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# Performance modelling

## with PEPA nets and PRISM

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# Outline

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- PEPA nets: informal introduction
- PEPA nets: few formal definitions
- Simple example: mobile agent
- From PEPA nets to PRISM
- Complex example: mobile IP



# PEPA nets: informal introduction

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Stochastic  
(coloured)  
Petri Nets



PEPA



PEPA nets

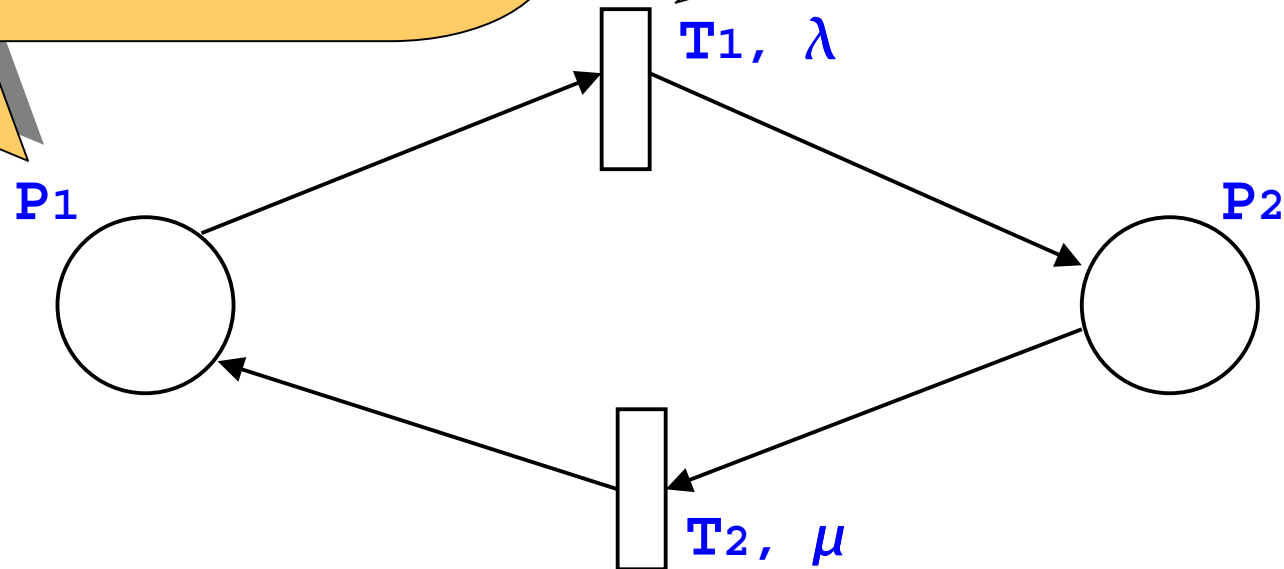


# PEPA nets: informal introduction

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The tokens are PEPA components that perform local activities and can move from one place to another ...

Transition names are labelled with activity names



# PEPA nets: informal introduction

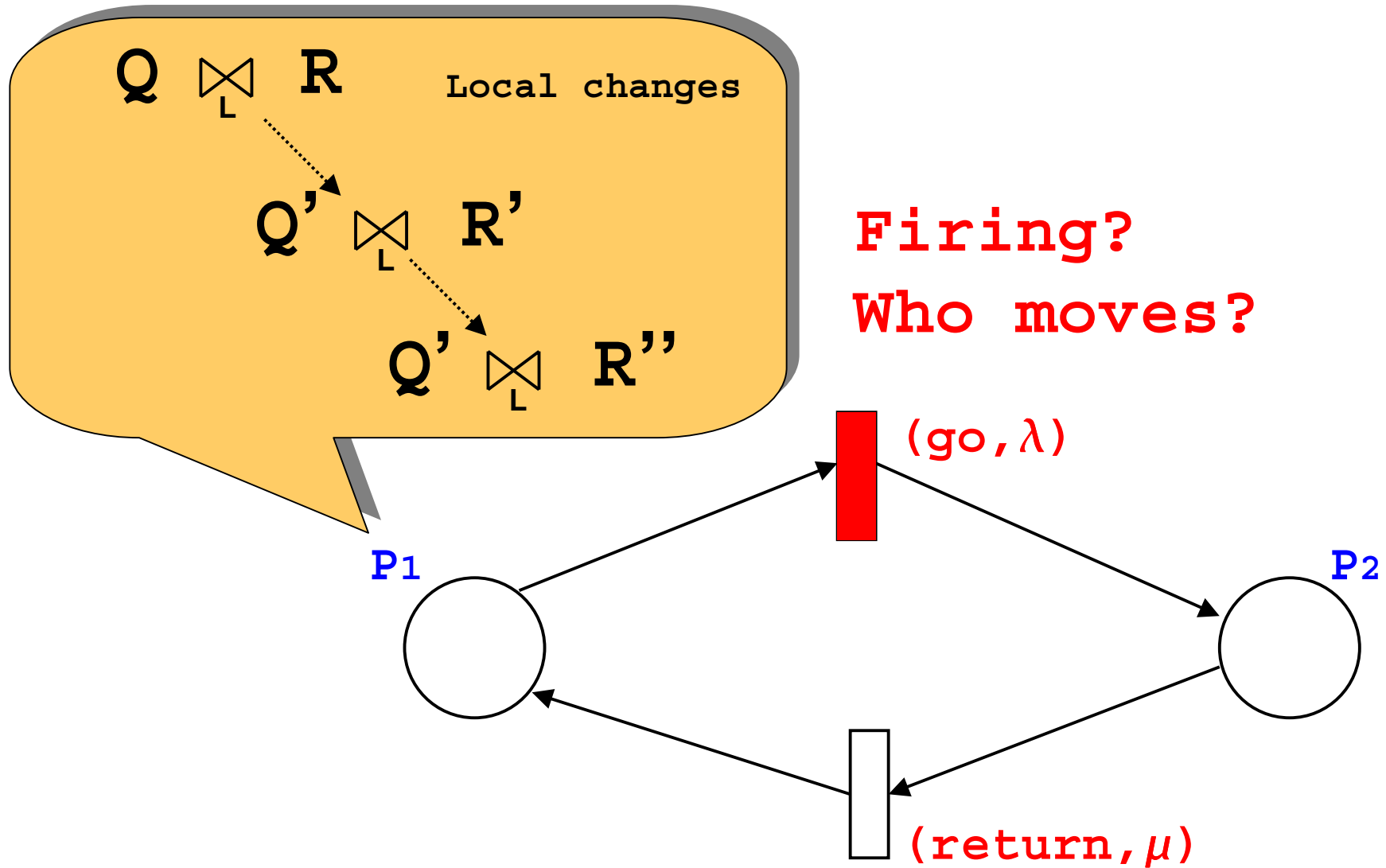
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- PEPA components perform activities to represent state changes
- ... in PEPA nets we distinguish between two types of changes ...
  - ✓ “local” changes  
(transitions of PEPA components)
  - ✓ “global” changes  
(net firings)



# PEPA nets: informal introduction

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# PEPA nets: informal introduction

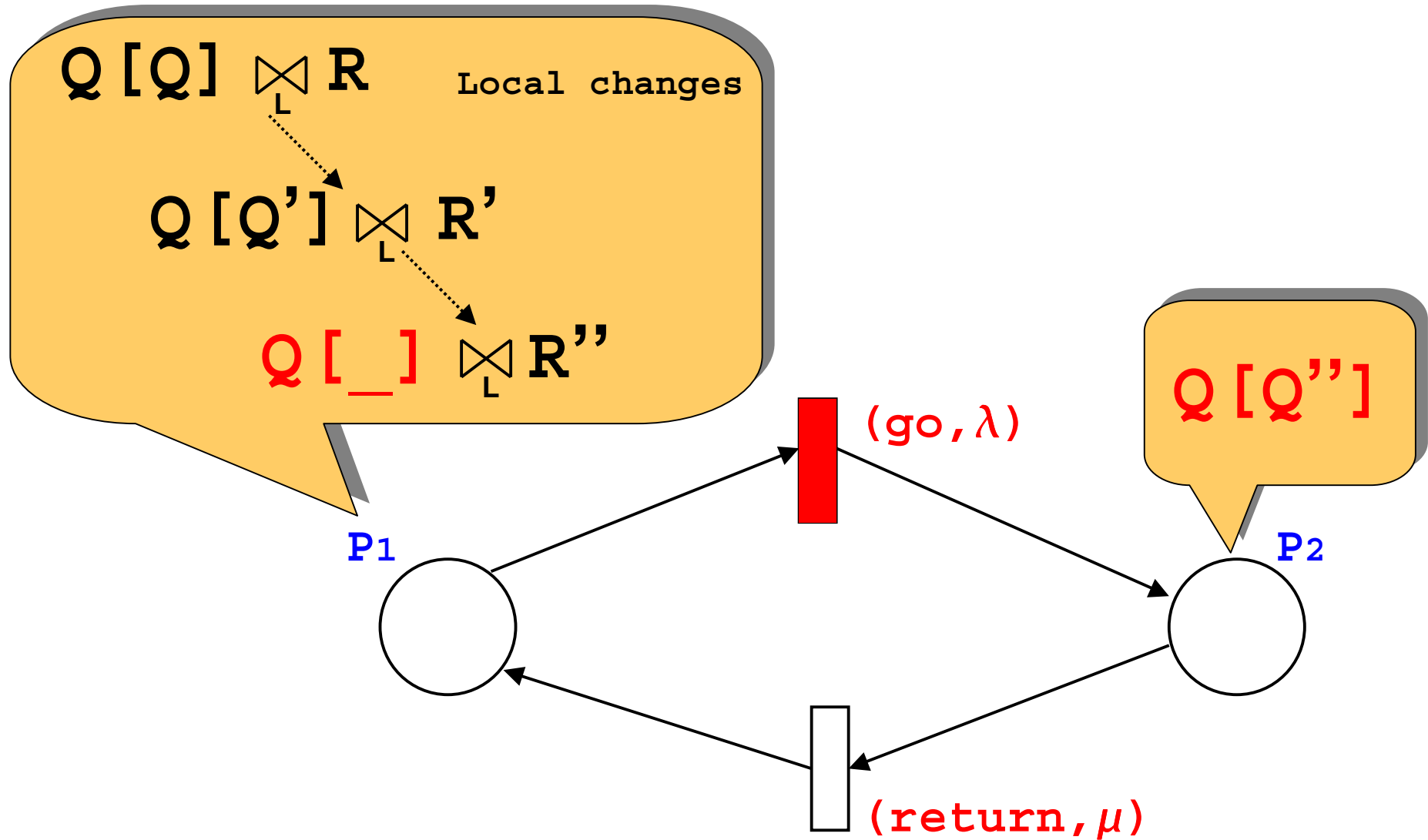
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- There is a PEPA **context** at each place of the net
- A PEPA context consists of
  - ✓ **static components**
  - ✓ **cells [ ]**



# PEPA nets: informal introduction

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# PEPA nets: informal introduction

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- Some assumptions


- ✓ Components can cooperate only when they are in the same place

- ✓ It is NOT possible for one component to cooperate with another component AND transfer to another place



# PEPA nets: syntax

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$$\begin{array}{l} S ::= (\alpha, r).S \quad (\text{prefix}) \\ | S + S \quad (\text{choice}) \\ | I \quad (\text{identifier}) \end{array}$$
$$\begin{array}{l} P ::= P \boxtimes_L P \quad (\text{cooperation}) \\ | P/L \quad (\text{hiding}) \\ | P[C] \\ | I \end{array}$$
$$\begin{array}{l} C ::= \_ \quad (\text{empty}) \\ | S \quad (\text{full}) \end{array}$$


# PEPA nets: markings and places

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$$M ::= (M_{\mathbf{P}}, \dots) \quad (\text{marking})$$
$$M_{\mathbf{P}} ::= \mathbf{P}[C, \dots] \quad (\text{place marking})$$
$$\mathbf{P}[C, \dots] \stackrel{\text{def}}{=} P[C] \underset{L}{\boxtimes} P \quad (\text{place defn})$$

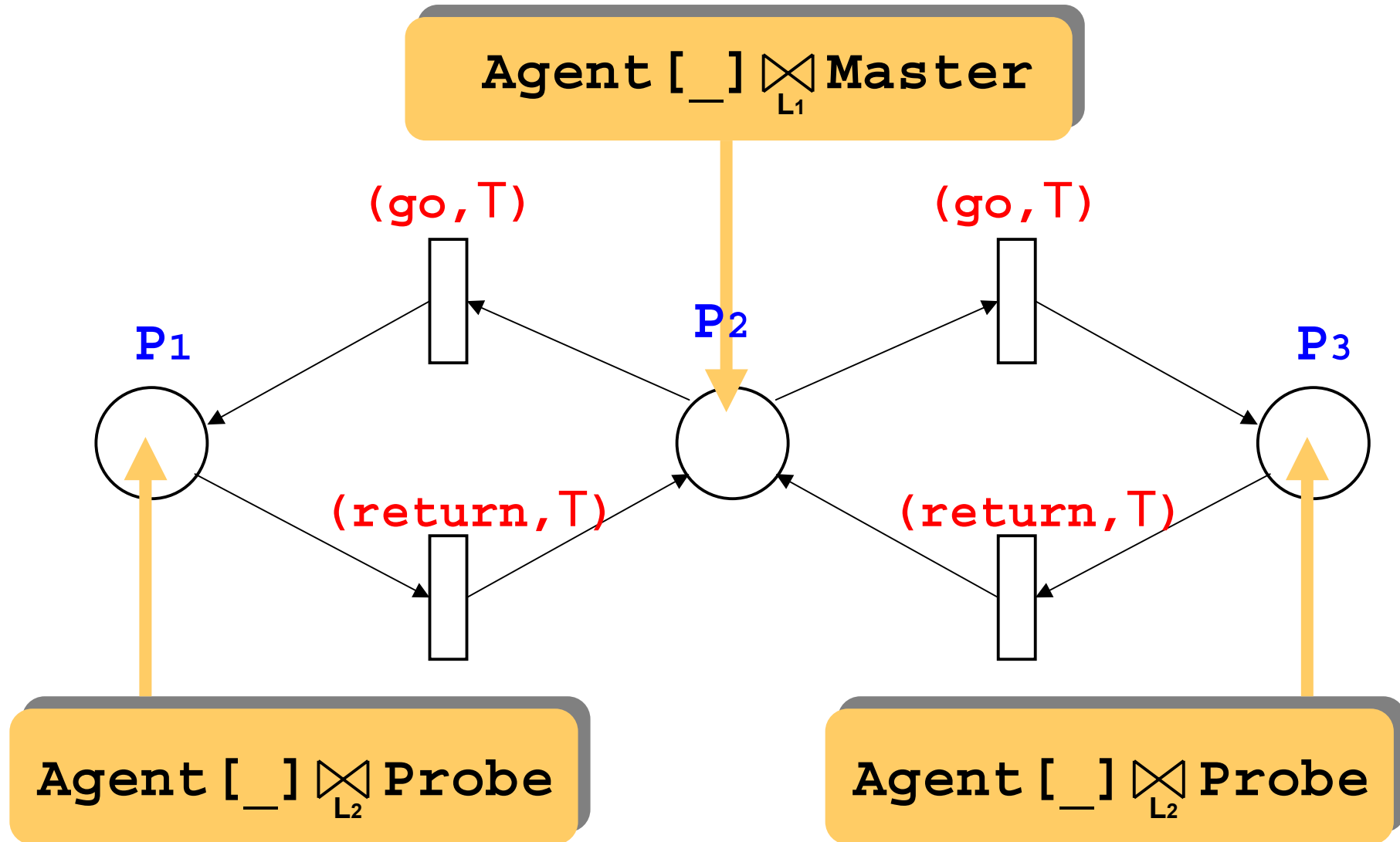

## Simple example: mobile agent

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- A **mobile software agent visits three sites**, where it interacts with static software components
- In two sites the agent **interrogates a network sensor for data** (on recent patterns of network traffic)
- In the other site, the agent **dumps the data to a master sensor**



# Simple example: Mobile agent



# Simple example: Mobile agent

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|                 |            |  |
|-----------------|------------|--|
| <b>Agent</b>    | <i>def</i> | <b>(go, λ) .Agent'</b>                       |
| <b>Agent'</b>   | <i>def</i> | <b>(interrogate, r<sub>i</sub>) .Agent''</b> |
| <b>Agent''</b>  | <i>def</i> | <b>(return, μ) .Agent'''</b>                 |
| <b>Agent'''</b> | <i>def</i> | <b>(dump, r<sub>d</sub>) .Agent</b>          |

|                |            |   |
|----------------|------------|---|
| <b>Master</b>  | <i>def</i> | <b>(dump, T) .Master'</b>               |
| <b>Master'</b> | <i>def</i> | <b>(analyse, r<sub>a</sub>) .Master</b> |

|              |            |  |
|--------------|------------|--|
| <b>Probe</b> | <i>def</i> | <b>(monitor, r<sub>m</sub>) .Probe +<br/>(interrogate, T) .Probe</b> |
|--------------|------------|--|



# PRISM

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- **Probabilistic model checker**  
probabilistic temporal logic, PCTL and CSL
  
- **Supports three models**  
DTMC, MDP, CTMC
  
- **Compact state representation (BDD)**
  
- **Input to the PRISM tool**
  1. description of the system
  2. set of properties to be checked



## From PEPA to PRISM

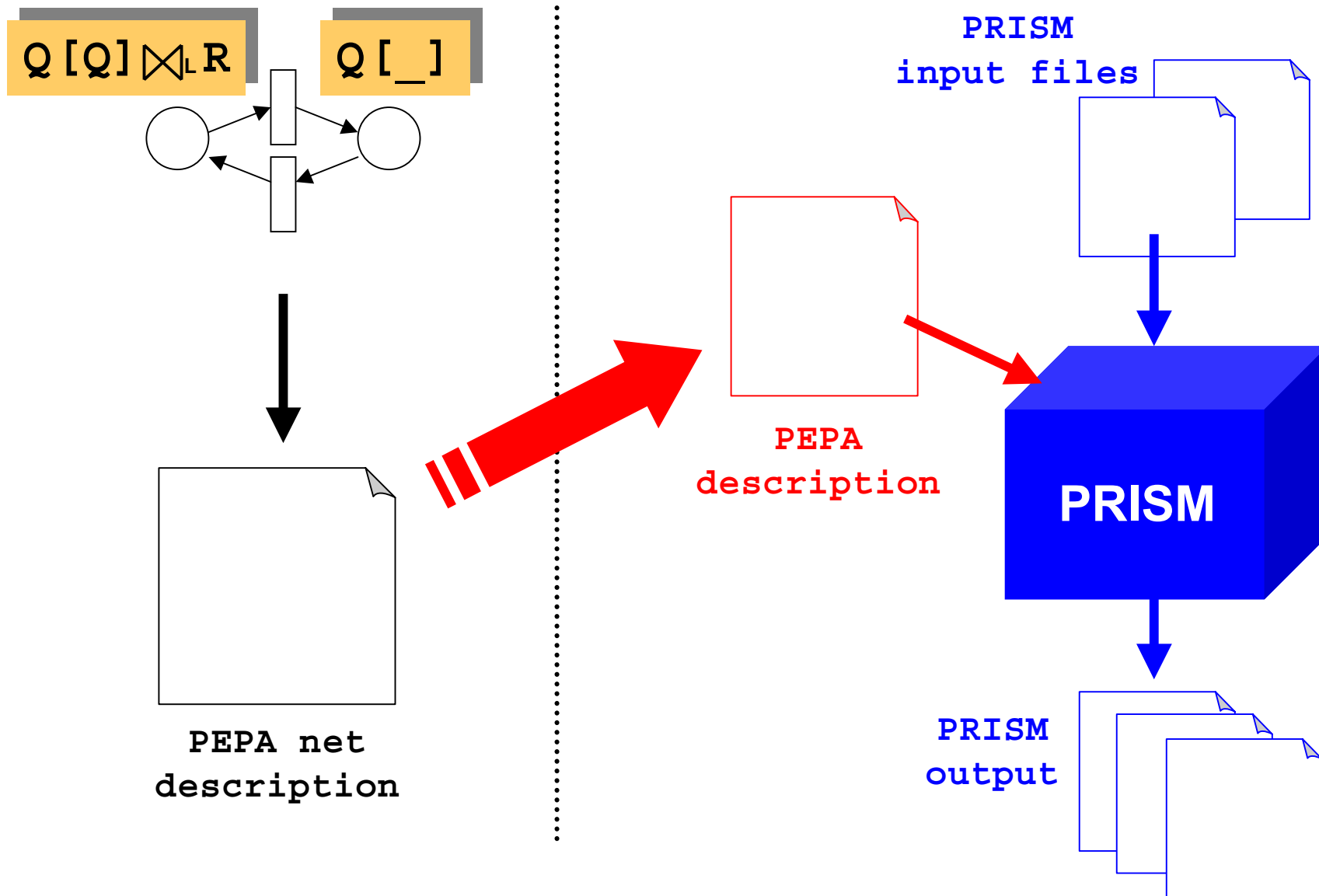
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- A compiler exists for translating PEPA models (a subset of PEPA) into PRISM models
- ... then the models can be analysed with the PRISM tool
  - ✓ The steady-state probability distribution for the underlying CTMC can be automatically derived
  - ✓ Properties can be verified





# From PEPA nets to PRISM



# From PEPA nets to PEPA

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- We need to map the net structure into (standard) PEPA components
- Problems
  - ✓ What happens to **different transitions** with the **same** associated **label**?
  - ✓ What happens to **replica** of the same **static component**, resident in different places
  - ✓ What about **cells**?

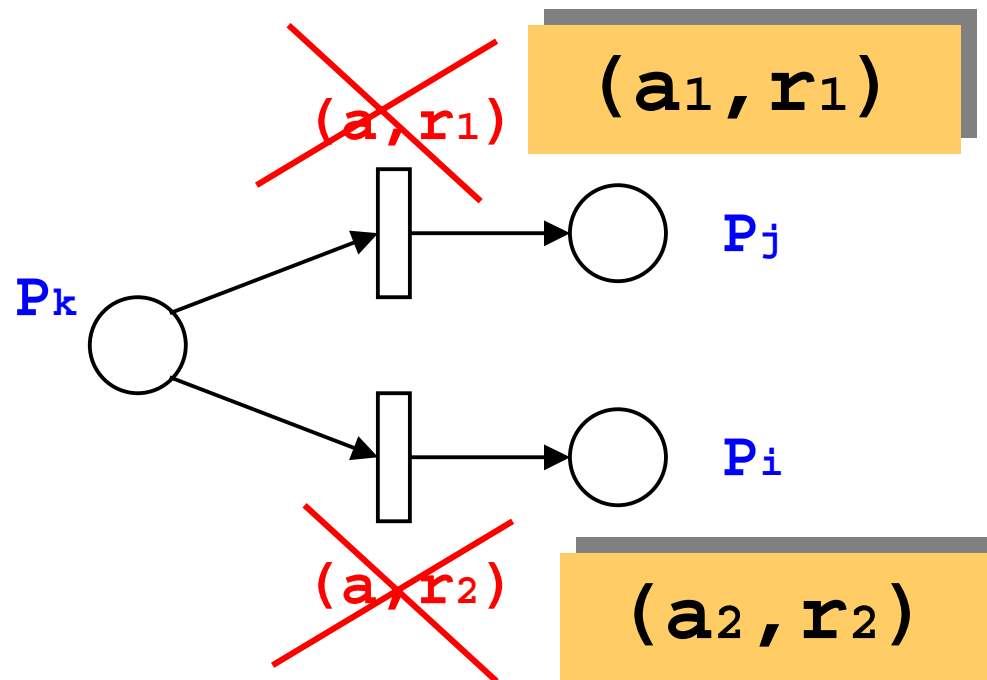


# From PEPA nets to PEPA

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- Steps of the translation algorithm

## 0. Preprocessing



# From PEPA nets to PEPA

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## 1. Translation of static components

In order to **avoid wrong synchronisations** we need to distinguish replicas of the same static component.

This is done by **renaming** action types and derivatives

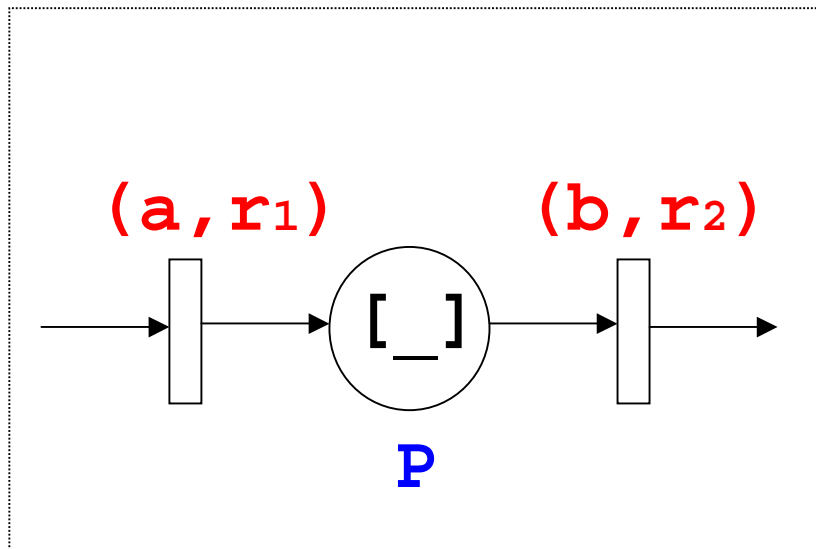


# From PEPA nets to PEPA

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## 2. Translation of cells

A new PEPA component need to be defined for each cell  $i$  within each place  $P$


$$\text{Cell}_{i0} \stackrel{\text{def}}{=} (a, r1) . \text{Cell}_{i1}$$
$$\text{Cell}_{i1} \stackrel{\text{def}}{=} (b, r2) . \text{Cell}_{i0}$$

# From PEPA nets to PEPA

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## 3. Translation of tokens

The movement of a token in a new place and its interaction with static components are considered

To allow correct synchronisations the new names introduced in the previous steps are introduced in the token as well



# From PEPA nets to PEPA

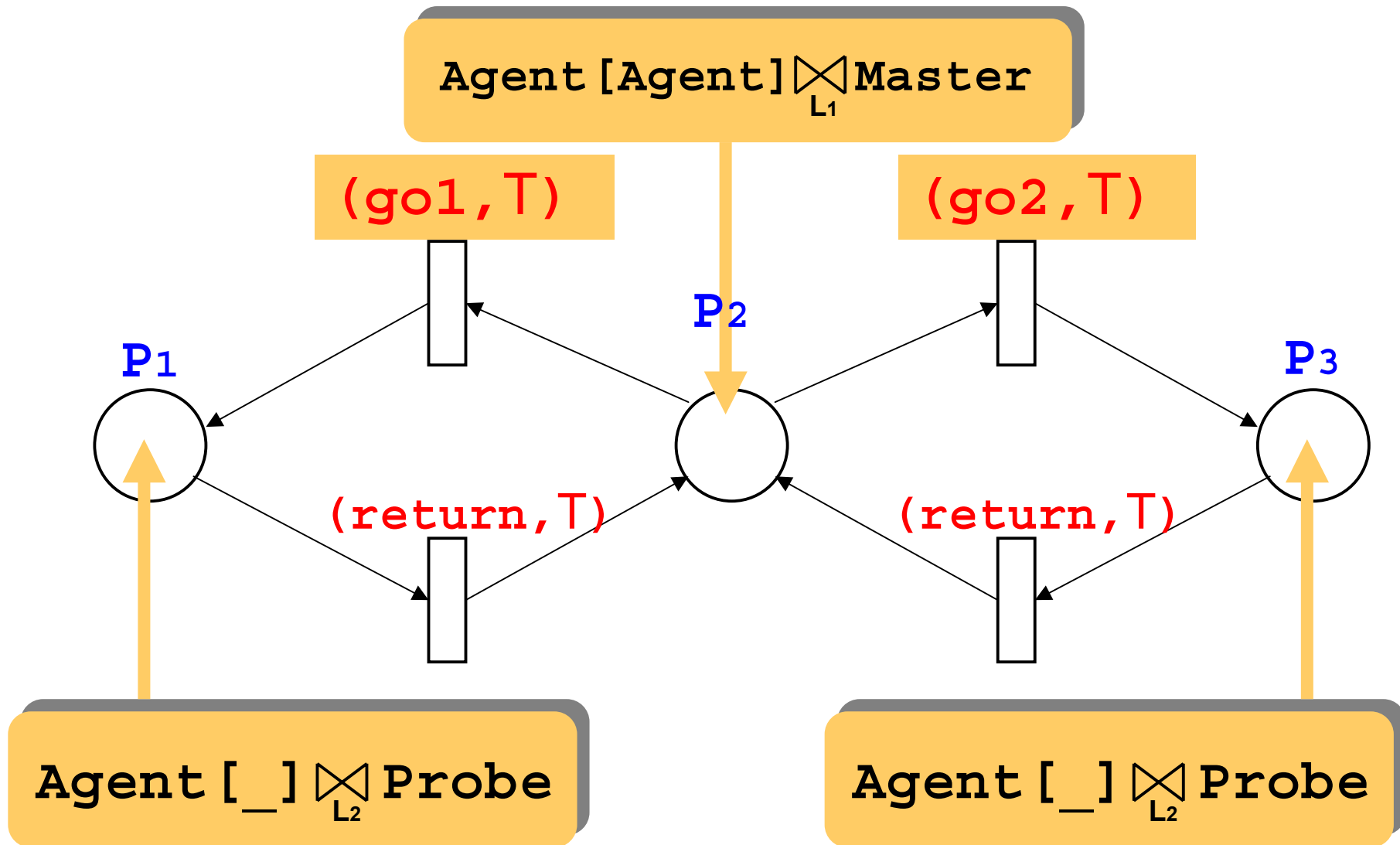
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## 4. Building the system equation

All PEPA components built in the previous steps are put in parallel and forced to synchronise on common action types



# A simple example: preprocessing





# A simple example: static components

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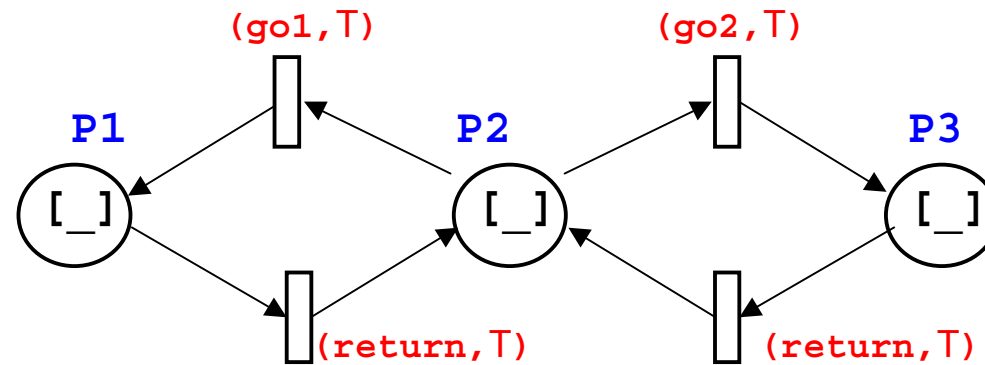
**Master**  $\stackrel{def}{=} (\text{dump}, T) . \text{Master}' +$   
**Master'**  $\stackrel{def}{=} (\text{analyse}, r_a) . \text{Master}$

**Probe<sub>1</sub>**  $\stackrel{def}{=} (\text{monitor}_1, r_m) . \text{Probe}_1 +$   
 $(\text{interrogate}_1, T) . \text{Probe}_1$

**Probe<sub>2</sub>**  $\stackrel{def}{=} (\text{monitor}_2, r_m) . \text{Probe}_2 +$   
 $(\text{interrogate}_2, T) . \text{Probe}_2$



# A simple example: cells



|               |            |  |
|---------------|------------|--|
| <b>Cell10</b> | <i>def</i> | <b>= (go1, T) . Cell11</b>                         |
| <b>Cell11</b> | <i>def</i> | <b>= (return, T) . Cell10</b>                      |
| <b>Cell20</b> | <i>def</i> | <b>= (return, T) . Cell21</b>                      |
| <b>Cell21</b> | <i>def</i> | <b>= (go1, T) . Cell20 +<br/>(go2, T) . Cell20</b> |
| <b>Cell30</b> | <i>def</i> | <b>= (go2, T) . Cell31</b>                         |
| <b>Cell31</b> | <i>def</i> | <b>= (return, T) . Cell30</b>                      |



# A simple example: tokens

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|                 |                            |  |
|-----------------|----------------------------|--|
| <b>Agent</b>    | <sup>def</sup><br><b>=</b> | <b>(go1, λ) .Agent1' +<br/>(go2, λ) .Agent2'</b> |
| <b>Agent1'</b>  | <sup>def</sup><br><b>=</b> | <b>(interrogate1, r<sub>i</sub>) .Agent''</b>    |
| <b>Agent2'</b>  | <sup>def</sup><br><b>=</b> | <b>(interrogate2, r<sub>i</sub>) .Agent''</b>    |
| <b>Agent''</b>  | <sup>def</sup><br><b>=</b> | <b>(return, μ) .Agent'''</b>                     |
| <b>Agent'''</b> | <sup>def</sup><br><b>=</b> | <b>(dump, r<sub>d</sub>) .Agent</b>              |



# A simple example: model equation

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**System** <sup>def</sup> =

**(Cell<sub>10</sub>  $\bowtie_{K_1}$  (Probe<sub>1</sub>  $\bowtie_{K_2}$  (Agent  $\bowtie_{K_3}$  (Cell<sub>21</sub>  $\bowtie_{K_4}$   
  
(Master  $\bowtie_{K_5}$  (Probe<sub>2</sub>  $\bowtie_{K_6}$  Cell<sub>30</sub>))))))**

**K<sub>1</sub> = {go1, return}**

**K<sub>3</sub> = {go1, go2, return}**

**K<sub>6</sub> = {go2, return}**

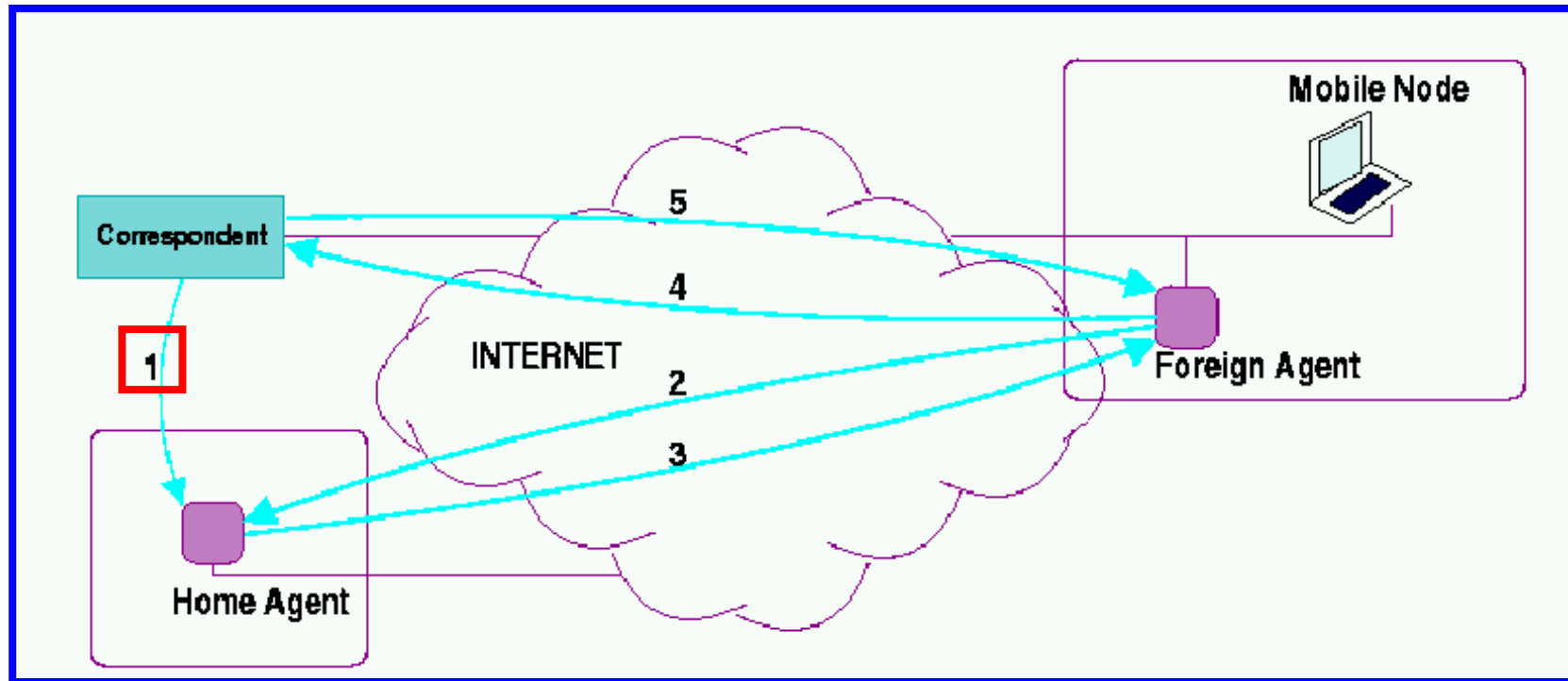
**K<sub>2</sub> = {interrogate1}**

**K<sub>4</sub> = {dump}**

**K<sub>5</sub> = {interrogate2}**

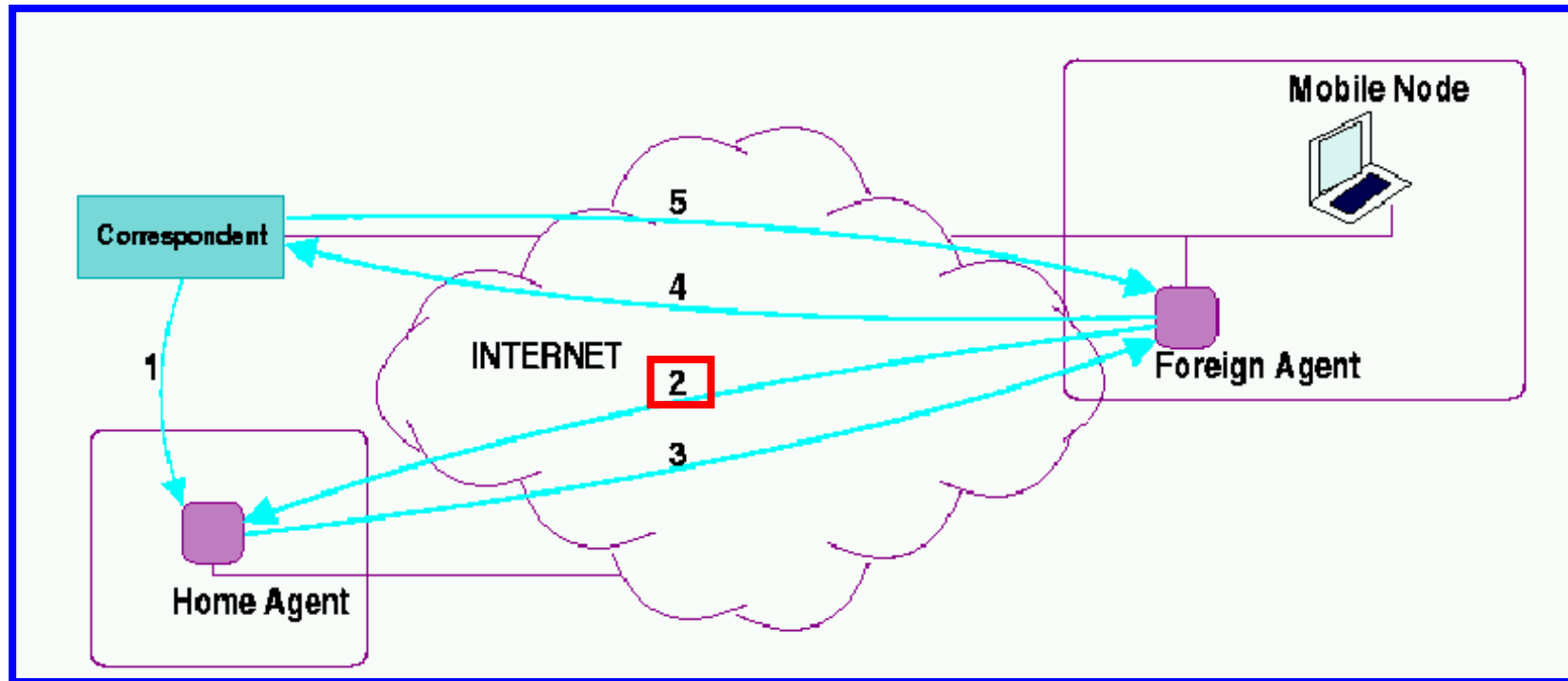


# Complex example: mobile IP



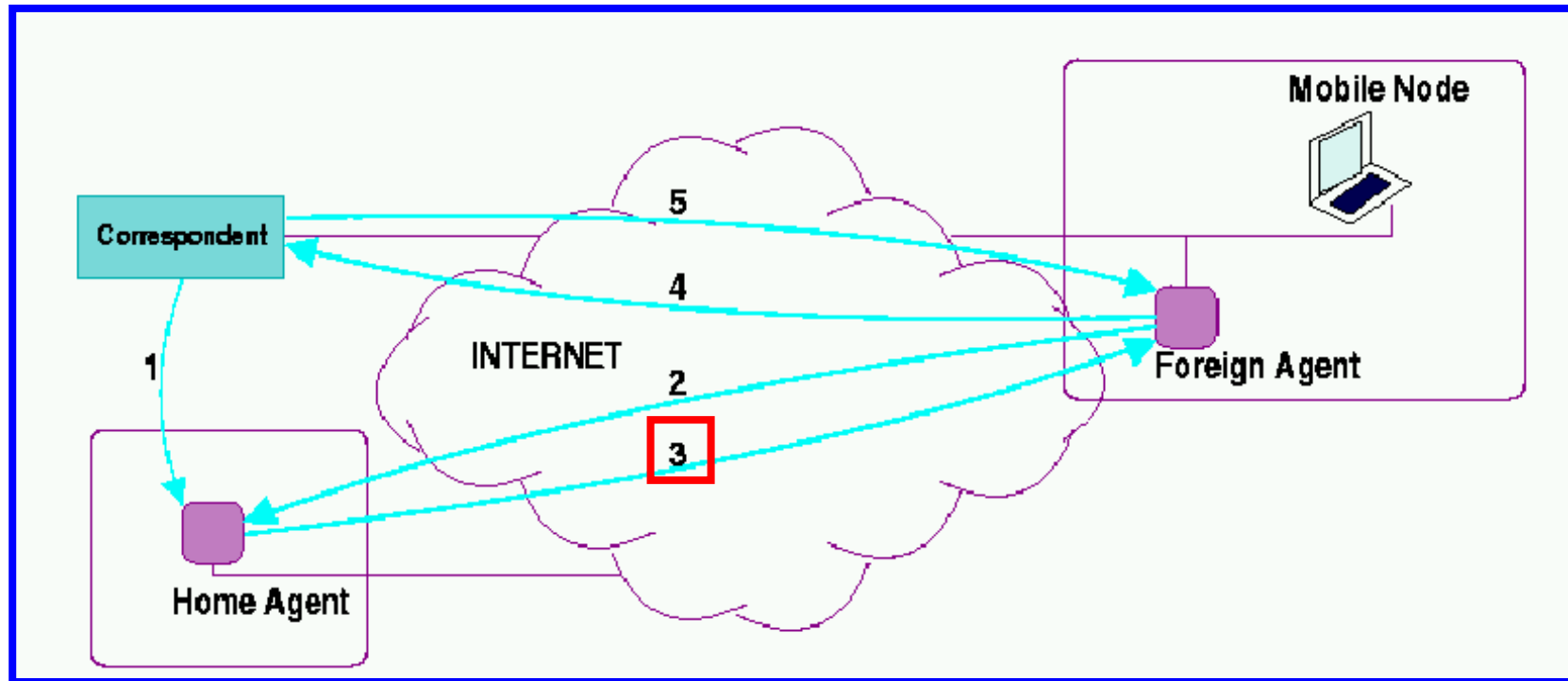
1. Correspondent sends IP packets to the mobile node at his home address

# Complex example: mobile IP



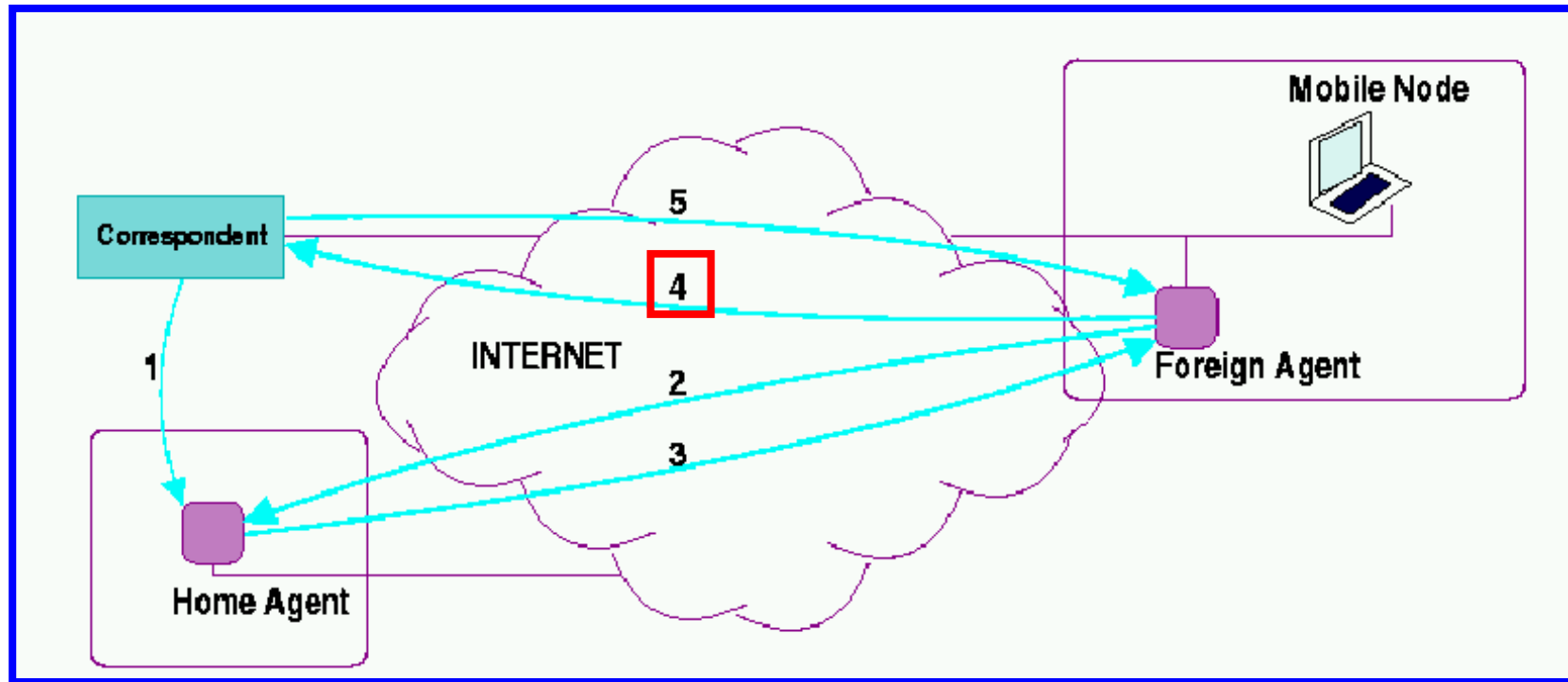
2. The Mobile Node sends its new IP address to the Home Agent

# Complex example: mobile IP



3. The Home Agent forwards packets to the Mobile Node

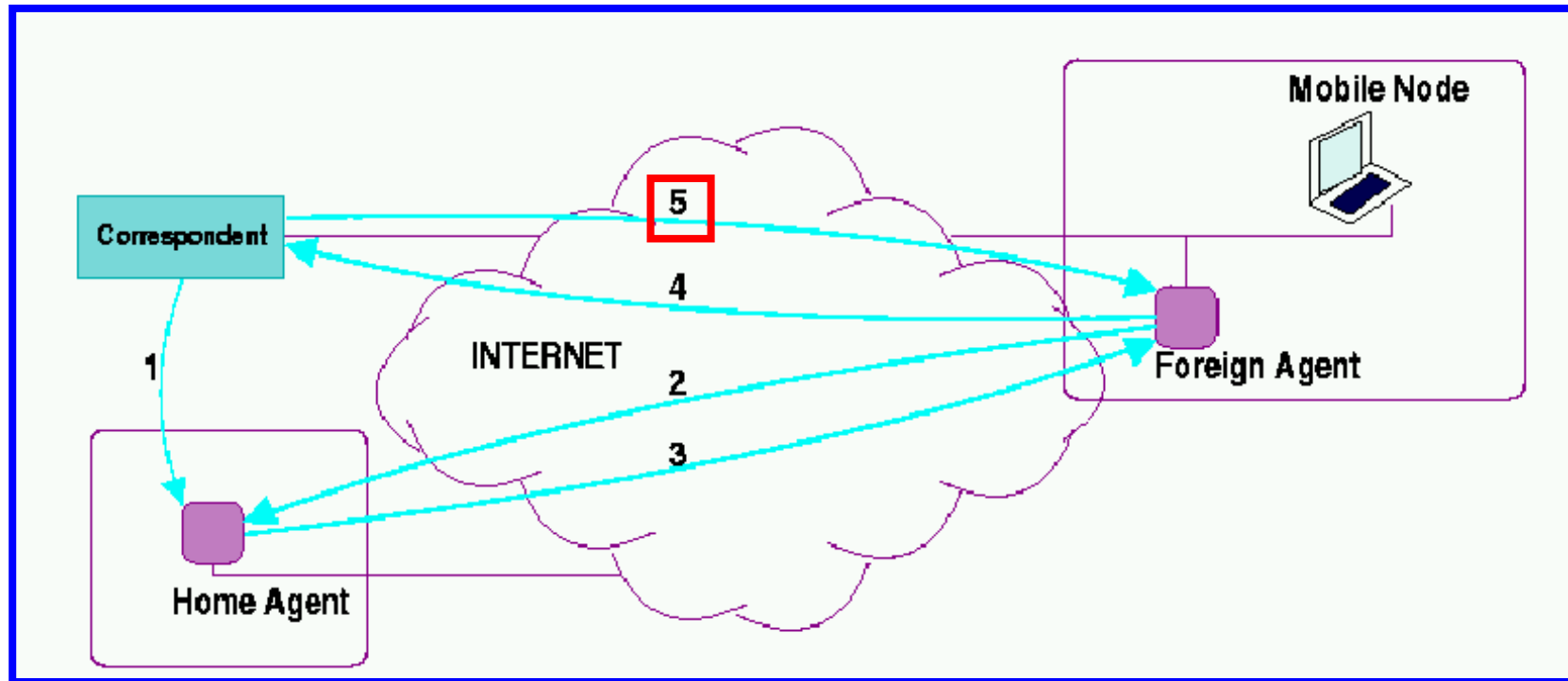
# Complex example: mobile IP



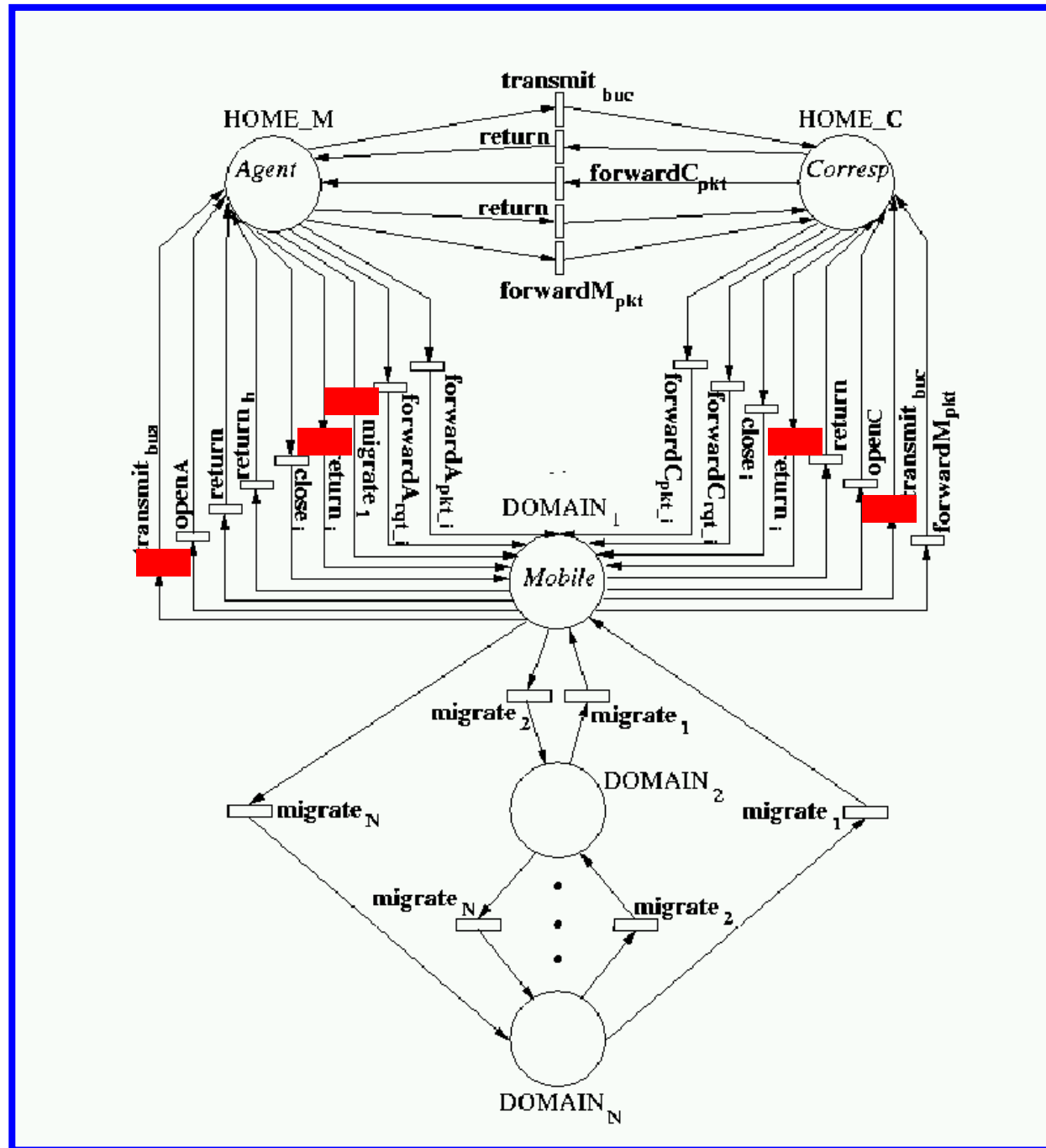
4. The Mobile Node sends its new IP address to the Correspondent



# Complex example: mobile IP



5. The Correspondent sends packets directly to the Mobile Node



## Complex example: mobile IP

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- 1 Mobile Node, 1 Correspondent,  
1 Domain
  - ✓ 2.8 million of states
  - ✓ 16 million of transitions
  - ✓ in 13.2 seconds, 1.6GHz Pentium IV  
with 256 MB of RAM



# Conclusions

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- PEPA nets is relatively new but we think that it can provide a framework for modelling systems characterised by some mobility

## Future work

- ✓ Synchronisation over net transitions
- ✓ Movement of more than one token
  
- ✓ Graphical interface (done!)

