

A semi-quantitative equivalence for abstracting from fast reactions

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Outline

Competitive inhibition

Bio-PEPA

Fast-slow bisimilarity

Slow bisimilarity

Conclusions

Example: competitive inhibition, bimolecular



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- ▶ reactions produce two **intermediate species** and product

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$$\frac{dS}{dt} = \dots$$

$$\frac{dE}{dt} = \dots$$

$$\frac{dI}{dt} = \dots$$

$$\frac{dEI}{dt} = \dots$$

$$\frac{dSE}{dt} = \dots$$

$$\frac{dP}{dt} = \dots$$



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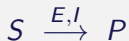
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 - ▶ derive new model with new rates by setting some ODEs to zero
- ▶ example: Michaelis-Menten kinetics for substrate and enzyme
- ▶ need to understand limitations of abstracted version

Example: competitive inhibition, abstract



- ▶ using QSSA, express competitive inhibition as a single reaction
- ▶ results in fewer ODEs
- ▶ reaction rate depends on S , E and I

Example: competitive inhibition, abstract



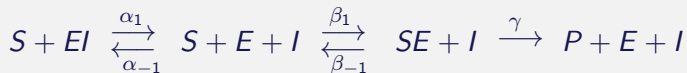
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How can these ideas be used to develop a behavioural equivalence for a biological process algebra?

Bio-PEPA: competitive inhibition, bimolecular

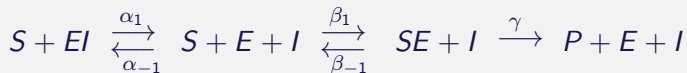


Bio-PEPA: competitive inhibition, bimolecular



- name each reaction

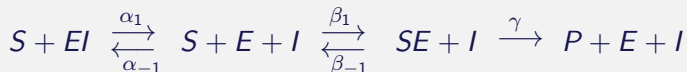
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- ▶ well-defined Bio-PEPA species/sequential components

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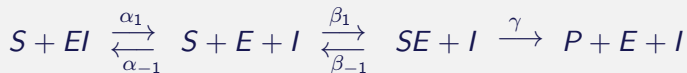
Prefix notation: $(\alpha, \kappa) \text{ op}$

α reaction name

κ stoichiometry for reaction

op role of species in reaction

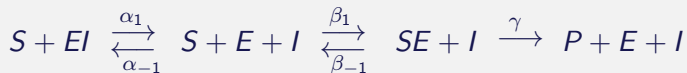
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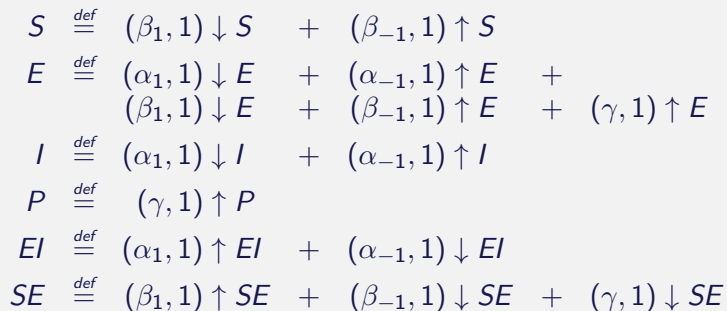
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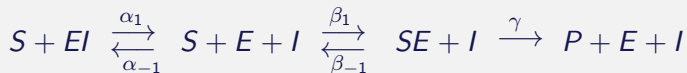
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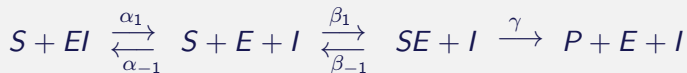


- ▶ well-defined Bio-PEPA model component, includes quantities/levels for each species

$$M \stackrel{\text{def}}{=} S(\ell_S) \boxtimes_* E(\ell_E) \boxtimes_* I(\ell_I) \boxtimes_* P(\ell_P) \boxtimes_* EI(\ell_{EI}) \boxtimes_* SE(\ell_{SE})$$



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- ▶ well-defined Bio-PEPA system, defines context for model

$$\mathcal{P} = \langle \mathcal{V}, \mathcal{N}, \mathcal{K}, \mathcal{F}, \text{Comp}, M \rangle$$

- \mathcal{V} volume and location information for model
- \mathcal{N} quantitative information for each species
- \mathcal{K} constant definitions
- \mathcal{F} rate equations for each reaction
- Comp* Bio-PEPA definition for each species

Bio-PEPA: competitive inhibition, abstract



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$$S' \stackrel{\text{def}}{=} (\gamma, 1) \downarrow S' \quad P' \stackrel{\text{def}}{=} (\gamma, 1) \uparrow P'$$

Bio-PEPA: competitive inhibition, abstract

$$S \xrightarrow[\gamma]{E, I} P$$

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$$((\alpha, \kappa) \downarrow S)(\ell) \xrightarrow{(\alpha, [S:\downarrow(\ell, \kappa)])}_c S(\ell - \kappa) \quad \kappa \leq \ell \leq N_S$$

$$((\alpha, \kappa) \uparrow S)(\ell) \xrightarrow{(\alpha, [S:\uparrow(\ell, \kappa)])}_c S(\ell + \kappa) \quad 0 \leq \ell \leq N_S - \kappa$$

$$((\alpha, \kappa) \oplus S)(\ell) \xrightarrow{(\alpha, [S:\oplus(\ell, \kappa)])}_c S(\ell) \quad \kappa \leq \ell \leq N_S$$

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Bio-PEPA semantics (continued)

- Cooperation for $\alpha \in L$

$$\frac{P \xrightarrow{(\alpha, v)}_c P' \quad Q \xrightarrow{(\alpha, u)}_c Q'}{P \boxtimes_L Q \xrightarrow{(\alpha, v::u)}_c P' \boxtimes_L Q'} \quad \alpha \in L$$

Bio-PEPA semantics (continued)

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- ▶ can define stochastic relation which is quantitative
- ▶ we work semi-quantitatively with capability relation
- ▶ reaction names are partitioned into two sets

\mathcal{A}_f : fast reactions
 \mathcal{A}_s : slow reactions

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- similar definition to Milner's weak bisimilarity
- fast reactions play same role as τ labelled transitions



Congruence of fast-slow bisimilarity

- congruence for cooperation if no shared fast reactions

$$P_1 \approx_{\mathcal{A}_f} P_2 \Rightarrow P_1 \boxtimes_L Q \approx_{\mathcal{A}_f} P_2 \boxtimes_L Q, \quad Q \boxtimes_L P_1 \approx_{\mathcal{A}_f} Q \boxtimes_L P_2$$

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- congruence for extension operator if no shared reactions

$$A_1 \approx_{\mathcal{A}_f} A_2 \Rightarrow A_1\{B\} \approx_{\mathcal{A}_f} A_2\{B\} \quad \text{and} \quad B\{A_1\} \approx_{\mathcal{A}_f} B\{A_2\}$$

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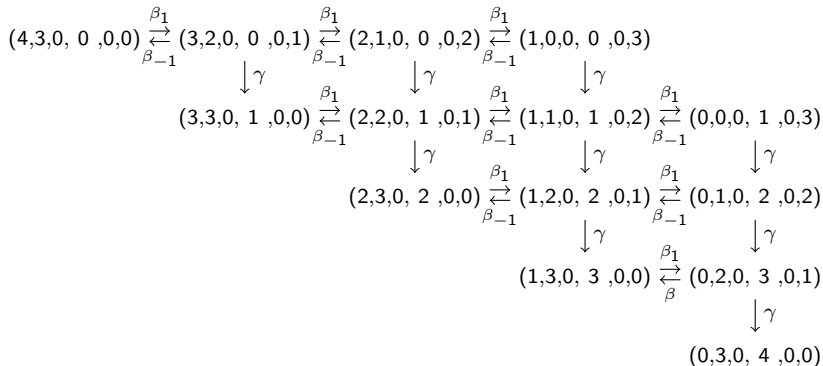
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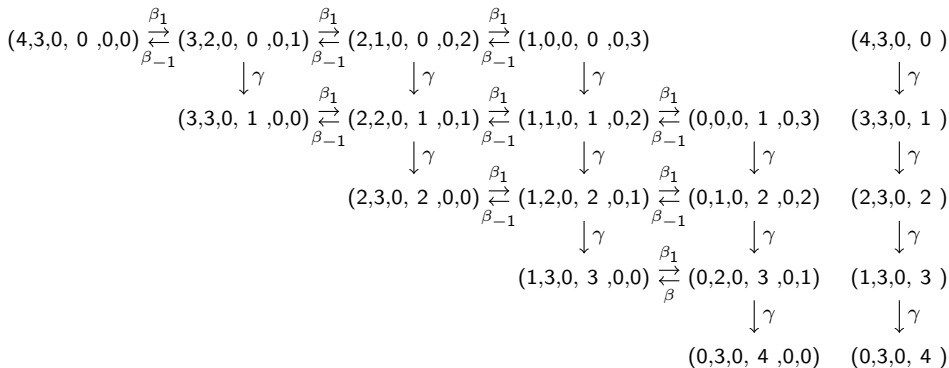
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- ▶ use in model checking with appropriate logics

Applying bisimulation to competitive inhibition



bimolecular

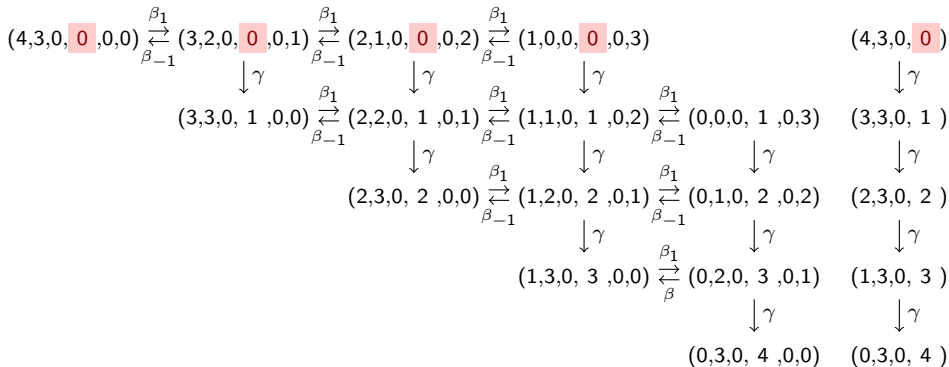
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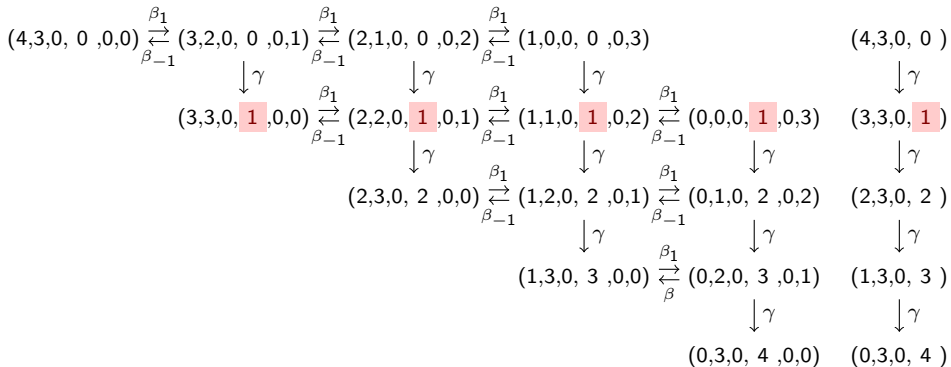
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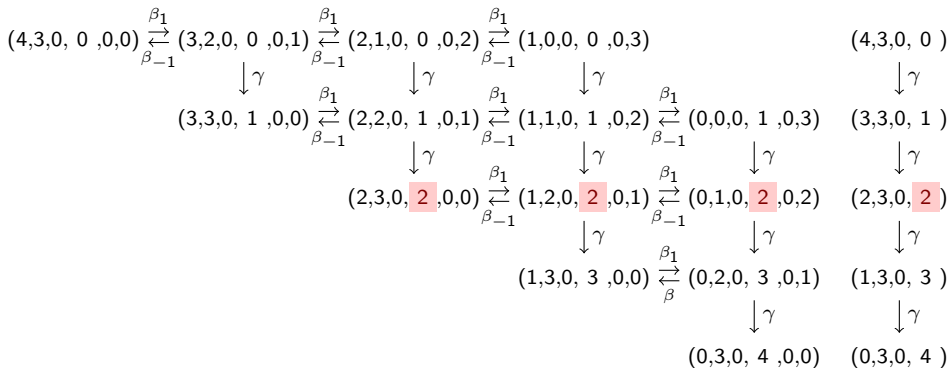
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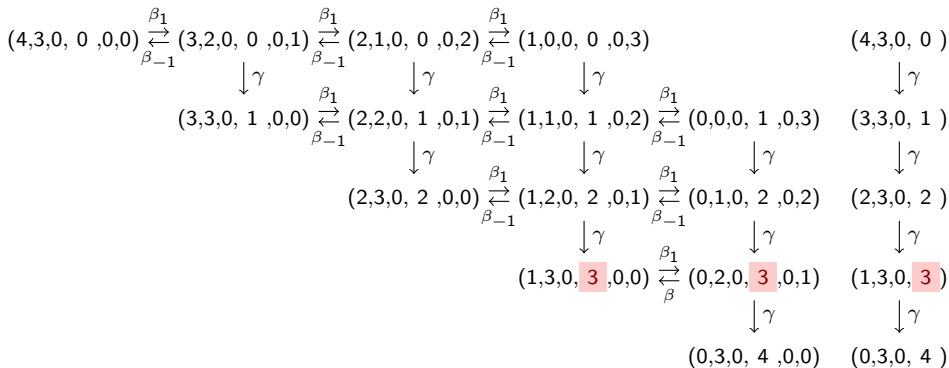
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bimolecular

abstract

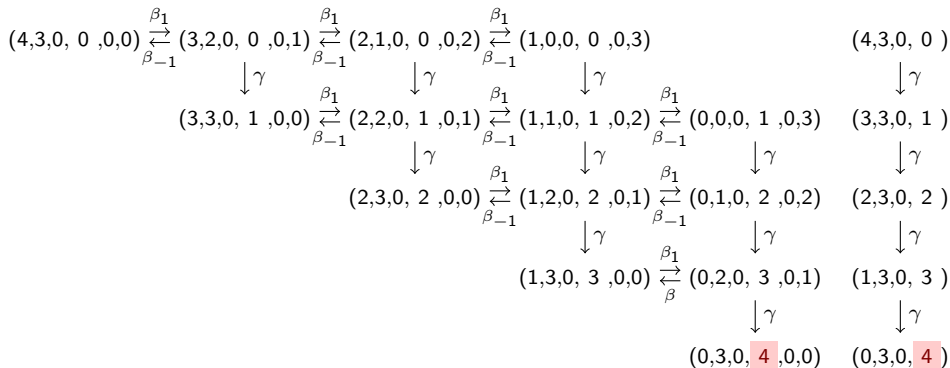
Applying bisimulation to competitive inhibition



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can this be done more efficiently by just checking slow reactions?



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- ▶ identify conserved, fast and slow variables by transforming stoichiometry matrix [Gómez-Urbe *et al*, 2008]
- ▶ result about relationship between bisimulations

M_1 has slow variables and fast variables

M_2 has same slow variables, no fast variable and $\Delta_2 = \Delta_1$

$\mathcal{R} = \{((s_1, \dots, s_n, f_1, \dots, f_m), (s_1, \dots, s_n)) \mid \text{ranges for } s_i, f_j\}$

\mathcal{R} slow bisimulation $\Rightarrow \mathcal{R}$ fast-slow bisimulation.

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Columns: reactions

groups: slow, fast

Rows: variables

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- ▶ species values can be recovered from transformed matrix

Variable classification (cont.)

- ▶ how to determine classification

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- ▶ use Bio-PEPA Eclipse Plug-in to discover potential variables
- ▶ choose independent variables
- ▶ choose sufficient fast variables so that there are the same number of variables as species

Applying classification to competitive inhibition



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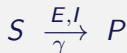


► bimolecular model

$$\begin{array}{llllll} X_{S_T} & = & S + SE + P & = & S_0 & = & n & \text{conserved} \\ X_{E_T} & = & E + EI + SE & = & E_0 & = & m & \text{conserved} \\ X_{I_T} & = & EI + I & = & I_0 & = & p & \text{conserved} \\ X_P & = & P & & & = & k & \text{slow} \\ X_{EI} & = & EI & & & = & l & \text{fast} \\ X_{SE} & = & SE & & & = & j & \text{fast} \end{array}$$

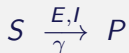


Applying classification to competitive inhibition (cont.)



- abstract model

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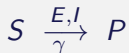


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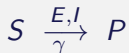
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- ignore conserved variables as they never change
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$$\Delta = \{P\}$$

Transformation of transition systems

- ▶ without loss of information, states can be transformed

bimolecular model: $(\ell_S, \ell_E, \ell_I, \ell_P, \ell_{EI}, \ell_{SE})$ to $(\ell_P, \ell_{EI}, \ell_{SE})$
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- ▶ hence M and M' are fast-slow bisimilar

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- ▶ investigate application to other process algebras

Thank you

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- ▶ how to decide which behaviours are the same?
 1. different abstractions of the same model – discretisation
 2. ideas from biology – fast/slow reactions, grouping of species
 3. existing equivalences – PEPA, bisimulation-based

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- ▶ work with a more constrained form

Well-defined Bio-PEPA systems

- well-defined Bio-PEPA species

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Well-defined Bio-PEPA systems

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$$C \stackrel{\text{def}}{=} (\alpha_1, \kappa_1) \text{op}_1 C + \dots + (\alpha_n, \kappa_n) \text{op}_n C \text{ with all } \alpha_i \text{'s distinct}$$

- ▶ well-defined Bio-PEPA model

$$P \stackrel{\text{def}}{=} C_1(\ell_1) \boxtimes_{\mathcal{L}_1} \dots \boxtimes_{\mathcal{L}_{m-1}} C_m(\ell_m) \text{ with all } C_i \text{'s distinct}$$

- ▶ well-defined Bio-PEPA system

$$\mathcal{P} = \langle \mathcal{V}, \mathcal{N}, \mathcal{K}, \mathcal{F}, \text{Comp}, P \rangle$$

Well-defined Bio-PEPA systems

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- ▶ well-defined Bio-PEPA system

$$\mathcal{P} = \langle \mathcal{V}, \mathcal{N}, \mathcal{K}, \mathcal{F}, \text{Comp}, P \rangle$$

- ▶ well-defined Bio-PEPA model component with levels
 - ▶ minimum and maximum concentrations/number of molecules
 - ▶ fix step size, convert to minimum and maximum levels
 - ▶ species S : 0 to N_S levels