Introduction Dynamic ACL Semantics Conclusions

Dynamic Semantics for Agent Communication Languages

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What this paper is about

- We present formal framework for dynamic semantics (DS) for agent communication languages (ACLs)
- A DS describes how the meaning of utterances changes depending on previous behaviour of agents
- Example:
 - ► "A makes a promise to B to perform X" initially means: "A intends to perform X and intends that B believes this"
 - If A doesn't perform X, it will come to mean "A does not intend to perform X but intends that B believes she does"
- Fairly complex formal framework, bottom line: DS consists of different states of ACL semantics with transitions contingent on commitment stores and agents' actions

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What for?

- A means of describing how agents will respond to what agents do depending on what has been said
- ► A tool for tracking the evolution of semantics in an adaptive sense
- A way to exploit *communication-inherent* sanctioning and rewarding mechanisms that work in open systems
- A mechanism that allows for linking communication to practical reasoning from the agent's point of view
- Not so much a contribution to agent communication language research, but more about social reasoning in open systems

"Static" ACL semantics

- Standard, "static" ACL semantics come in two flavours:
 - mentalistic: linked to agent reasoning but not verifiable
 - social (commitment-based): verifiable but not grounded in agent reasoning
- We present DS for commitment-based semantics, but have to fix the grounding issue first
- DS for mentalistic ACL semantics are also possible (in the sense of tracking "ostensible mental states")

Commitment-based semantics

► The Fornara/Colombetti approach to commitments:



- Solid lines indicate state transitions brought about by agents, dashed lines stand for transitions caused by external events
- ACL semantics are defined based on state transitions (e.g. request creates unset commitment, accept makes it pending, reject cancels it)

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Commitments

▶ We use the following notation for commitments:

$$\langle \iota, \mathbf{s} : \chi \oplus \varphi \ominus \psi \rangle_t^{i o j}$$

where

- ι identifier, s state (unset, pending, active, violated, fulfilled)
- $\blacktriangleright~\chi$ is the debitum, φ/ψ are activation/deactivation conditions,
- i/j is the debtor/creditor, t timestamp of transition to s
- χ, φ, ψ are taken from some propositional logical language (we often abbreviate Γ = χ ⊕ φ ⊖ ψ)
- Example:

 $\langle x, \mathbf{v} : received(5, \$500) \oplus received(3, toys) \ominus returned(3, toys)
angle_{12}^{3
ightarrow 5}$

Semantics of commitments (informal version)

- Apply the following transition rules after each step to a commitment store CS:
 - D: cancel any $\langle \iota, *: \chi \oplus \varphi \ominus \psi \rangle_t$ for which ψ occurs
 - A: make any $\langle \iota, \mathbf{p}: \chi \oplus \varphi \ominus \psi \rangle$ active for which φ holds
 - S: make any $\langle \iota, \mathbf{a} : \chi \oplus \varphi \ominus \psi \rangle$ fulfilled if χ becomes true (serendipity)
 - F/V: make any $\langle \iota, \mathbf{a} : \chi \oplus \varphi \ominus \psi \rangle_{t-1}^{i \to j}$, fulfilled/violated if $Done(i, a) \land causes(a, \chi)$ is true/false
- Similar to Fornara and Colombetti's operational view
- But how about grounding, i.e. what do these commitments actually mean in terms of agent behaviour?

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Grounding (informal version)

Basic grounding rule: The behaviour of agent *i* is said to be compliant with CS iff

any commitment in CS that becomes active is immediately fulfilled by the respective debtor

- ▶ Assume system defined in terms of runs $r = e_1 \stackrel{\vec{a}_1}{\rightarrow} \dots \stackrel{\vec{a}_{t-1}}{\rightarrow} e_t \in \mathcal{R}$ with environment states e_i and joint actions \vec{a}_i
- Agent behaviour given by functions $g_i : \mathcal{R} \to Ac_i$
- We can define

 $compliant(CS) := the set of agent functions for i that will always execute an action that causes <math>\chi$ if a commitment became pending for i during run r in any "extension" r' of r in which that commitment becomes active

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Compliant vs. expected behaviour

- This tells us what agents *ought* to do, but not what we *expect* them to do
- Introduce second type of commitments called *expectations* that override other commitments
- ▶ Distinguish from "normative" commitments by using round brackets (ι, s : Γ)^{i→j}_t, same semantics in terms of processing rules

Define

$$\begin{bmatrix} CS \end{bmatrix} := \{ \langle \iota, \mathbf{s} : \Gamma \rangle \in CS | \mathbf{s} \in \{ \mathbf{u}, \mathbf{p}, \mathbf{a}, \mathbf{f}, \mathbf{v} \} \} \\ \lfloor CS \rfloor := \{ (\iota, \mathbf{s} : \Gamma) \in CS | (\iota, \mathbf{s} : \Gamma) \in CS, \\ \langle \iota, \mathbf{s}' : \Gamma \rangle \in CS, \mathbf{s}, \mathbf{s}' \in \{ \mathbf{u}, \mathbf{p}, \mathbf{a}, \mathbf{f}, \mathbf{v} \} \}$$

• $compliant(\lceil CS \rceil)$ expresses what agents are supposed to do

Compliant vs. expected behaviour

Expected behaviour defined as

 $expected(CS) := compliant(\lfloor CS \rfloor)$

i.e. behaviour that adheres to expectations where such expectations exist (and is compliant otherwise)

- Advantages of separate treatment of compliant and expected behaviour
 - 1. We can respond to "unexpected" compliant behaviour
 - 2. We can make concrete predictions about others' behaviours in different ways
- Planning perspective: different rules can be imagined, the simplest is to always predict expected behaviour

Definition of dynamic semantics

- Basic idea: provide different "versions" of an ACL semantics and switch from one to the other depending on agent behaviour
- Each version is a semantic state, i.e. a collection of definitions for the semantics of individual speech acts
- A semantic transition relation tracks the evolution of semantic states (for all possible pairs of communicating agents separately)
- As an example, assume a small fragment of an ACL with two semantic variants for accept

A minimal ACL fragment

$$\begin{array}{c|c} \hline time(t), new(\iota) \\ RQ: \boxed{request(i, j, \iota : \Gamma)} \\ CS \leftarrow CS \cup \{\langle \iota, \mathbf{u} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{u} : \Gamma \rangle_t^{i \to j}\} \\ \hline \langle \iota, \mathbf{u} : \Gamma \rangle_{t'}^{j \to i} \in CS, time(t) \\ AC: \boxed{accept(i, j, \iota : \Gamma)} \\ CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \\ \hline CS \leftarrow CS \cup \{\langle \iota, \mathbf{p} : \Gamma \rangle_t^{i \to j}\} \cup \{\langle \iota, \mathbf{c} : \Gamma \rangle_t^{i \to j}\}\} \\ \hline \end{array}$$

 Publicly verifiable, "action schemata" specified in a way that makes them directly usable for planning

Example

- Semantic states $s_0 = \{RQ, RJ, AC\}$ and $s_1 = \{RQ, RJ, AC2\}$
- Initial state is s_0 for all pairs of agents (i, j)
- Transitions depend on constraints on commitment store contents
 Example:

$$\begin{array}{c} (i, \mathbf{v} : \Gamma)^{i \to j} \in CS : \ \{(i, *)\} \cup \{(j, i)\} \\ \hline \\ (s_0) \\ \hline \\ \forall \langle \iota, \mathbf{v} : \Gamma \rangle_t^{i \to j} \in CS \ \exists \langle \iota, \mathbf{f} : \Gamma' \rangle_{t'}^{i \to j'} \in CS.t' > t : \ \{(i, *)\} \end{array}$$

• Semantics: for every transition $s \xrightarrow{c} s'$ "move" all pairs currently in s for which c applies to state s'

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Summary

- Proposed a mechanism for adapting the semantics of speech acts depending on observed behaviour and previous utterances
- This version based on commitment-based approach (but working on application to mentalistic version)
- Provided grounding for commitment-based semantics based on actual expectations about future behaviour ("compliance")
- Introduced distinction between (normative) commitments and (predictive) expectations and defined relationship between them
- Described semantics of DS framework using a state transition system approach



Conclusions

- Our work allows for the definition of communication-inherent sanctioning and rewarding mechanisms
- Paper discusses desiderata for DS (e.g. respect for commitment autonomy, avoiding commitment inconsistencies, unprejudiced judgement, convergence, forgiveness, equality)
- Evaluation of advantages (contingent on agent reasoning about communication, of course) yielded good results
- Research perspectives: finding appropriate DSs of common ACLs for classes of application domains
- Currently we are focussing on applying framework to strategic lying using mentalistic semantics

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The End

Thank you for your attention!

