

# Local Perspectives on Actions

Michael P. Fourman

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## Abstract

Giving an account of agents acting in the world—sensing, planning, communicating, doing—requires a coordinated account of, at least, three different kinds of action: ontic, epistemic, and communicative, which focus, respectively, on fact, knowledge and communication.

In this note we are concerned primarily with ontic actions. The motivating example is the STRIPS approach to the frame problem, where actions are restricted to change only specified fluents.

We present an algebraic setting for ontic actions modeled as relations describing controlled state change. We start from a standard model that encodes STRIPS updates to address the frame problem in logical terms. This model is a *system*, in the sense of Resende and Baltag.

We describe a structure that introduces a notion of local perspective or experiment. This provides a novel treatment of causal relations (which are closely related to integrity constraints and domain axioms in the AI-planning literature).

We show how this local structure arises naturally from the semantic structure of the set of possible states, and suggest that it may also help in modeling agents with different perspectives.

## 1 Introduction

Algebras of actions can be traced back to Tarski's calculus of relations [26], and beyond (see [22]). This provided the foundation for fifty years of fertile study (see, for example, [21, 10, 22, 14]). More recently, there has been revived interest, placing these algebraic ideas in a more abstract setting, and extending them—for example, to encompass semantics of knowledge and give an account of epistemic actions [23, 2].

Meanwhile, the AI planning community has worked on syntax and semantics of languages for specifying planning domains and actions, and on algorithms for solving for generating plans—sequences of actions that lead from presumed initial (or pre-)conditions to desired outcomes, or post-conditions.

The *frame-problem* [24] has been a recurrent issue for practical planning. Early syntactic approaches are now superseded, or subsumed, by more robust semantic accounts. However, these assume that all fluents are independent.

Inspired by [9], [20], and [2], we introduce a local structure on the set of fluents that can be derived from dependencies inherent in the world model, or can be established beforehand to impose such dependencies.

## 2 Properties and Actions

We consider *systems* in the sense of Resende [23] (see also Baltag *et al.* [3, 1, 2, 4]). A complete  $\bigvee$ -lattice of *properties*,  $\varphi, \psi, \dots$ , is operated on by a quantale of *actions*,  $\alpha, \beta, \dots$

$$(\varphi \bullet \alpha) \bullet \beta = \varphi \bullet (\alpha ; \beta) \quad (1)$$

$$\left(\bigvee_i \varphi_i\right) \bullet \alpha = \bigvee_i (\varphi_i \bullet \alpha) \quad \varphi \bullet \bigvee_i \alpha_i = \bigvee_i (\varphi \bullet \alpha_i) \quad (2)$$

## 3 The Standard Model

Let  $I$  be a set of *state-variables*, with each  $x \in I$  ranging over a non-empty set  $\mathcal{S}_x$  of possible values. States are *valuations*,  $\mathbf{v}$ —functions that assign a value to each state-variable. So, valuations are elements of the product *state-space*,  $\mathcal{S}$ :

$$\mathbf{v} \in \mathcal{S} = \prod_{x \in I} \mathcal{S}_x \quad \mathbf{v} = \langle \mathbf{v}_x \in \mathcal{S}_x \mid x \in I \rangle \quad (3)$$

We often consider only a subset  $\mathcal{P} \subseteq \mathcal{S}$  of states. A *property*,  $\varphi \subseteq \mathcal{S}$ , is a subset of the state space. ( $\mathcal{P}$  will be an invariant property.) The set of properties is a complete lattice—indeed, in this case, a complete boolean algebra.

A many-sorted first-order predicate calculus with a sort corresponding to each state variable provides a convenient language for defining properties. For any first-order formula  $\theta$  whose free variables are contained in  $I$ , a standard definition of satisfaction defines an associated property,  $\llbracket \theta \rrbracket$ :<sup>1</sup>

$$\llbracket \theta \rrbracket = \{ \mathbf{v} \mid \mathbf{v} \models \theta \} \quad (4)$$

A non-deterministic *action* is represented by a relation  $\alpha \subseteq \mathcal{S} \times \mathcal{S}$  from states to states. The relationship,  $x \alpha y$ , has an operational reading as, “*doing  $\alpha$  in state  $x$  may lead to state  $y$* ”. Actions, composed sequentially by relation composition, and ordered by inclusion, form a *quantale*.

An alternative, but isomorphic, representation of the quantale of actions, as sup-preserving functions  $\wp(\mathcal{S}) \rightarrow \wp(\mathcal{S})$ , equipped with function composition and point-wise ordering, has a logical reading: the result,  $\varphi \bullet \alpha$ , is the strongest post-condition including all states that *may be reached by doing  $\alpha$  in a state satisfying  $\varphi$* .

<sup>1</sup>Context is usually sufficient to distinguish  $\theta$  from  $\llbracket \theta \rrbracket$ , so we may abuse notation by omitting the *Scott brackets*,  $\llbracket \cdot \rrbracket$ . We will also let a boolean variable stand for itself:  $b \equiv b = \top$ .

We call this the *standard model*, since it has emerged as a common foundation for much work on actions and planning (sometimes explicitly, often implicitly: see, e.g. [8, 19, 17, 18, 11, 16, 6, 15, 7, 13]). In this setting, we define some basic actions.<sup>2</sup>

For every property,  $\varphi$ , we have an action  $?\varphi$ , called *require*  $\varphi$ :

$$\psi? \varphi = \psi \wedge \varphi \quad (5)$$

For every map,  $\sigma : I \rightarrow I$ , we have an action  $[\sigma(x)/x]$  (*substitute*  $\sigma(x)$  for  $x$ ):

$$\varphi[\sigma(x)/x] = \{\mathbf{w} \mid \mathbf{w} \circ \sigma \in \varphi\} \quad (6)$$

For every set of state variables,  $U \subseteq I$ , we have an action,  $\| \| U$ :

$$\psi \| \| U = \{\mathbf{w} \mid \mathbf{w} \upharpoonright U \in \psi \upharpoonright U\} \quad (7)$$

this action, *preserve*  $U$ , is the largest action that preserves the values of state variables in  $U$ . It allows other variables to change non-deterministically. We also write  $\Downarrow V$ , called *change*  $V$ , for  $\| \| (I \setminus V)$ . This action may be expressed using quantification over state variables:  $\varphi \Downarrow V \equiv \exists V. \varphi$ . A more refined version, permitting change in one direction only, can be expressed using bounded quantification.

### 3.1 Local Change

The standard solution to the frame-problem is to confine attention to actions that leave most observations unchanged.

The STRIPS representation of actions was originally defined operationally, for a setting with only boolean state variables. A state is represented by the set of *true* variables. The STRIPS action  $\alpha$ , given by a triple,  $(\mathbf{pre}, \mathbf{add}, \mathbf{delete})$ , of sets of variables, can be applied to a set  $U \subseteq I$  of state variables, iff  $\mathbf{pre} \subseteq U$ ; if applied, it produces the result  $(U \setminus \mathbf{delete}) \cup \mathbf{add}$ .

In the present setting, we can define the action corresponding to a STRIPS triple by

$$\text{STRIPS}(\mathbf{pre}, \mathbf{add}, \mathbf{delete}) = \\ (? \bigwedge \mathbf{pre}; \Downarrow (\mathbf{add} \cup \mathbf{delete}); ?(\bigwedge \mathbf{add} \wedge \neg \bigvee \mathbf{delete})) \quad (8)$$

It is straightforward to calculate that applying this action to the property corresponding to the single state in which  $U$  is the set of *true* variables, simulates the STRIPS procedure.

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<sup>2</sup>We often omit the  $\bullet$  and use simple postfix notation in place of  $-\bullet\alpha$ .

## 4 Causal Models

The treatment of domain axioms or integrity constraints remains an outstanding problem. The issue is, how to model a world where the values of state variables may be inter-related? When one variable is changed, others may require adjustment to accommodate some universal constraint. It has seemed hard to restrict the ability to change to the appropriate variables in a semantically principled way. Here, we suggest an approach to this problem that may initially appear *ad hoc*. In the following section we derive our treatment from purely semantic principles.

Rather than considering all state variables as independent atoms, we identify a structure that encodes functional dependencies. This structure places the atomic observations of fluents within a poset of perspectives, which is a complete  $\vee$ -lattice with distributive joins—a *frame*.

Starting from any poset, the downsets constitute the frame of open sets of a point-set topology on the poset. The poset is embedded, in the frame, by taking each element,  $x$  to the down-set  $\downarrow x$ , composed of elements  $y \leq x$ . This embedding preserves existing meets ( $\wedge$ ), but not joins ( $\vee$ ).

The original poset has a natural map,  $\nu$ , to any quotient of this frame (by  $\wedge$ ,  $\vee$ -congruences). Indeed, a quotient may be specified by stipulating *covering conditions*, of the form  $\nu(y) \leq \bigvee_i \nu(x_i)$  (algebraically  $\nu(y) \wedge \bigvee_i \nu(x_i) = \bigvee_i \nu(x_i)$ ).

Consider an example. Pearl [20] models causal relationships as functional dependencies, and defines a *causal model* (*op cit.* p27.) to be

a set of equations of the form

$$x_i = f_i(pa_i, u_i), \quad i = 1, \dots, n, \quad (9)$$

where  $pa_i$  (connoting *parents*) stands for the set of variables judged to be immediate causes of  $X_i$  and where the  $U_i$  represent errors (or “disturbances”) due to omitted factors.

Such a model is represented as a directed graph, with a node for each variable, and an arrow  $x \rightarrow y$  if  $x$  is a parameter of an equation defining  $y$ . In this note, we confine our attention to causal models represented by a directed acyclic graph (DAG). We take this graph as a pre-order on variables, and take the reflexive, transitive closure of  $x \rightarrow y$ , to give a partial order.

Consider a specific example: suppose that we want to include in our model a particular dependency — that  $y = f(x_1, \dots, x_n)$ . In general, if one of the  $x_i$  changes, then  $y$  may change. So, the set of fluents that may be changed by an action cannot be entirely arbitrary. Moreover,  $y$  cannot change unless at least one of the  $x_i$  does.

We define a (point-set) topology on the set of fluents such that any closed set of fluents  $a$  is a suitable candidate for changes.

First, we equip our poset with a topology in which down-sets are open. So any open set containing  $y$  contains all of the  $x_i$ . Second, we choose the topology a

coarser topology, with fewer open sets, generated by letting the family  $\langle x_i \rightarrow y_i \mid i = 1, \dots, n \rangle$  cover  $y$ . In other words, we enforce the condition that any open set containing all of the  $x_i$  contains  $y$ .

A closed set contains  $y$  iff it contains at least one of the  $x_i$ ; We allow changes to a closed set—equivalently, we allow the action  $|||U$  only where  $U$  is an open set.

#### 4.1 Grothendieck Topology

Let  $P$  be a poset, and  $\nu$  a  $\wedge$ -preserving map from  $P$  to a frame. For  $y \in P$ , we say a family  $x_i \leq y$   $\nu$ -covers  $y$  iff  $\nu(y) \leq \bigvee_i \nu(x_i)$ . The collection of  $\nu$ -covers is a *Grothendieck topology* on the poset. Every Grothendieck topology arises in this way—we can construct a suitable frame, and the map,  $\nu$ , from the topology. An intersection of Grothendieck topologies is a Grothendieck topology—since a product of frames is a frame. The coarse topology, in which every family is covering, is a Grothendieck topology—since the singleton lattice is a frame. So any collection of covering families generates a Grothendieck topology.

For details of Grothendieck topologies on posets, see [25], and, for a more general account, [12].

### 5 Canonical Causal Dependencies

We revisit the example of the previous section. Instead of imposing a partial order and topology, we derive these from the semantic structure. To do this we introduce new structure, conservatively extending the standard model. In the next section we introduce an abstract presentation of this extension, which is in general substantive.

We start from the standard model, but consider only those valuations that satisfy the equations of the causal model:  $\mathcal{P} = \{\mathbf{v} \mid \mathbf{v} \text{ satisfies the causal constraints.}\}$

We introduce a family  $U \in \mathcal{O}$  of *perspectives*—these may be thought of as experiments. In the standard model, these are just sets of variables,  $U \subseteq I$ . The intuition is that these “perspectives” represent a local focus on some aspects of the world, an isolated environment within which we may experiment independently of outside influences. We write  $\mathcal{P}(U)$  for the local states that may be hypothesized from this perspective  $U$ . In the standard model these are just the restrictions of global states to  $U$ .

Perspectives are partially-ordered:  $V \rightarrow U$  if  $V$  represents a narrower focus than  $U$ , and we require that this be evidenced by a *restriction map*  $\upharpoonright V : \mathcal{P}(U) \rightarrow \mathcal{P}(V)$ . For the standard model this order is just set inclusion, and the restriction maps are immediate.

We say a family of hypotheses  $a_i \in \mathcal{P}(U_i)$  is *coherent* iff, whenever there is a perspective  $W$  such that  $U_i \leftarrow W \rightarrow U_j$  then  $a_i \upharpoonright W = a_j \upharpoonright W$ . A family  $U_i \rightarrow W$   $P$ -covers  $W$  iff for every coherent family  $a_i \in \mathcal{P}(U_i)$ , there is a unique hypothesis  $\bar{a} \in \mathcal{P}(W)$  such that, for all  $i$ ,  $\bar{a} \upharpoonright U_i = a_i$ . We call  $\bar{a}$  an extension of the  $a_i$ .

Now suppose we have a functional relationship  $y = f(x_1, \dots, x_n)$ . We have a inclusions  $\{x_i\} \rightarrow \{x_1, \dots, x_n, y\}$ . Any selection of values  $a_i \in \mathcal{P}(\{x_i\})$  is a coherent family and it has a unique extension  $\bar{a} = \langle a_1, \dots, a_n, f(a_1, \dots, a_n) \rangle \in \mathcal{P}(x_1, \dots, x_n, y)$ .

The collection of all covering families generates a Grothendieck topology on the poset of perspectives. The sets  $\mathcal{P}(U)$  together with restriction maps form a presheaf whose sheafification is our substitute for the presentation-derived construction of the previous section.

In the simple example with a single functional dependence,  $\{x_1, \dots, x_n, y\}$ , in the sheafification, plays the rôle of  $y$  in our earlier construction;  $\{y\}$  turns out to be covered by the empty family—and so behaves as a minimal element accommodating a trivial, singleton-set of possibilities in the sheafification.

## 6 Perspectives

Algebraically, perspectives form a partially ordered set, equipped with a Grothendieck topology; we take the state space  $\mathcal{P}$  to be a sheaf over this site. Equivalently, let  $\mathcal{X}$  be a locale, whose open sets we call *perspectives*, and  $\mathcal{S}$  a sheaf on  $\mathcal{X}$ . We say that the space of perspectives is *irredundant* if, whenever  $U_i \mathcal{P}$ -covers  $W$  then  $\bigvee_i U_i = W$ .

Global states arise as global sections of  $\mathcal{P}$ —which correspond to consistent sets of hypotheses from each local perspective. Each perspective  $U$  induces an equivalence relation on global states:  $\mathbf{v} \equiv \mathbf{w}$  iff  $\mathbf{v} \upharpoonright U = \mathbf{w} \upharpoonright U$  these states appear equivalent from this perspective. There may also be states, elements of  $\mathcal{P}(U)$ , that appear possible from this perspective, but that do not extend to global states. It may be interesting to link the appearance maps of [1] to these structures.

## 7 Conclusions

This is a preliminary report on ongoing work.

We have shown how a causal model, in the sense of Pearl, can arise naturally from semantic analysis of a set of possible states. This leads to an account of local perspectives as the open sets of a space of observations.

The structures introduced to this end suggest the beginnings of an account of local perspective, hypothesis and experiment.

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