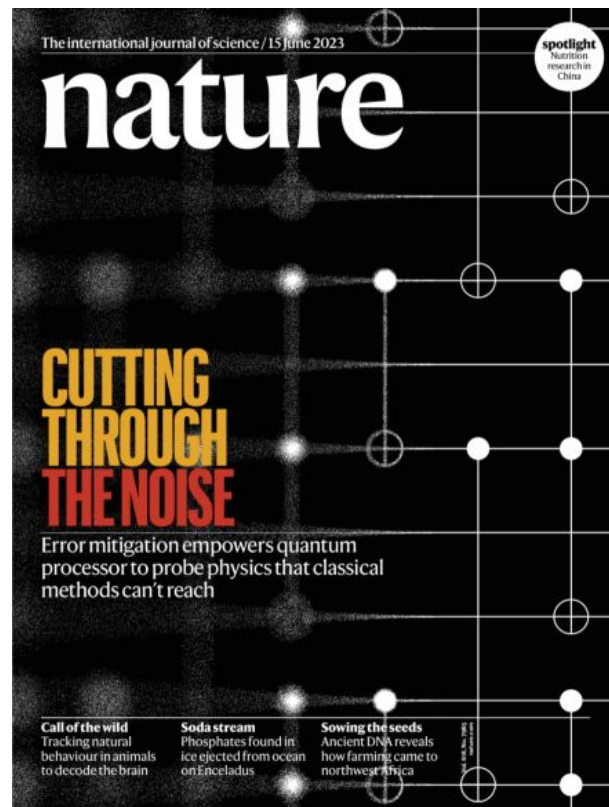


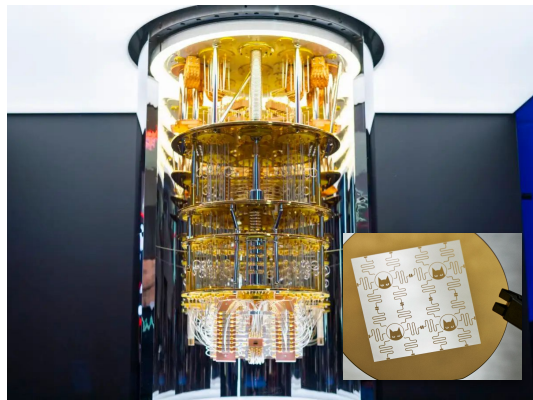
How to use a quantum computer



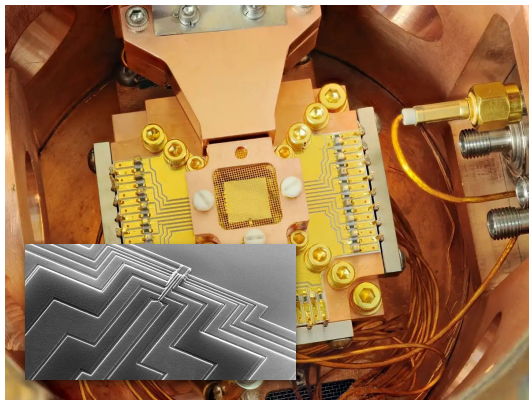
Chris Heunen

What is a quantum computer?

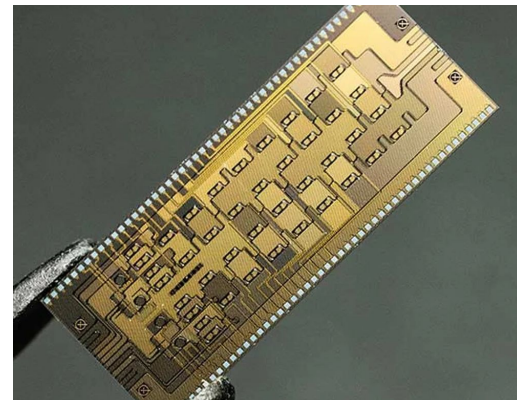




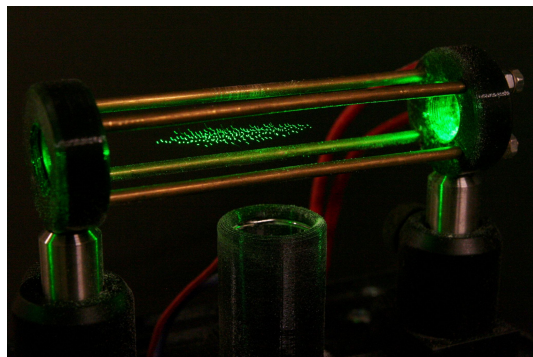
Superconducting



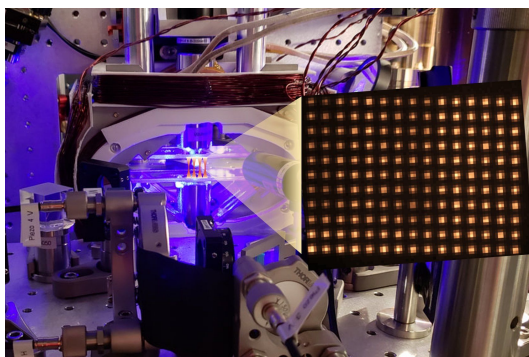
Spin



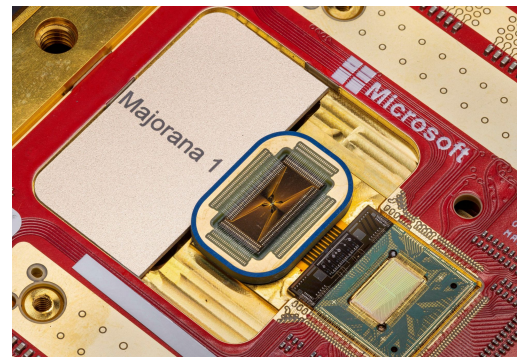
Photons



Trapped ions



Neutral atoms



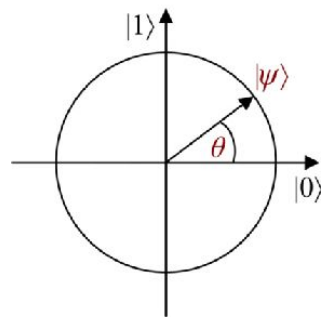
Majorana fermions

What is quantum information?



Qubits

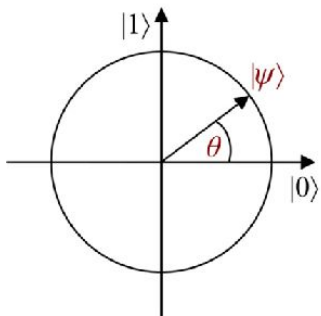
- Elements of unit circle in \mathbb{C}^2
- false = (1,0) true = (0,1)



- Also $|+\rangle = (1,1)/\sqrt{2}$ $|-\rangle = (1,-1)/\sqrt{2}$
- Manipulate by 2x2 unitary matrices
- Cannot clone or delete

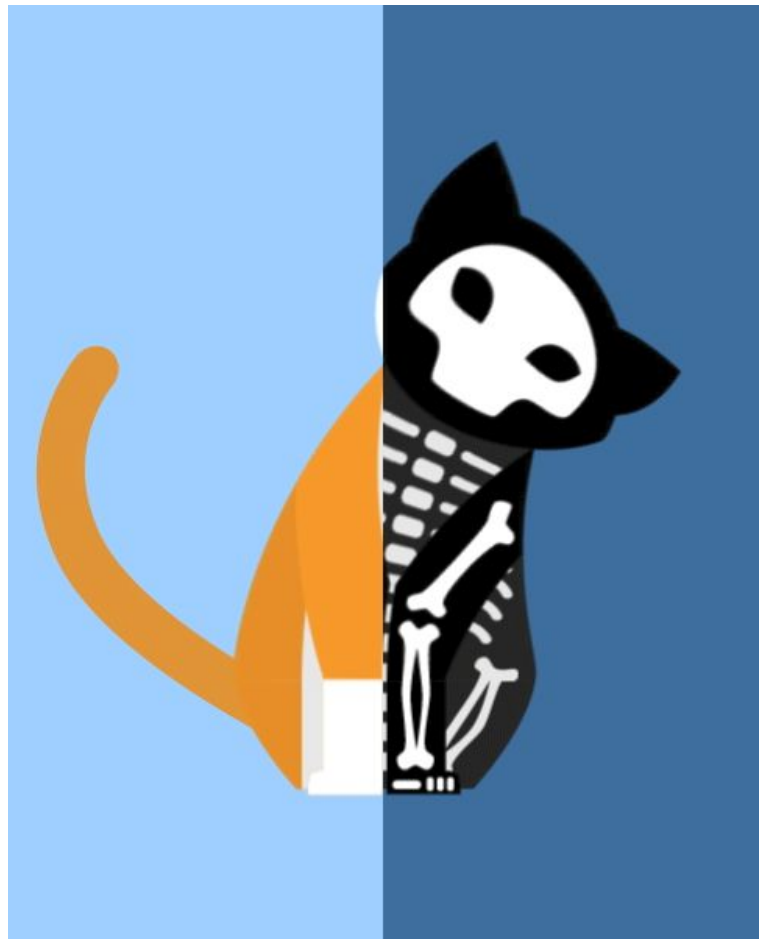
Measurement

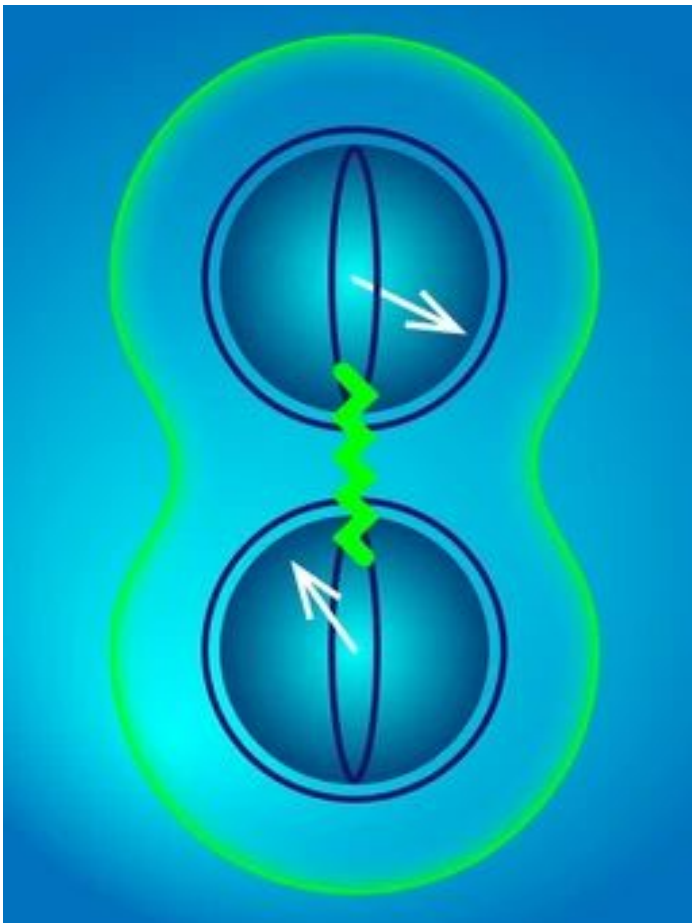
- Probabilistic operation qubit \rightarrow bit



true with probability $|\langle\psi|1\rangle|^2$

- Depends on basis/angle
- Collapse after measurement





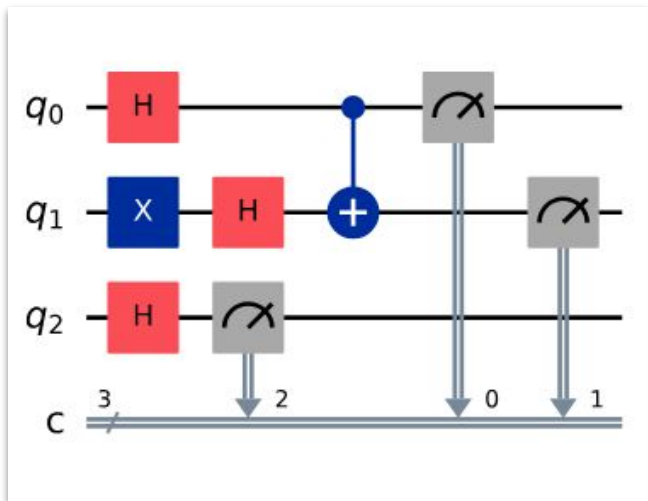
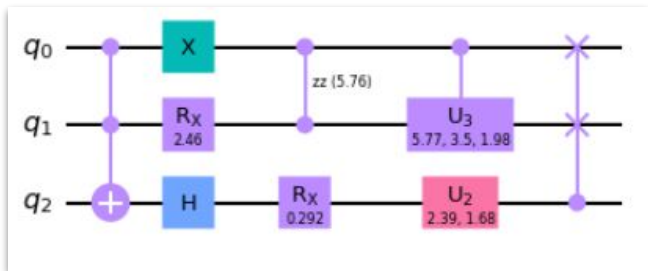
Multiple qubits

- Compound systems given by tensor product

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \otimes \begin{pmatrix} e & f \\ g & h \end{pmatrix} = \begin{pmatrix} ae & af & be & bf \\ ag & ah & bg & bh \\ ce & cf & de & df \\ cg & ch & dg & dh \end{pmatrix}$$

- State not determined by factors
- Dimension multiplies
- Measurement of one qubit influences other

How not to use a quantum computer



Circuits

- Time flows left to right
- Space goes up and down
- Qubits undergo unitary gates
- Gates can be controlled by multiple qubits
- Bits may influence control flow

$$* \left(H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \quad X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix} \quad CX = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \right)$$

Universal:

- Computational universality:
can compute every computable function

But:

- Hard to verify
- Hard to manipulate
- Hard to discover algorithms
- Hard to scale
- Hard to reuse classical infrastructure



Scale

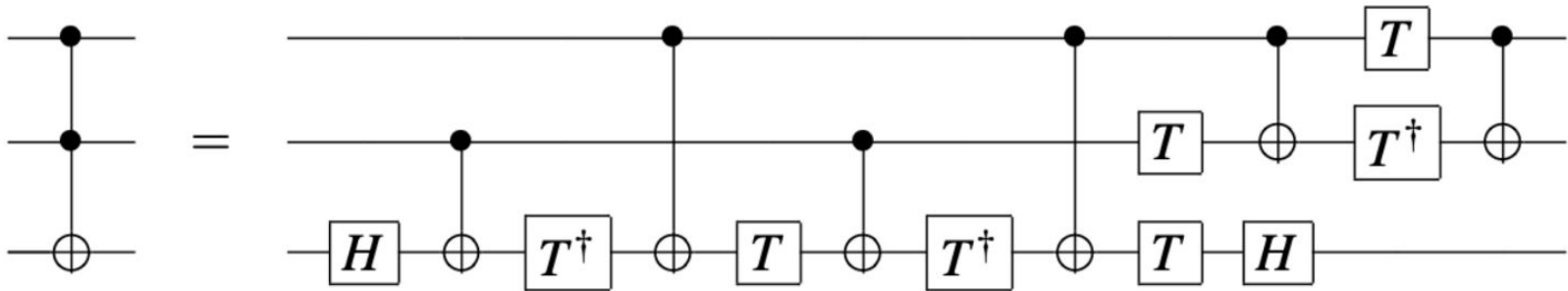
This circuit adds two 8-bit numbers:



Reasoning

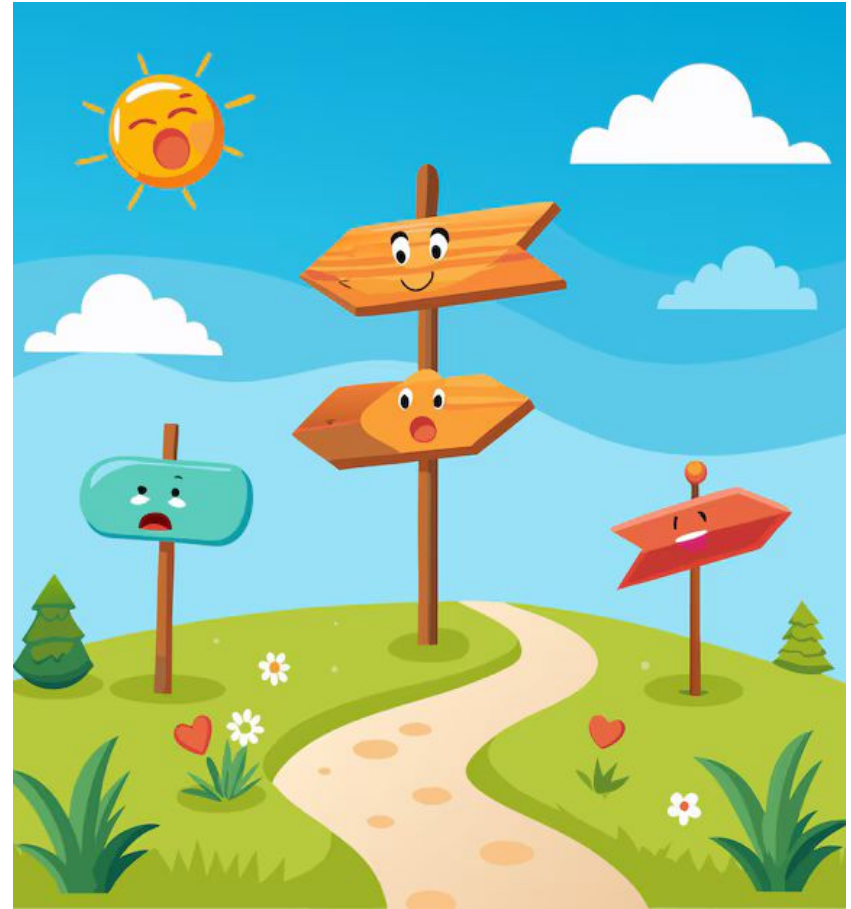
Toffoli gate has many 'obvious' properties ...

... that are completely obscured once expressed by elementary gates

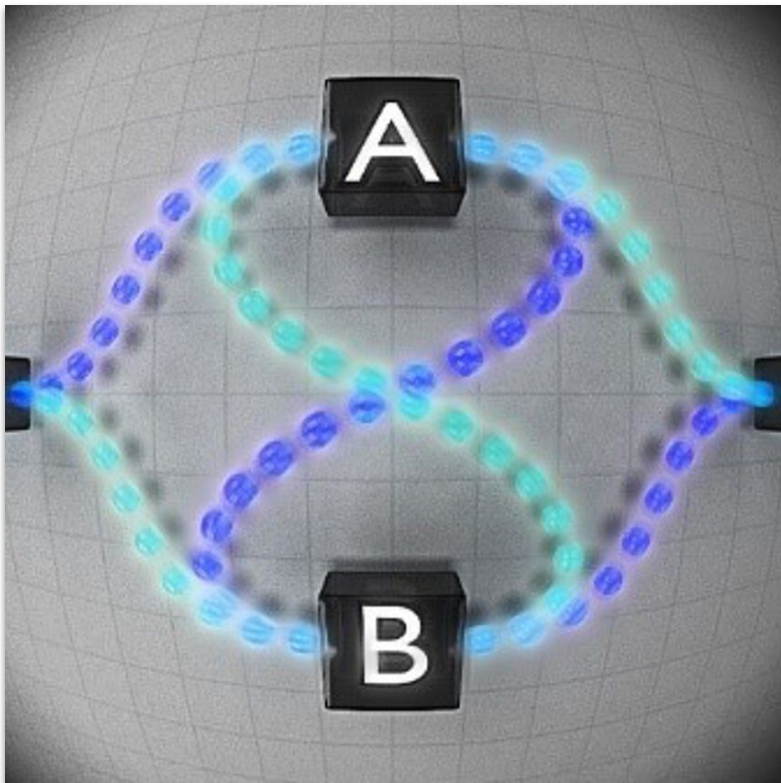


Quantum conditional

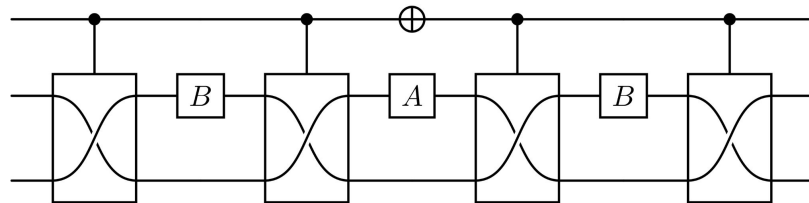
- $\llbracket \text{if } q \text{ then } U \text{ else } V \rrbracket$
 $= \begin{pmatrix} \llbracket U \rrbracket & 0 \\ 0 & \llbracket V \rrbracket \end{pmatrix}$
- Both branches in superposition
- Not monotone
(unlike probabilistic computing)



Coherent Control



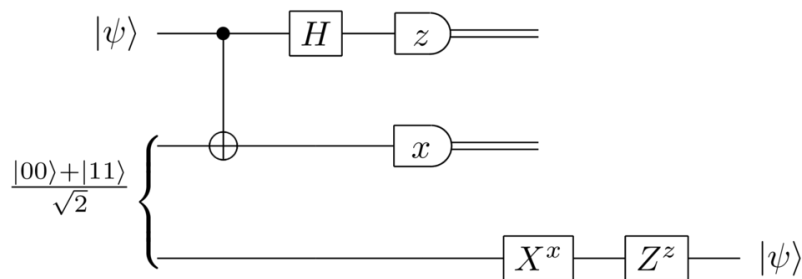
- Quantum switch:
if q then AB else BA
- Possible as circuit
(with multiple uses of oracles A and B)



- Possible experimentally
(with single use of oracles)
- Impossible as circuit
(with single use of oracles)

Teleportation

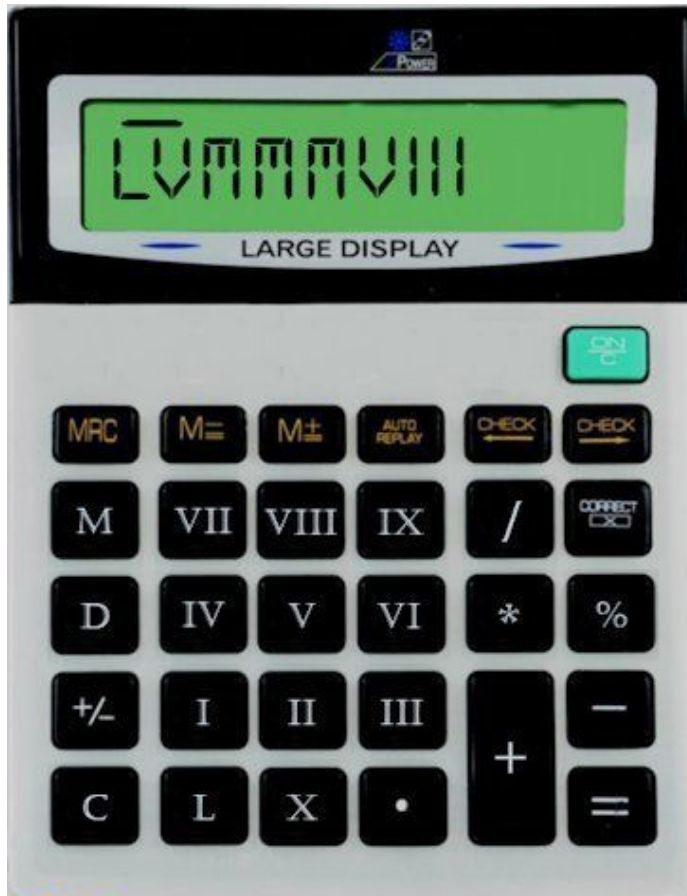
- Use entanglement as communication channel



- Transmit 2 bits (finite amount of information) to teleport 1 qubit (uncountable)
- Pay cost ahead of time



How to use quantum information



Notation

Notation isn't just a way to write ideas.

It stimulates ideas you can have.

Abstraction

- *Scalable*
- *Automatable*
- *Optimisable*
- *Understandable*



Abstraction

- *Scalable*
- *Automatable*
- *Optimisable*
- *Understandable*



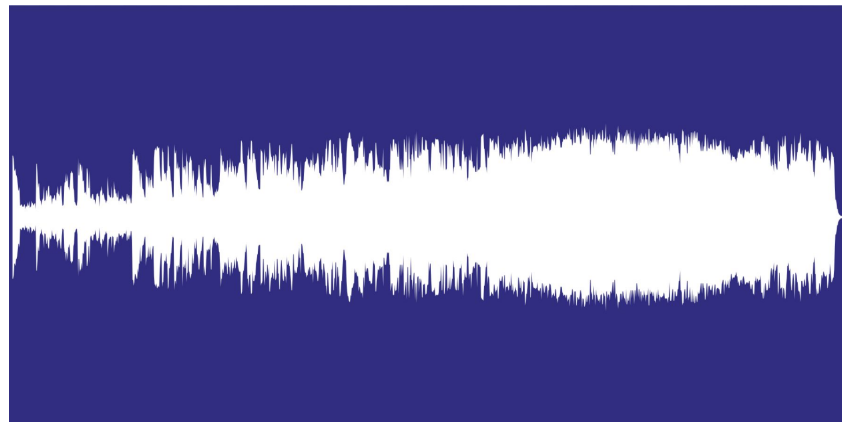
Abstraction

- *Scalable*
- *Automatable*
- *Optimisable*
- *Understandable*

AULD LANG SYNE

ROBERT BURNS

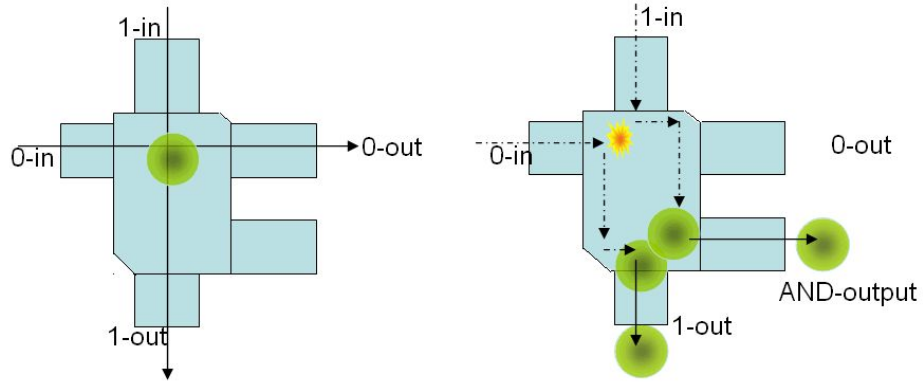
Scotch Air

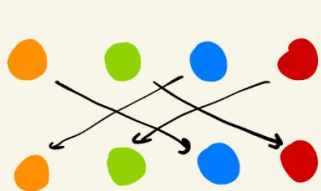


How to use a quantum computer

Billiard Ball Computing

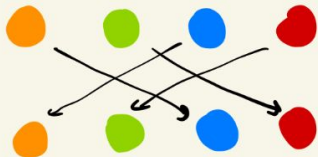
Universal for reversible computation [Fredkin & Toffoli, 1982]



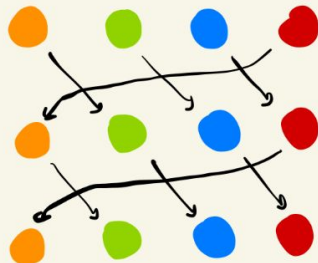


=

$$\begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}$$



=



=

??

Permutations

Permutation with n elements is $n \times n$ matrix with boolean entries

Permutations give semantics for **reversible classical programs**

Not all polynomials $x^2 + f$ have solutions:
not every permutation has a square root

Rig categories

Categorification of natural numbers: finite sets and bijections

$$\frac{f: A \rightarrow B \quad g: B \rightarrow C}{g \circ f: A \rightarrow C}$$

$$\overline{\text{id}: A \rightarrow A}$$

$$(h \circ g) \circ f = h \circ (g \circ f)$$

$$\text{id} \circ f = f = f \circ \text{id}$$

$$\frac{f: A \rightarrow B \quad f': A' \rightarrow B'}{f \otimes f': A \otimes A' \rightarrow B \otimes B'}$$

$$(A \otimes B) \otimes C \simeq A \otimes (B \otimes C)$$

$$I \otimes A \simeq A \simeq A \otimes I$$

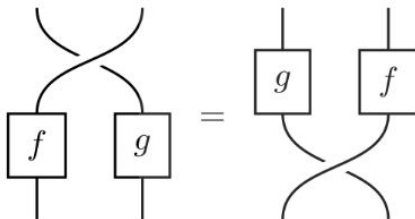
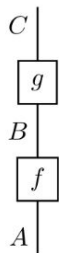
$$A \otimes B \simeq B \otimes A$$

$$\frac{f: A \rightarrow B \quad f': A' \rightarrow B'}{f \oplus f': A \oplus A' \rightarrow B \oplus B'}$$

$$(A \oplus B) \oplus C \simeq A \oplus (B \oplus C)$$

$$O \oplus A \simeq A \simeq A \oplus O$$

$$A \oplus B \simeq B \oplus A$$

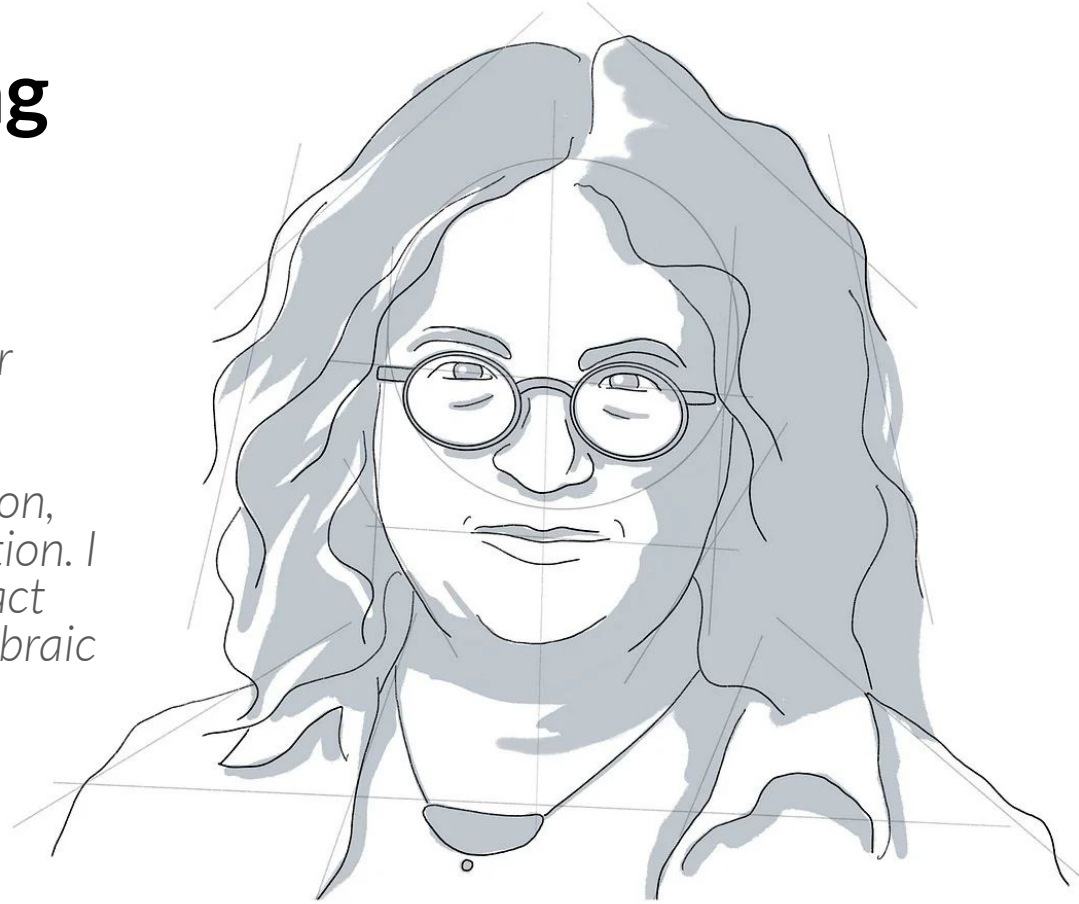


$$A \otimes (B \oplus C) \simeq (A \otimes B) \oplus (A \otimes C)$$

$$A \otimes O \simeq O$$

Quantum computing as a completion

"It's really something that is special for quantum computation because it's somehow 'complete' — quantum computation is some kind of completion, mathematically, of classical computation. I think of this as maybe similar to the fact that the complex numbers are an algebraic closure of the real numbers."





Universality

Categorical universality
*characterise property up to isomorphism
by behaviour rather than construction*

$$\begin{array}{ccc} V \times W & \longrightarrow & V \otimes W \\ & \searrow \text{bilinear} & \downarrow \text{linear} \\ & & Z \end{array}$$

Finite sets and permutations are initial rig category

$$\begin{array}{ccc} \text{Axioms} & \longrightarrow & \text{Free model} \\ & \searrow & \downarrow \exists! \\ & & \text{Any model} \end{array}$$

A Few Square Roots

Add two generators

$$w: 1 \leftrightarrow 1$$

$$v: 1+1 \leftrightarrow 1+1$$

And impose three equations

$$v^2 = \text{swap}$$

$$w^8 = \text{id}$$

$$vsv = sv$$



where $s = \text{id} + w^2$



Can build Clifford+T: $T = \text{id} + w$ $S = \text{id} + w^2$ $Z = \text{id} + w^4$ $H = w^7 v s v$



Freeq

Theorem: Free model Π exists

Theorem: If \mathbb{D} is dyadic rationals, ζ is 8th root of unity, then $\Pi = \text{Unitary}(\mathbb{D}[\zeta])$

Theorem: There is inclusion $\llbracket - \rrbracket: \Pi \rightarrow \text{Unitary}(\mathbb{C})$, and it is dense

Theorem: $\llbracket f \rrbracket = \llbracket g \rrbracket$ iff $\langle\langle f \rangle\rangle = \langle\langle g \rangle\rangle$ for all interpretations $\langle\langle f \rangle\rangle$

Theorem: There is faithful $F(\Pi) \rightarrow \text{FCstar}_{\text{cp}}$ for universal construction F

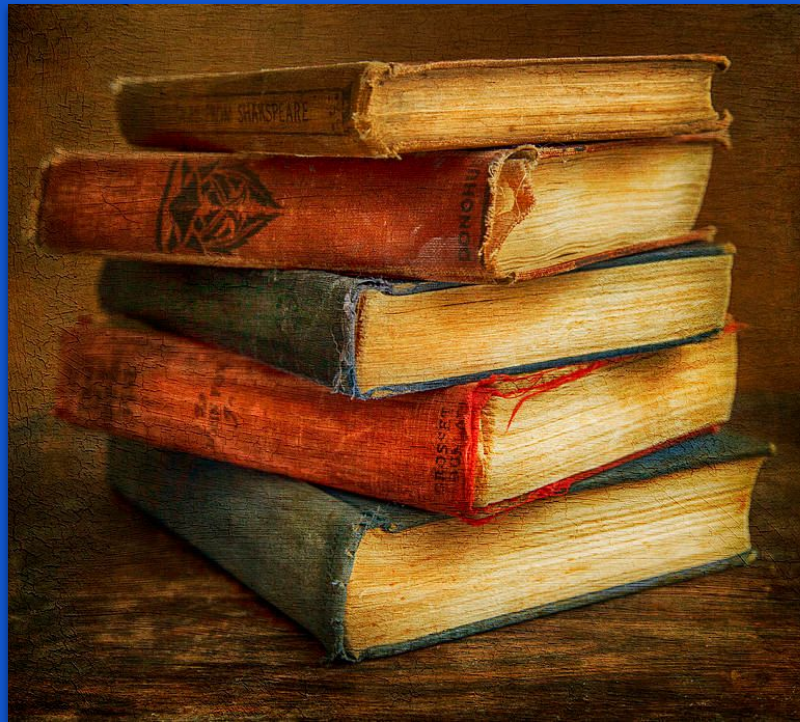
Conclusion

- How to use a quantum computer?
I don't know.
- If scalable, optimisable, and understandable,
then programming language must be abstract.
- Universality: no alarms and no surprises.
Can discover primitives and algorithms.
- Can uncover nature of quantum information.
- Suggestion: rig categories, square roots.
Next: syntax, optimisation. Hamiltonians.

Many exciting questions for next 100 years!



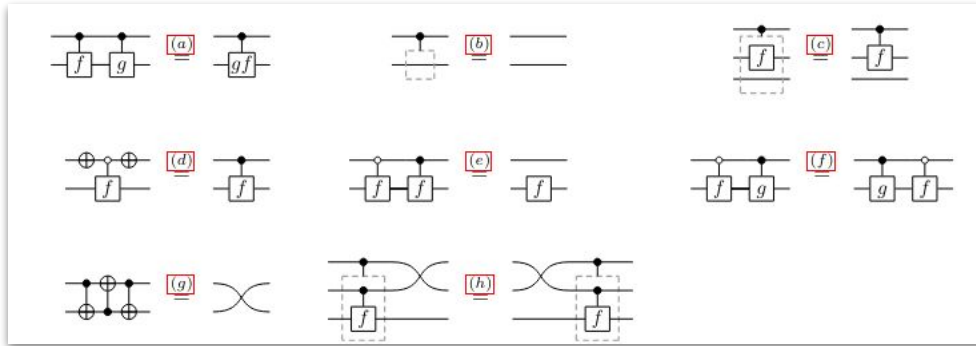
References



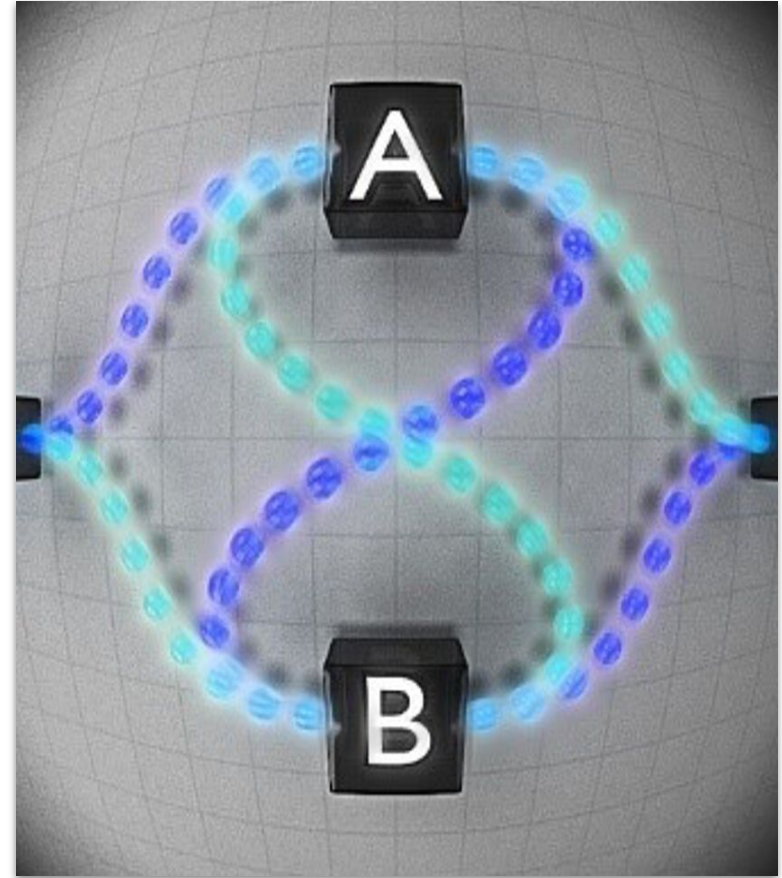
- “Categories for Quantum Theory”
Oxford University Press, 2019
- “Axioms for the Category of Hilbert Spaces”
Proceedings of the National Academy of Sciences, 2022
- “Weakly Measured While Loops: Peeking at Quantum States”
Quantum Science and Technology, 2022
- “With a Few Square Roots, Quantum Computing is as Easy as π ”
Principles of Programming Languages, 2024
- “Qurts: Automatic Quantum Uncomputation by Affine Types with Lifetimes”
Principles of Programming Languages, 2025
- “Quantum Circuits are Just a Phase”
Principles of Programming Languages, 2026
- “One Rig to Control Them All”
arXiv, 2026

Taking back control

- Start with any 'circuit theory'
- Add control



- Get complete, structured language



It's just a phase

- Many (most?) quantum algorithms come down to eigenspace decomposition and eigenvector manipulation
- Lift to primitive
- Universal
- Grover now one-liner:

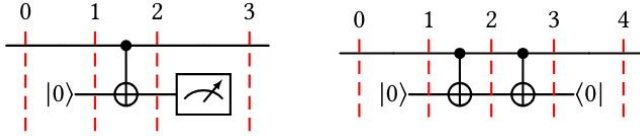
$Z := \text{if let } 1\rangle \text{ then Ph}(\pi)$	$X := \text{if let } -\rangle \text{ then Ph}(\pi)$
$T := \text{if let } 1\rangle \text{ then Ph}(\pi/4)$	$Y := \text{if let } S \cdot -\rangle \text{ then Ph}(\pi)$
$H := \text{if let } Y^{1/4} \cdot 1\rangle \text{ then Ph}(\pi)$	$CX := \text{if let } 1\rangle \otimes \text{id}_1 \text{ then } X$

$\text{if let } |\omega_1\rangle \otimes \cdots \otimes |\omega_n\rangle \text{ then Ph}(\pi)$



More ancillae, more problems

- Cannot reuse dirty qubits

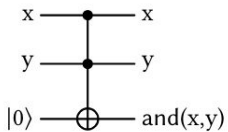


- Rust type system: ownership, borrowing

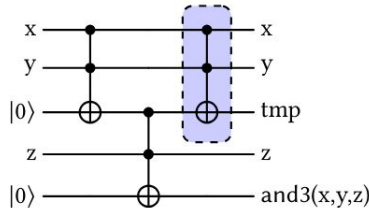
```
1 fn main() {
2   let x = [1,2,3];
3   let y = f(x); // value of x is moved
4   print!("{}",{x},{y}); // so x lost access
5 }
6 fn f<T>(x : T) -> T {
7   x
8 }
```

```
1 fn main() {
2   let x = [1,2,3];
3   let y = f(&x); // x is borrowed
4   print!("{}",{x},{y}); // x can be read
5 }
6 fn f<T>(x : T) -> T {
7   x
8 }
```

- Compiler automatically inserts uncomputation



```
1 fn and3<'a>(<
2   x:&'a qbit,
3   y:&'a qbit,
4   z:&'a qbit
5 ) -> #'a qbit {
6   let tmp = and(x,y);
7   let ref = &tmp;
8   and(ref,z)
9 }
```



Becoming measured

- classical computation
= classical reversible computation
+ information effects
- Can copy and delete classical bits with
erase : $b \rightsquigarrow 1$
create : $1 \rightsquigarrow b$
- Dynamic quantum programming language
- Compiles measurements anywhere to
single standard measurement at end



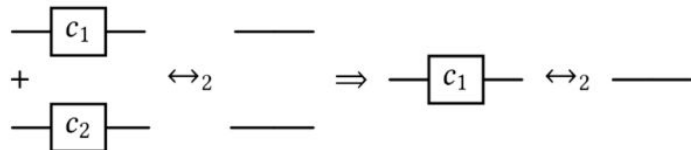
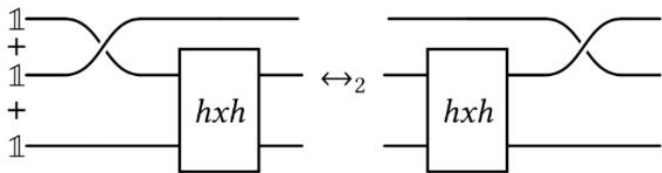
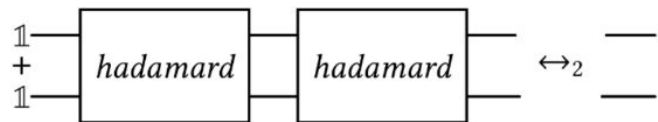
Syntax

$iso ::= \dots \mid hadamard$

Types

$hadamard : \mathbb{1} + \mathbb{1} \leftrightarrow \mathbb{1} + \mathbb{1}$

Equations



Programming with equations

- Word problem decidable
- Normalisation by evaluation
- Fewer generators than Π
- **Theorem (completeness):** $\llbracket f \rrbracket = \llbracket g \rrbracket$ iff $\langle\langle f \rangle\rangle = \langle\langle g \rangle\rangle$

Orthogonal group presentations

Proposition: there is finite presentation for unitary matrices $O_n(\mathbb{Z}[1/\sqrt{2}])$

$$(-1)_{[a]}^2 \approx \varepsilon$$

$$X_{[a,b]}^2 \approx \varepsilon$$

$$(-1)_{[a]}(-1)_{[b]} \approx (-1)_{[b]}(-1)_{[a]}$$

$$(-1)_{[a]}X_{[b,c]} \approx X_{[b,c]}(-1)_{[a]}$$

$$X_{[a,b]}X_{[c,d]} \approx X_{[c,d]}X_{[a,b]}$$

$$(-1)_{[a]}X_{[a,b]} \approx X_{[a,b]}(-1)_{[b]}$$

$$X_{[b,c]}X_{[a,b]} \approx X_{[a,b]}X_{[a,c]}$$

$$X_{[a,c]}X_{[b,c]} \approx X_{[b,c]}X_{[a,b]}$$

$$H_{[a,b]}^2 \approx \varepsilon$$

$$(-1)_{[a]}H_{[b,c]} \approx H_{[b,c]}(-1)_{[a]}$$

$$X_{[a,b]}H_{[c,d]} \approx H_{[c,d]}X_{[a,b]}$$

$$H_{[a,b]}H_{[c,d]} \approx H_{[c,d]}H_{[a,b]}$$

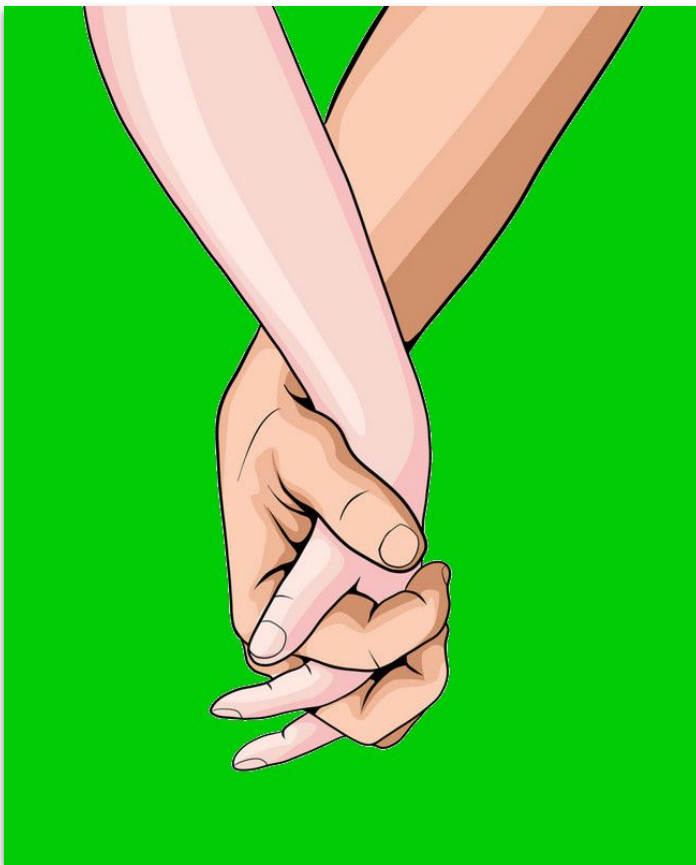
$$H_{[b,c]}X_{[a,b]} \approx X_{[a,b]}H_{[a,c]}$$

$$H_{[a,c]}X_{[b,c]} \approx X_{[b,c]}H_{[a,b]}$$

$$(-1)_{[a]}(-1)_{[b]}H_{[a,b]} \approx H_{[a,b]}(-1)_{[a]}(-1)_{[b]}$$

$$(-1)_{[b]}H_{[a,b]} \approx H_{[a,b]}X_{[a,b]}$$

$$X_{[b,c]}H_{[a,b]}X_{[a,c]}H_{[a,d]}H_{[a,b]}X_{[a,c]}H_{[a,d]} \approx H_{[a,b]}X_{[a,c]}H_{[a,d]}H_{[a,b]}X_{[a,c]}H_{[a,d]}X_{[c,d]}$$



Bit commitment

- Alice commits to value hidden from Bob
- Cryptographic primitive, essential in
 - zero-knowledge proofs
 - secret sharing
 - secure multi-party computation
- Impossible with quantum values