

Example 1

SKS - theory of k successors

$$\Sigma = \{0, \dots, k-1\}$$

$$T_k = (\Sigma^*, \text{succ}_0, \dots, \text{succ}_{k-1})$$

$$\text{succ}_i = \{(s, s+1)\}$$

SKS is decidable because it can be interpreted
in S2S

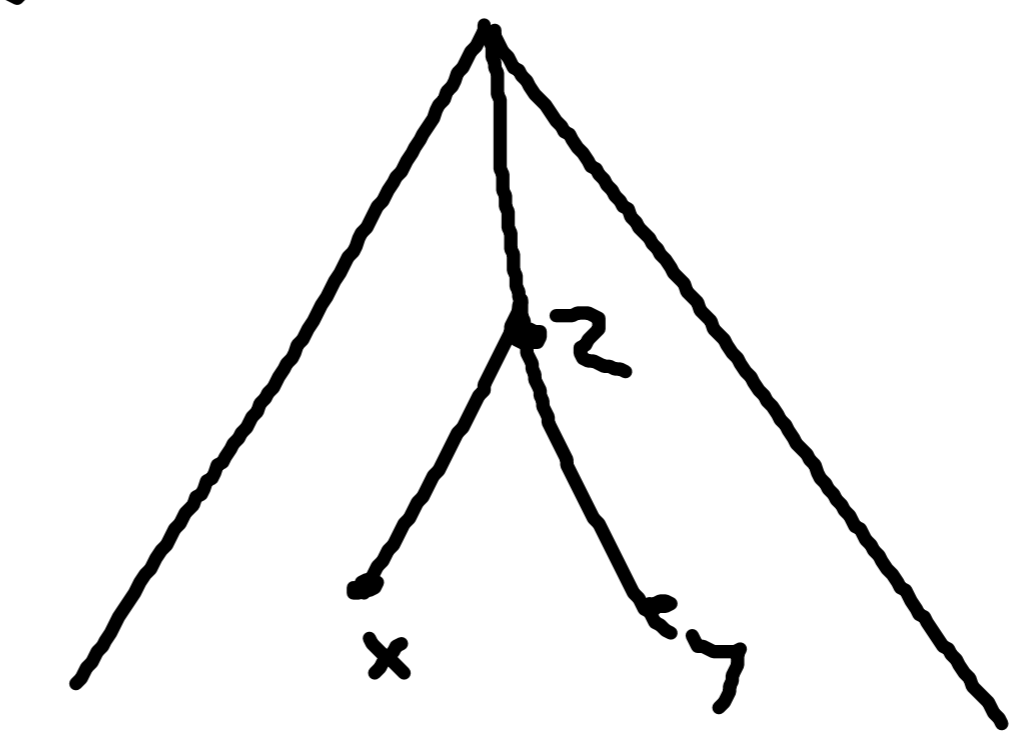
Expressible over T_2

- $x < y$ (y is a descendant of x)

- $x \cap y = z$

$z < x \wedge$

$z < y \wedge$



$(\text{succ}_0(z) < x \wedge \text{succ}_1(z) < y)$

$(\text{succ}_0(z) < y \wedge \text{succ}_1(z) < x)$

- $x <_{\text{lex}} y \Leftrightarrow$

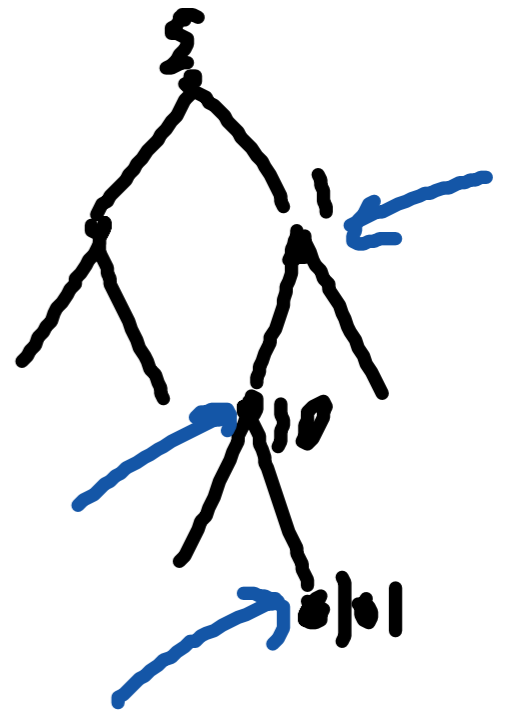
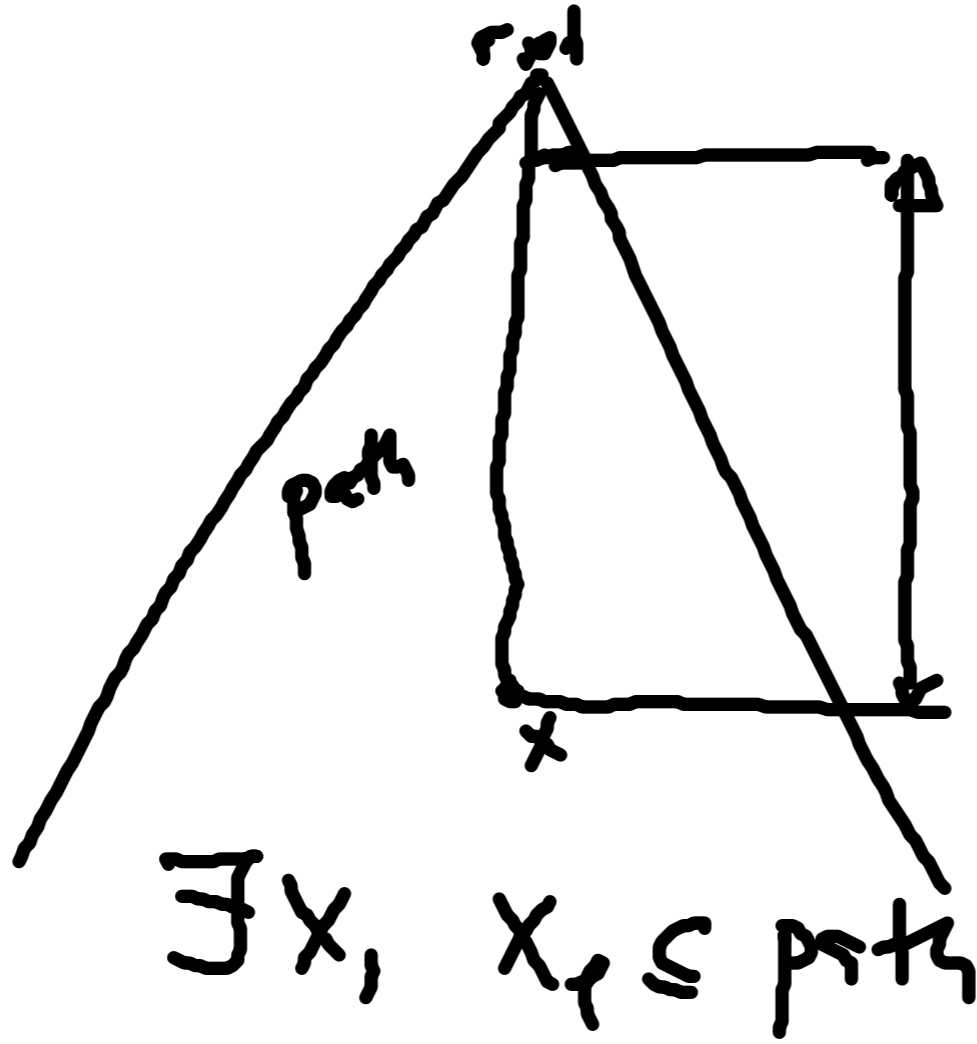
$x < y \vee \exists z (z < x \wedge z < y \wedge \text{succ}_0(z) < x \wedge \text{succ}_1(z) < y)$

L is a regular language over $\{0, 1\}^*$

$d_L(x)$

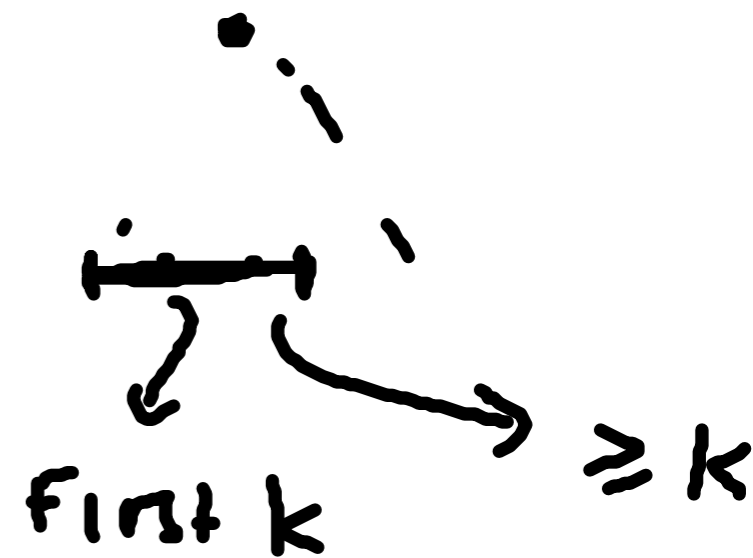
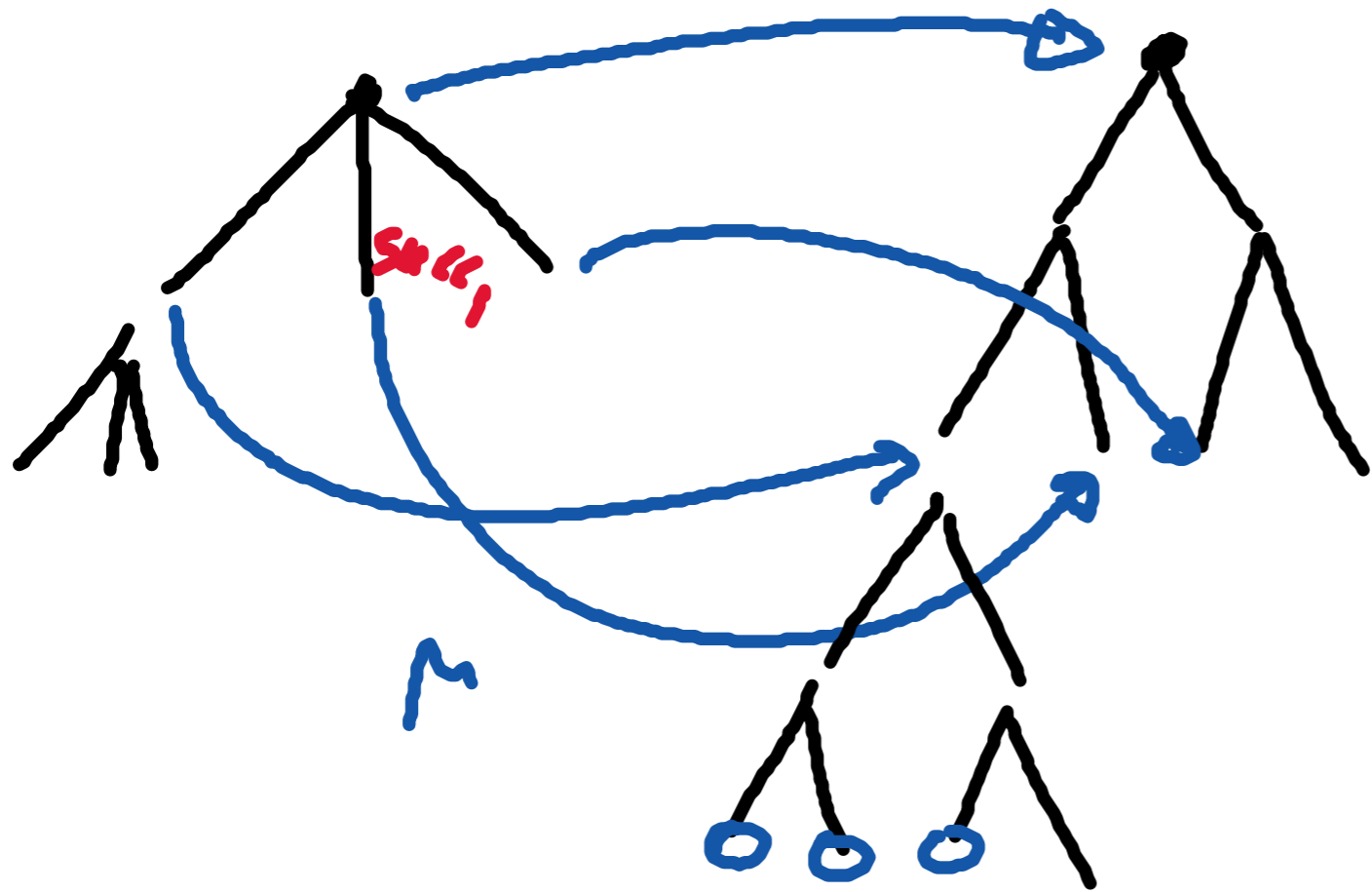
$x \in L$

L
 $\exists x_1, x_2 \beta$



β (in place of p_0 left side, in place of p_1 right side)

Interpreting S3S in S2S

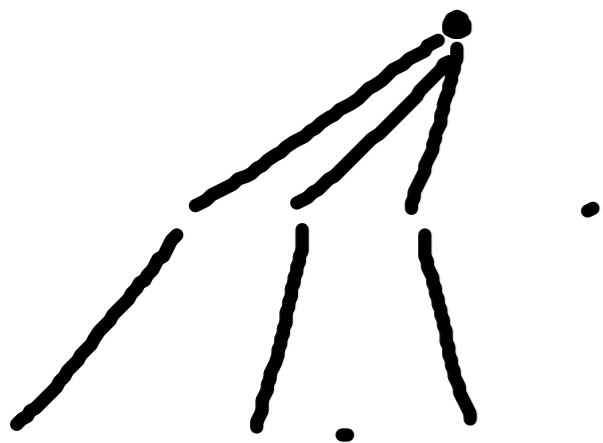


$\mu \models \mu$. $\alpha_{(\infty)^*}^*(x) \wedge \neg \exists y z \begin{matrix} y = \text{succ}_1(z) \wedge \\ x = \text{succ}_1(y) \end{matrix}$

$\Psi_{\text{succ}_1}(x, y) = \exists z \quad z = \text{succ}_1(x) \wedge y = \text{succ}_1(z)$

SWS

$(\mathbb{N}^*, \alpha, <_{lex})$

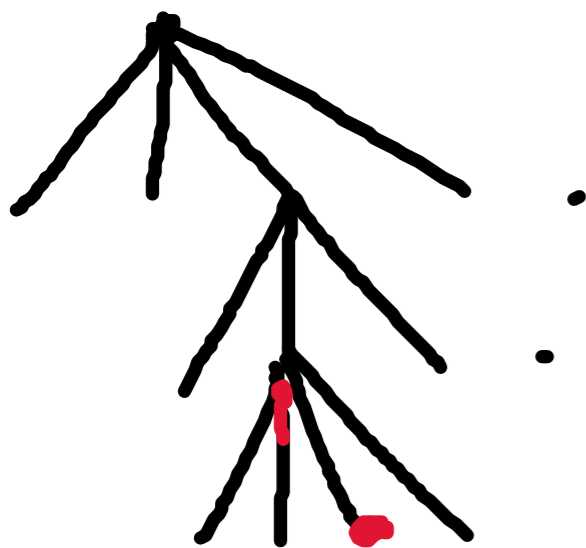


SWS is
interpretable
in S2S

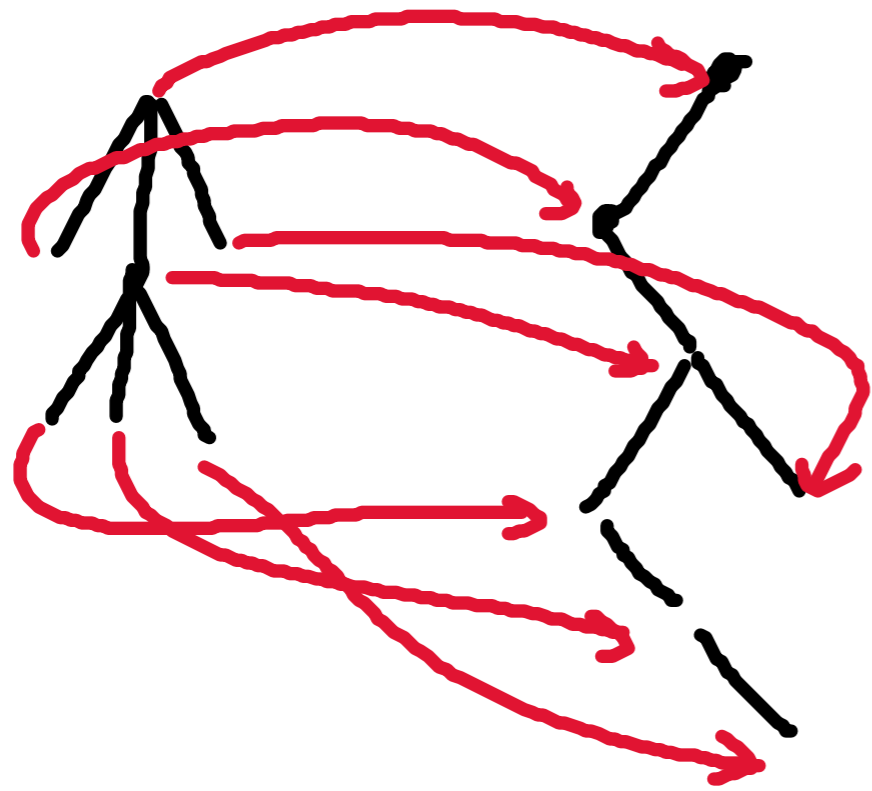


$y <_{lex} x$

\exists exactly $n-1$
of them



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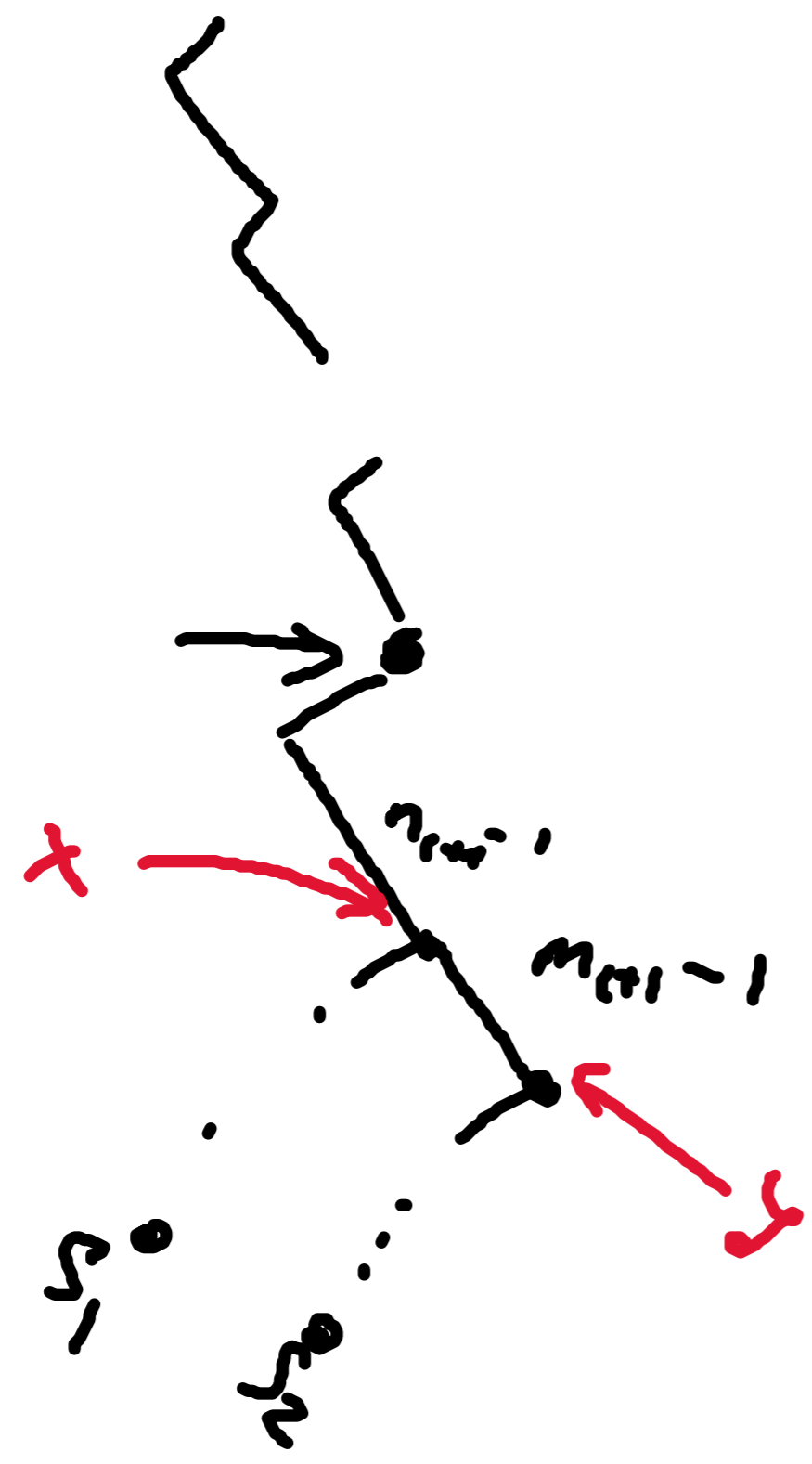
$n_0 n_1 n_2 \dots n_k \in \mathbb{N}^*$

$$I_\mu(\mu) = \varepsilon \cup 0 \cup \{0, 1\}^*$$

$<_{lex}$

<lex

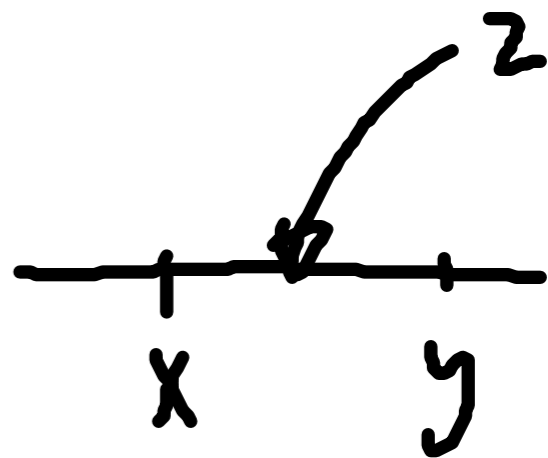
$$\frac{n_0 \cdot n_1 \dots n_k}{m_{i+1}} = S_1 \dots m_{l+1} > n_{i+1}$$



$\text{Th}_{\text{FO}}(\mathbb{Q}, <)$

$\text{Th}_{\text{MSO}}(\mathbb{Q}, <)$

$(\mathbb{Q}, <)$ $\xrightarrow{\text{interpret}}$ S2S



$(\{0, 1\}^{\mathbb{N}}, <_{\text{lex}})$ \cong $(\mathbb{Q}, <)$

There is exactly one countable dense linear ordering without endpoints

$\{0, 1\}^*$, $<_{lex}$

is dense and has no
min and no max elt

0101 1
↑
0

$\Sigma_0 1 <_{lex} \Sigma_1 1$

