

Lecture 11 February 18, 2002

















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H	uffman Coding			
Huffn	fman coding is an optimal, lossless compression scheme			
	ery pixel is mapped onto a variable-length bit string ording to a probability table, as follows: For an example, I will show an example of Huffman coding digits (0-9), but for an 8-bit image you would Huffman code 0-255			
Digit	0 1 2 3 4 5 6 7 8 9			
Prob.	.1 .3 .1 .03 .1 .05 .1 .06 .05 .1			
Digit	0 1 2 4 6 7 8 9 3/5			
Prob.	.1 .3 .1 .1 .1 .06 .05 .1 .08			
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4	Huffi	nan (II)
	Digit	0 1 2 4 6 9 3/5 7/8
	Proh	
	1100.	.1 .5 .1 .1 .1 .1 .06 .11
	Digit	0 1 2 4 6 3/5/9 7/8
	Prob.	.1 .3 .1 .1 .1 .18 .11
	Digit	0 1 2 4/6 3/5/9 7/8
	Prob.	.1 .3 .1 .2 .18 .11
	Digit	1 0/2 4/6 3/5/9 7/8
	Prob.	.3 .2 .2 .18 .11
	Digit	1 0/2 4/6 3/5/7/8/9
	Prob.	.3 .2 .2 .29
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Huf	fman (III)
Digit	1 0/2/4/6 3/5/7/8/9
Prob.	.3 .4 .29
Digit	0/2/4/6 1/3/5/7/8/9
Prob.	.4 .59
Digit	0/1/2/3/4/5/6/7/8/9
Prob.	.99
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