

# Advanced Combinatorics - 2016 Fall

## Solutions to Exercise 2

Comments and corrections are welcome.

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1. Prove that if  $G$  is a graph of order  $n \geq 3$  such that  $d(u)+d(v) \geq n$  for any  $(u, v) \notin E(G)$ , then  $G$  is Hamiltonian.

**Solution:** It is pretty much the same proof as the Dirac's theorem. For the connectedness, pick  $u$  and  $v$  from two different components, then  $d(u) + d(v) \leq n - 2$ . Contradiction.

The rest of the proof is almost identical. The only thing to note is that  $x_0$  and  $x_\ell$  cannot be adjacent, as otherwise we either have a Hamiltonian cycle or can construct a longer path.

2. Show that Dirac's theorem is the best possible in terms of minimal degrees. In other words, for any odd  $n \geq 3$ , construct a graph  $G$  such that  $\delta(G) = \frac{n-1}{2}$  and  $G$  is not Hamiltonian.

**Solution:** Note that we want to avoid a cycle of length  $n$ , which is odd. It is natural to construct a bipartite graph. Then, to spread degrees evenly, consider  $K_{\lfloor n/2 \rfloor, \lfloor n/2 \rfloor}$ .

3. Let  $k < n$  be two positive integers. Let  $G$  be a connected graph of order  $n$  such that  $\delta(G) \geq \frac{k}{2}$ . Prove that  $G$  must have a cycle of length  $\lceil \frac{k}{2} \rceil + 1$ , and this is tight.

**Solution:** Once again, consider the maximal length path  $P_\ell = \{x_0, x_1, \dots, x_\ell\}$ . Consider  $\Gamma(x_0)$ . By maximality,  $\Gamma(x_0) \subset P_\ell \setminus \{x_0\}$ . Since  $|\Gamma(x_0)| \geq k/2$ , there exists a  $t \geq \lceil k/2 \rceil$  such that  $x_t \in \Gamma(x_0)$  (otherwise there is not enough "room" for  $\Gamma(x_0)$ ). This creates a cycle of length  $\lceil k/2 \rceil + 1$ .

This is best possible. Consider a bunch of copies of  $K_{\lceil k/2 \rceil + 1}$ . This graph satisfies the degree constraint and has no cycles of length  $\lceil \frac{k}{2} \rceil + 2$ . However it is not connected. To make it connected without creating cycles, connect these copies by one edge in a tree-like structure.

In technical terms, we "contract" each  $K_{\lceil k/2 \rceil + 1}$  into a vertex. Then construct a spanning tree.