Towards a Principle of Least Surprise for bidirectional transformations

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Plan

- Why do we need such a Principle?
- Optimality: Least Change via metrics
- Why isn’t that just the right answer?
- Reasonableness: Least Surprise
- What next?
Model-driven development

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Motivation
Manage information overload, by separating concerns and providing representations suitable for each set of decisions.
## Model-driven development

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Least Surprise is the whole point of MDD.

So we’d better get it right.
Informal idea

Meertens’ formulation

The action taken by the maintainer of a constraint after a violation should change no more than is needed to restore the constraint.

Extreme case is easy to formalise (and, usually, agree on):

Hippocraticness

If nothing needs to be changed, change nothing.
Let $X$ be how much we need to change to restore the constraint.
Let $Y$ be how much the $bx$ actually changes.
Require $Y \leq X$.
Oh, so $Y = X$.
Now: how do we measure $X$ and $Y$?
Given $m$ and $m'$, need to be able to say how much changed from $m$ to $m'$.

A sensible way to do this is called a metric.

**Definition**

A metric on $M$ is $d : M \times M \to \mathbb{R}$ satisfying

- $d(m, m') = 0$ iff $m = m'$
- $d(m, m') = d(m', m)$
- $d(m, m') + d(m', m'') \geq d(m, m'')$

E.g. minimal edit distance (typically).
Metric least change

Assume given:

- a bxi $\langle R, \vec{R}, \vec{R} \rangle : M \leftrightarrow N$;
- metrics $d_M$, $d_N$ on $M$ and $N$.

Then

**Definition**

$R$ is metric-least if for all $m \in M$ and for all $n, n' \in N$

$$R(m, n') \Rightarrow d_N(n, n') \geq d_N(n, \vec{R}(m, n))$$

and dually.
Implementing metric least change

Macedo and Cunha implemented this approach, providing a new semantics to the syntax of QVT-R.

Given consistency relation $R$, and models $m$, $n$, with the task to find a new $n'$ consistent with $m$:

1. set $d = 0$;
2. search exhaustively for consistent $n' \in N$ at distance $d$ from $n$;
3. if there are any, present them all to the user as options;
4. else, increment $d$ and goto 2.
Problems with metric least change

2. Usability: no easy way to ensure unique closest model, or to allow bx programmer to specify the choice.
3. Doesn’t compose (even where that makes sense).
4. Inflexibility (+/−?).

Can mitigate inflexibility by allowing bx programmer to choose metrics: but hmm.
Example where inflexibility bites

\[ M = \text{UML models} \]
\[ N = \text{test suite (in JUnit say)} \]

\[ R(m, n) \text{ iff every } \langle\langle \text{persistent} \rangle\rangle \text{ class in } m \text{ has a test class of the same name in } n, \text{ containing an “appropriate” set of tests for each public operation...} \]
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You modify the test class for a \( \langle \text{persistent} \rangle \) class: \( \text{int} \to \text{long} \) throughout.

Propagate back to UML model. What happens?

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- \text{int} \to \text{long} throughout?
- remove \langle \text{persistent} \rangle?
Mitigation attempt 1

Use a different metric: make removing the ⟨⟨persistent⟩⟩ be seen as an enormously expensive change.

But now we have a metric defined specifically for this bx.
Mitigation attempt 2

Why do we feel that removing «persistent» is not as cheap as it looks?

Because it disrupts the connection between the two models: breaks links.

It’s a small change to the model, but a large change to the witness structure explaining their consistency?
## Witness structures

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<td>A witness structure is a structure whose intention is to capture something about the relationship between the models being kept consistent.</td>
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- **Relational bx:** “they’re consistent because I say so”
- **TGGs:** “they’re consistent because this is how they are built up together using the TGG rules”
- **MDD with trace links:** “they’re consistent because this part of this links to that part of that”
- **QVT-R game:** “they’re consistent because here’s how Verifier wins a consistency game on them”
- **Or even:** “they’re consistent because Coq gave me this proof that they are”
So metric least change may yet be useful in some settings, or when its deficiencies can be mitigated as mentioned. But it does not deserve to be seen as the sole Right Answer. Back to the drawing board.
Principle of Least Surprise

or Least Astonishment, for user interface design, adapted to language design, etc. E.g. Saltzer and Kaashoek, via Wikipedia

People are part of the system. The design should match the user’s experience, expectations, and mental models.
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People are part of the system. The design should match the user’s experience, expectations, and mental models.

The essential point is not that user should always know what will be on the next screen.

Rather, that they should seldom be surprised, i.e., taken aback.

Despite “least”, we really seek reasonable, rather than optimal, behaviour.
Keep me informed, OK? No nasty surprises!

Yessir!
Hmm. I have to change my model.

How much change can Boss take? 

Better inform him before I cause his model to change.

So like, before my model changes by $\Delta$?
Continuity

**Definition**

$f : S \rightarrow T$ is continuous at $s$ iff

$$\forall \epsilon > 0 . \exists \delta > 0 . \forall s' . d_S(s, s') < \delta \Rightarrow d_T(f(s), f(s')) < \epsilon$$

We say just “$f$ is continuous” if it is continuous at all $s$. 
Continuity of $\overrightarrow{R}$ at $(m, n)$

Definition

$\overrightarrow{R}$ is continuous at $(m, n)$ iff

$$\forall \epsilon > 0 . \exists \delta > 0 . \forall m' . d_M(m, m') < \delta \Rightarrow d_N(\overrightarrow{R}(m, n), \overrightarrow{R}(m', n)) < \epsilon$$

i.e. iff $\overrightarrow{R}(-, n) : M \to N$ is continuous at $m$. 
Where do we want continuity?

Well, at the points where we want guarantees... e.g.

- Everywhere: strong continuity
- Only at consistent \((m, n)\): weak continuity
- Something else? Where we can get it? On “good” subspace pair?

In the history ignorant (vwb, \textsc{PutPut}) case, strong = weak.

+ Strongly (rsp. weakly) continuous bx compose, where composition makes sense.
On good subspace pairs?
The problem with continuity

- Every function is continuous at an isolated point :-(

Variants with better overall behaviour? See paper. Hölder continuity quite promising.

Prediction: way too strong to expect it everywhere in realistic bx, but identifying subspace pairs on which a bx is Hölder continuous may be worthwhile.
Other approaches/Future work

category theory
topology
dependent type theory (HoTT??)
subspaces
partiality
heuristics?
Lagrangian/Hamiltonian mechanics?
simulated annealing?
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