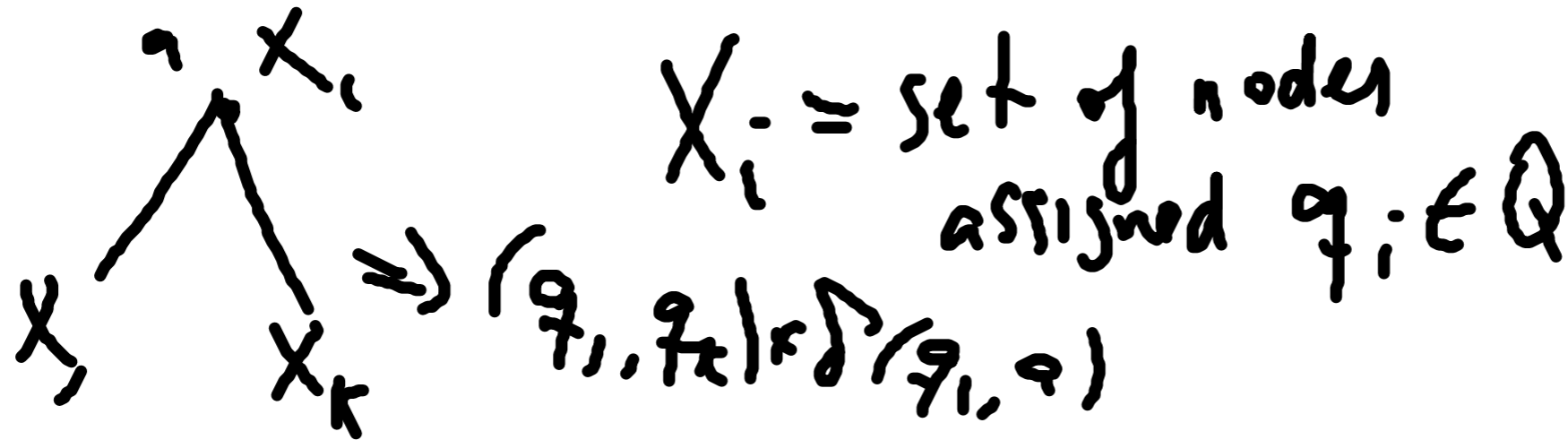
$\Pi(X, x, y) =$

Theorem (Thatcher / Wright 1969-1970)

A set of trees is MSO-definable iff
it is regular

Easy regular \Rightarrow MSO (code automata)

$A = (Q, q_0, F, \delta)$ $\exists X_0 \dots X_{n-1}$ (encode an accepting run)
 $|Q| = n$



MSO \Rightarrow automata

\hookrightarrow MSO²_{pre}

only 2nd order
variables

$\varphi(X_1, \dots, X_k)$

Succ(X)

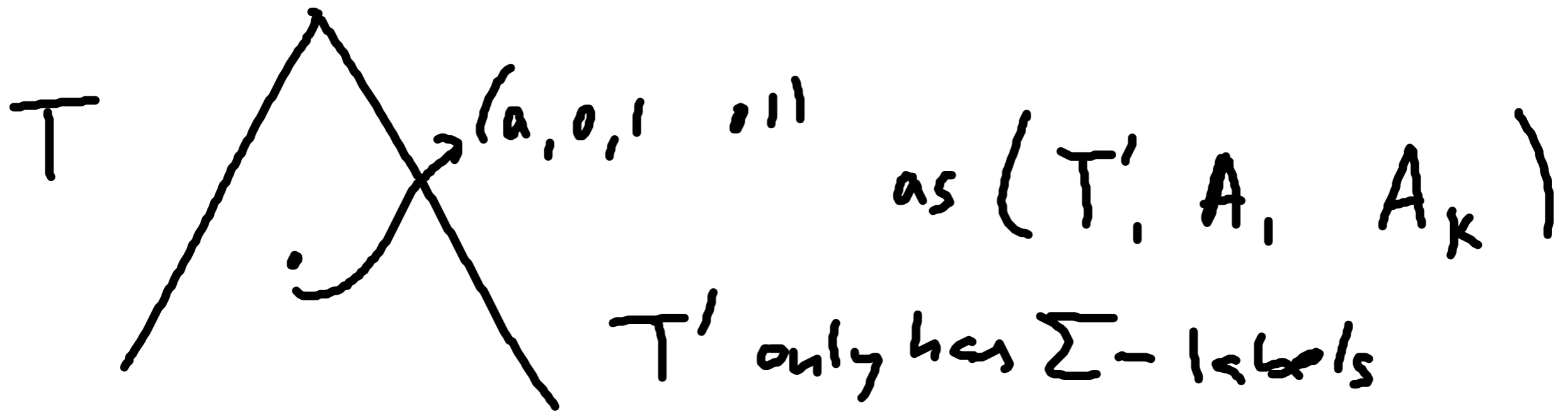
$X \subseteq Y \quad X \subseteq P_a, a \in \Sigma$

Succ_i(X, Y)

$X = \{x\} \quad Y = \{y\}$

$(x, y) \in \text{succ}_i, i: 0, 1$

$\varphi(x_1 \dots x_k) \mapsto A_\varphi$ over $\Sigma \times \{0,1\}^k$



T' only has Σ -labels

$A_i = \{s \mid \lambda_T = (a, 0, 1, \dots, s)\}$

A_φ accepts $T \iff T' \models \varphi(A_1, A_k)$ (it's position)

Over strings/trees, MSO = regular

FO¹

Thm A language is definable in FO

\iff it is star-free

star-free given by a reg expr built from

$\emptyset, a, a \in \Sigma, e_1 \cdot e_2, e_1 \cup e_2, \overline{e}$

$\Sigma = \{a, b\}$ $\overline{\emptyset} = \{a, b\}^*$ start with a $a \cdot \{a, b\}^*$

$(ab)^*$ is star-free. end with b $\{a, b\}^* \cdot b$

no aa $\overline{\emptyset} \cdot a \cdot a \cdot \overline{\emptyset}$

no bb