

PEPA models of Internet worm attacks

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8th September 2005

Joint work with Jeremy Bradley and Stephen Gilmore

Outline

Introduction

Internet worm models

Continuous Approximation

Quantified analysis

Conclusions

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PEPA

$$\begin{aligned} S &::= (\alpha, r).S \mid S + S \mid A \\ P &::= S \mid P \underset{L}{\boxtimes} P \mid P/L \end{aligned}$$

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| HIDING: | P/L | abstraction $\alpha \in L \Rightarrow \alpha \rightarrow \tau$ |

Generating a CTMC

The corresponding Continuous Time Markov Chain (CTMC) is derived automatically from the structured operational semantics which define the language:

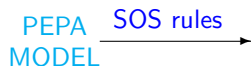
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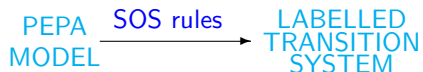
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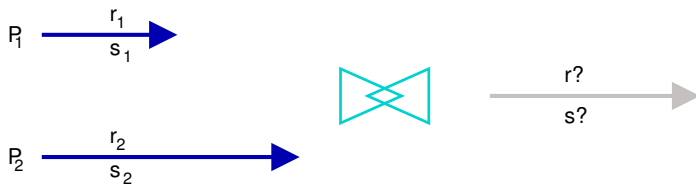
Solving the model has meant finding the steady state probability distribution over the entire state space.

Timed Synchronisation

- ▶ The issue of what it means for two timed activities to synchronise is a vexed one....

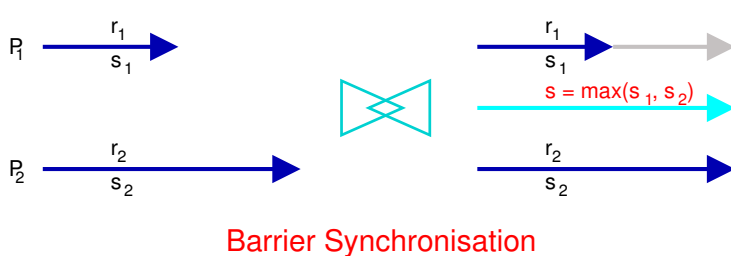
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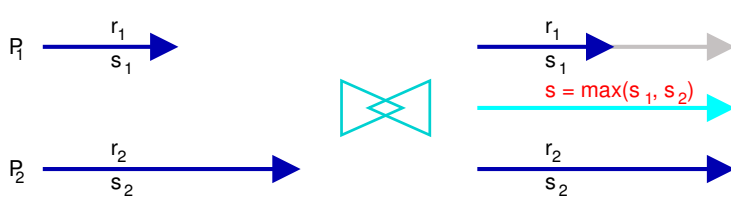
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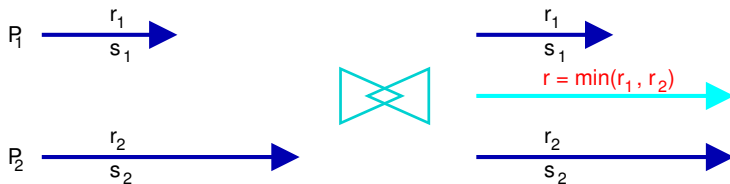
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- ▶ Synchronisation, or cooperation cannot make a component exceed its bounded capacity.
- ▶ Thus the apparent rate of a cooperation is the **minimum** of the apparent rates of the co-operands.

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In all the models we assume that the infection must pass over a network, which can sustain M independent concurrent connections.

Model 1

The Susceptible-Infective-Removed model.

$$S = (\text{infect}S, \top).I$$

$$I = (\text{infect}I, \beta).I + (\text{patch}, \gamma).R$$

$$R = \text{stop}$$

$$\text{Net} = (\text{infect}I, \top).\text{Net}'$$

$$\text{Net}' = (\text{infect}S, \beta).\text{Net}$$

$$\text{Sys} = (S[M] \parallel I) \bowtie_L \text{Net}[M]$$

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$$S = (\text{infect}S, \top).I$$

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The major limitations of the CTMC approach are the **state space explosion** problem and the reliance on **exponential distributions**.

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Only appropriate for some models, but results are promising in those cases. **large numbers of repeated components**

Differential equations from PEPA models

- ▶ In a PEPA model the state at any current time is the local derivative or **state of each component** of the model.
- ▶ When we have large numbers of repeated components it can make sense to represent the state of the system as the count of the current number of each possible local derivative or component type.
- ▶ We can approximate the behaviour of the model by treating the number of each component type as a continuous variable, and the state of the model as a whole as the set of such variables.
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Derivation of the system of ODES representing the PEPA model then proceeds via an activity matrix which keeps track of the impact of each activity type on each component type.

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In terms of Jeremy's classification yesterday, this means we use the passive synchronisation scheme in the ODEs.

Mapping to an ODE

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$$\frac{dv_{12}(t)}{dt} = -\gamma v_{12}(t) + \beta I_{11}(t)v_{22}(t)$$

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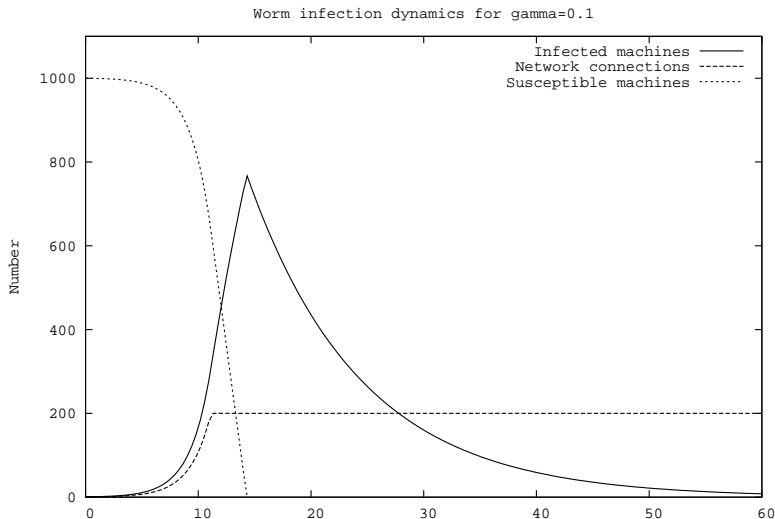
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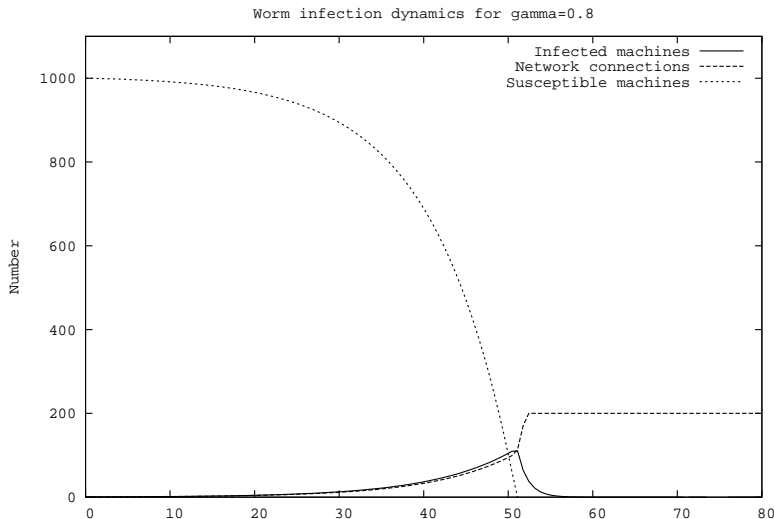
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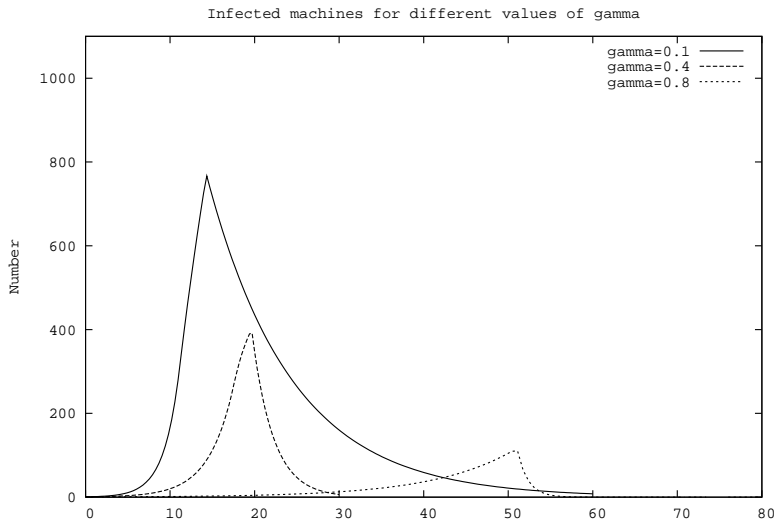
In the first experiment we varied the rate at which the patch is applied, γ , representing different (human) response rates to the infection.

Model 1: $\gamma = 0.1$ 

Model 1: $\gamma = 0.8$



Model 1: Number of infected machines as γ increases



Model 2: Susceptible-Infective-Removed-Reinfection model

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$$\begin{aligned}\frac{dv_{11}(t)}{dt} &= -\beta I_{11}(t)v_{22}(t) + \mu v_{13}(t) \\ \frac{dv_{12}(t)}{dt} &= -\gamma v_{12}(t) + \beta I_{11}(t)v_{22}(t) \\ \frac{dv_{13}(t)}{dt} &= -\mu v_{13}(t) + \gamma v_{12}(t) \\ \frac{dv_{21}(t)}{dt} &= -\beta I_{21}(t)v_{12}(t) + \beta I_{11}(t)v_{22}(t) \\ \frac{dv_{22}(t)}{dt} &= -\beta I_{11}(t)v_{22}(t) + \beta I_{21}(t)v_{12}(t)\end{aligned}$$

where $v_{11} \leftrightarrow S$, $v_{12} \leftrightarrow I$, $v_{13} \leftrightarrow R$, $v_{21} \leftrightarrow Net$, $v_{22} \leftrightarrow net'$.

Model 2: experiments

We assume a susceptible population of $N = 1000$ computers.

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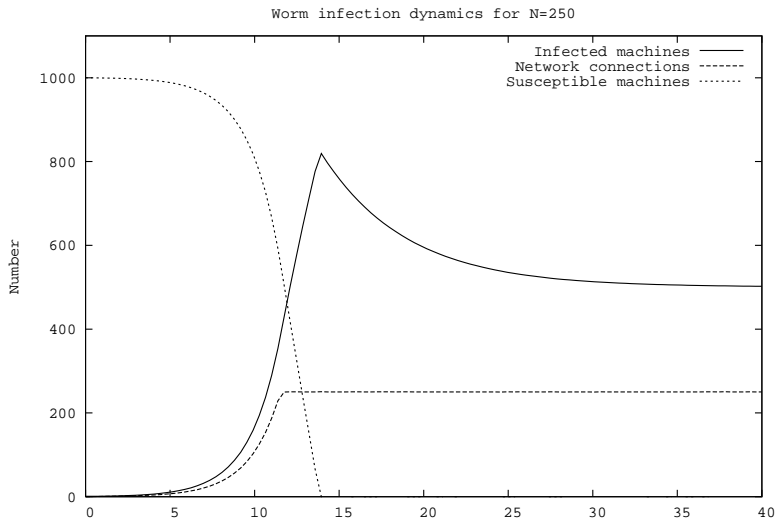
We assume that the system starts with one infected computer.

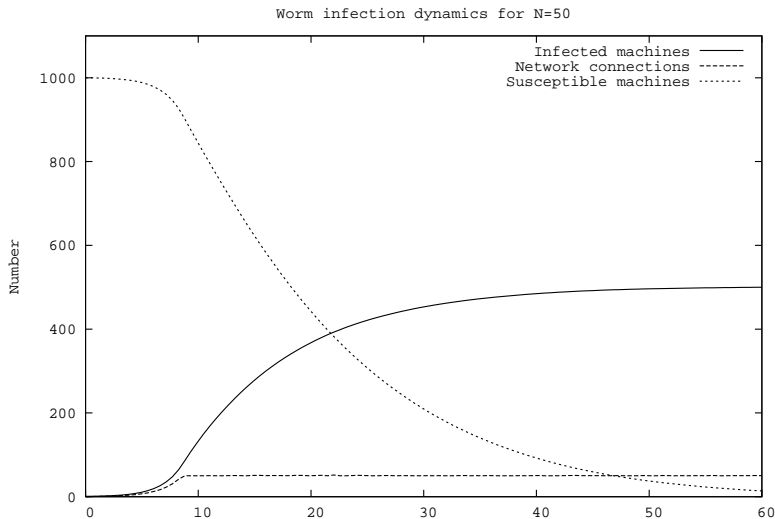
Model 2: experiments

We assume a susceptible population of $N = 1000$ computers.

We assume that the system starts with one infected computer.

In this experiment we varied the network capacity, i.e. M . This restricts the medium over which the infection is transmitted.

Model 2: $N = 250$ 

Model 2: $N = 50$ 

Model 3: Susceptible-Infective-Removed-Attack model

$$S = (\text{infect}S, \top).I$$

$$I = (\text{infect}I, \beta).I + (\text{attack}, \lambda).I + (\text{patch}, \gamma).R$$

$$R = \text{stop}$$

$$\text{Net} = (\text{infect}I, \top).\text{Net}'$$

$$\text{Net}' = (\text{infect}S, \beta).\text{Net}$$

$$A = (\text{attack}, \top).A'$$

$$A' = (\text{recover}, \mu).A$$

$$\text{Sys} = ((S[N] \parallel I) \bowtie_L \text{Net}[M]) \bowtie_{L'} A[T]$$

where $L = \{\text{infect}I, \text{infect}S\}$, $L' = \{\text{attack}\}$.

Mapping to an ODE

$$\frac{dv_{11}(t)}{dt} = -\beta I_{11}(t)v_{22}(t)$$

$$\frac{dv_{12}(t)}{dt} = -\gamma v_{12}(t) + \beta I_{11}(t)v_{22}(t)$$

$$\frac{dv_{13}(t)}{dt} = \gamma v_{12}(t)$$

$$\frac{dv_{21}(t)}{dt} = -\beta I_{21}(t)v_{12}(t) + \beta I_{11}(t)v_{22}(t)$$

$$\frac{dv_{22}(t)}{dt} = -\beta I_{11}(t)v_{22}(t) + \beta I_{21}(t)v_{12}(t)$$

$$\frac{dv_{31}(t)}{dt} = -\lambda I_{31}(t)v_{12}(t) + v_{32}(t)\mu$$

$$\frac{dv_{32}(t)}{dt} = -v_{32}(t)\mu + \lambda I_{31}(t)v_{12}(t)$$

$$v_{11} \leftrightarrow S, v_{12} \leftrightarrow I, v_{13} \leftrightarrow R, v_{21} \leftrightarrow Net, v_{22} \leftrightarrow net', v_{31} \leftrightarrow A, v_{32} \leftrightarrow A'.$$

Model 3: experiments

We assume a susceptible population of $N = 1000$ computers, a network capacity of $M = 200$

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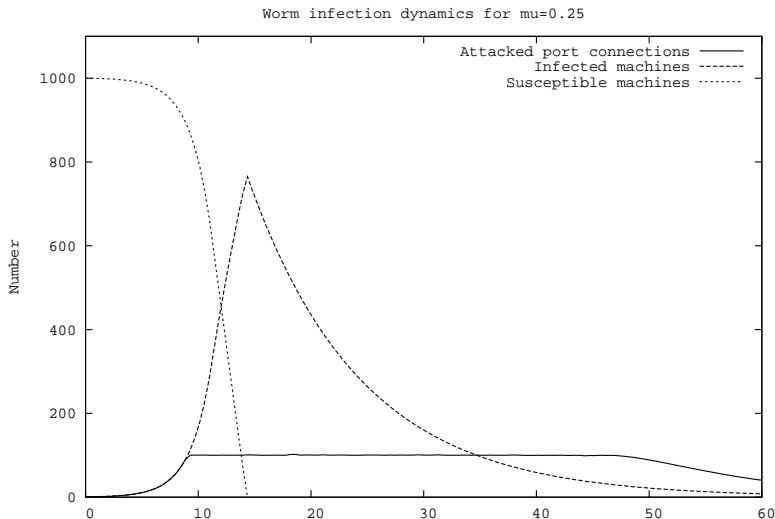
We assume that the system starts with one infected computer, and that the target of the attack has 100 ports on which it can accept connections.

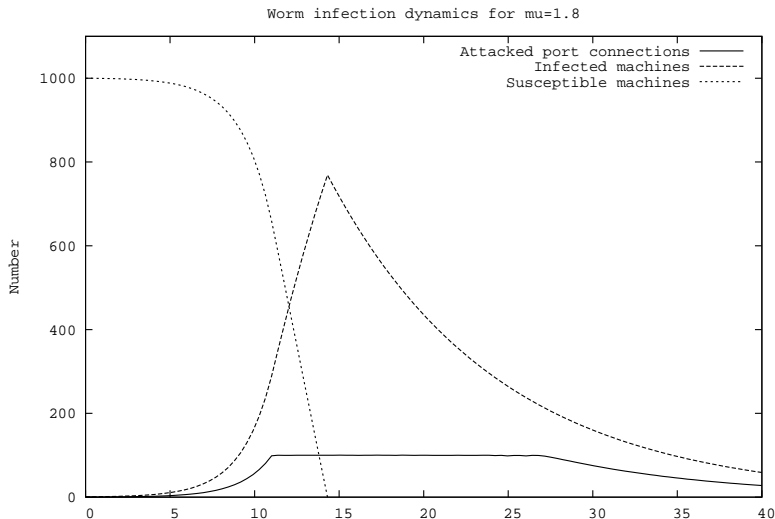
Model 3: experiments

We assume a susceptible population of $N = 1000$ computers, a network capacity of $M = 200$

We assume that the system starts with one infected computer, and that the target of the attack has 100 ports on which it can accept connections.

In this experiment we varied the rate μ at which a port timeouts and becomes usable again in the attacked machine.

Model 3: $\mu = 0.25$ 

Model 3: $\mu = 1.8$ 

Outline

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Continuous Approximation

Quantified analysis

Conclusions

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Conclusions

ODEs are great!

- ▶ We could evaluate small systems using the CTMC semantics but not with realistic populations
- ▶ We could construct the ODEs directly (eg. [Nicol et al]) but using the process algebra gives a more accessible model, and one which is amenable to other analyses such as model checking.
- ▶ For these models there are still many experiments to be considered and variations to the models which could be made.