

Average case local to global and applications

Thuy-Duong Vuong

Based on joint works with Yeganeh Alimohammadi, Nima Anari, Vishesh Jain, Yang P. Liu, Frederic Koehler, Huy Tuan Pham, Kirankumar Shiragur

Sampling from distributions

Density function $\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$, $\mathbb{P}_\mu(x) = \mu(x) / \sum_y \mu(y) \propto \mu(x)$

Approximate sampling: sample from $\hat{\mu}$ s.t.

$$d_{TV}(\mu, \hat{\mu}) = \frac{1}{2} \sum_x |\mathbb{P}_\mu(x) - \mathbb{P}_{\hat{\mu}}(x)| < 1/\text{poly}(n)$$

Counting: compute partition function $\sum_y \mu(y)$

Approximate sampling \Leftrightarrow approximate counting

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Approximate sampling \Leftrightarrow approximate counting

Spin system: $Q^n \equiv \binom{[n] \times Q}{n}$

Overview

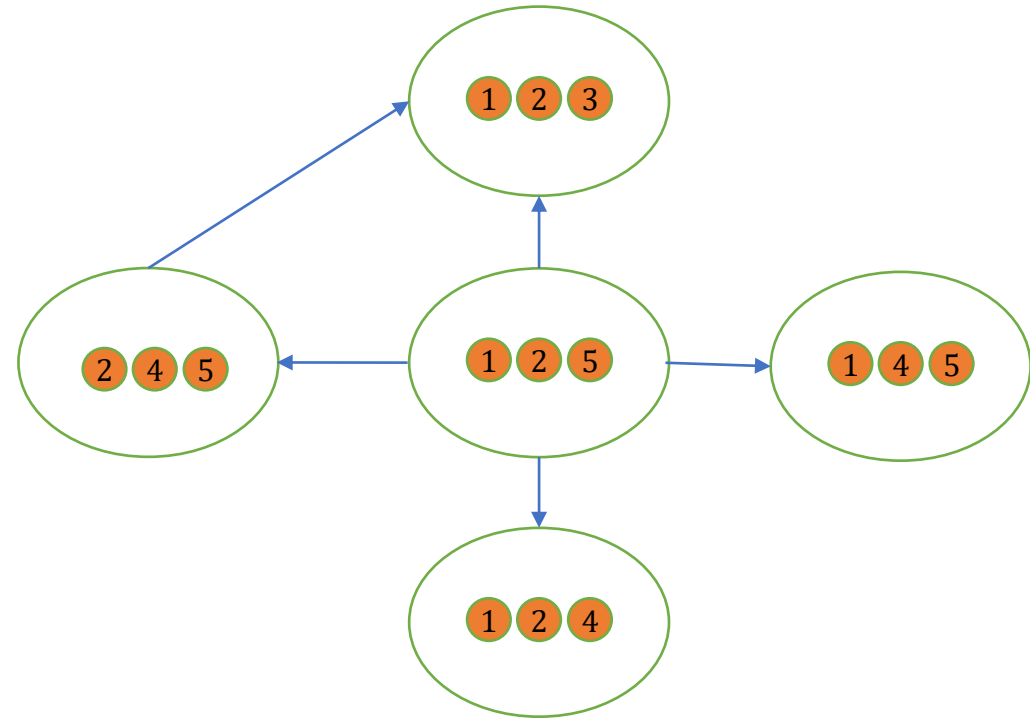
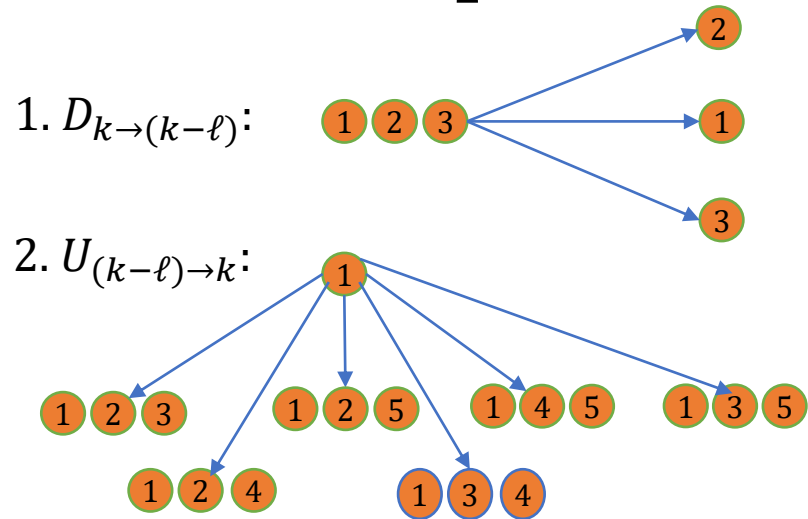
1. Background

- Functional inequality and mixing time
- Local to global: an inductive approach to prove functional inequality
- Building blocks: entropic/spectral independence

2. Average-case local to global

- Definition
- Application: p-spin system
- Application: spanning trees and strongly Rayleigh distributions

Multi-steps down-up walk for $\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$



$$P = D_{k \rightarrow (k-\ell)} U_{(k-\ell) \rightarrow k}$$

Glauber dynamics: $P = D_{n \rightarrow (n-1)} U_{(n-1) \rightarrow n}$

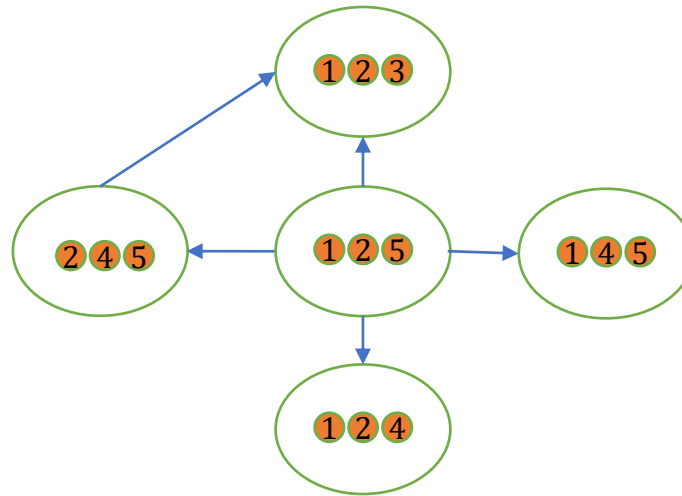
Variance/entropy contraction vs. mixing time

Variance contraction

- $\chi^2(\nu P || \mu P) \leq (1 - \rho_{\chi^2}) \chi^2(\nu || \mu)$
- $\rho_{\chi^2} = 1 - \lambda_2(P)$

Entropy contraction

- $\mathcal{D}_{KL}(\nu P || \mu P) \leq (1 - \rho_{KL}) \mathcal{D}_{KL}(\nu || \mu)$



$$P = D_{k \rightarrow (k-\ell)} U_{(k-\ell) \rightarrow k}$$

Variance/entropy contraction vs. mixing time

Variance contraction

- $\chi^2(\nu P || \mu P) \leq (1 - \rho_{\chi^2}) \chi^2(\nu || \mu)$
- $T_{mix} \leq \rho_{\chi^2}^{-1} \log(\min \mu(x))^{-1} \approx \rho_{\chi^2}^{-1} n$
- $\rho_{\chi^2} = 1 - \lambda_2(P)$

Entropy contraction

- $\mathcal{D}_{KL}(\nu P || \mu P) \leq (1 - \rho_{KL}) \mathcal{D}_{KL}(\nu || \mu)$
- $T_{mix} \leq \rho_{KL}^{-1} \log \log(\min \mu(x))^{-1} \approx \rho_{KL}^{-1} \log n$

Entropy contraction implies variance contraction: $\rho_{\chi^2} \geq \rho_{KL} \geq 0$

Typically for Glauber dynamics: $\rho_{\chi^2} \approx \rho_{KL} = \frac{1}{n}$

but $\log \min \mu(x)^{-1} \approx n$. Bounding $\rho_{KL} \Rightarrow$ quadratic improvement on T_{mix}

It is hard to bound ρ_{KL} !

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Local-to-global

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$

μ is reducible i.e. conditionals $\mu(\cdot | S): \binom{[n]}{k-|S|} \rightarrow \mathbb{R}_{\geq 0}$ of μ has same property as μ
 $\mu(S' | S) \propto \mu(S' \cup S)$

Local-to-global (variation of [Alev-Lau—STOC'21])

Entropy contraction

$$\forall \nu : \mathcal{D}_{\text{KL}}(\nu || \mu) \geq (1 + \gamma_k) \mathcal{D}_{\text{KL}}(\nu D_{k \rightarrow (k-\ell)} || \mu D_{k \rightarrow (k-\ell)})$$

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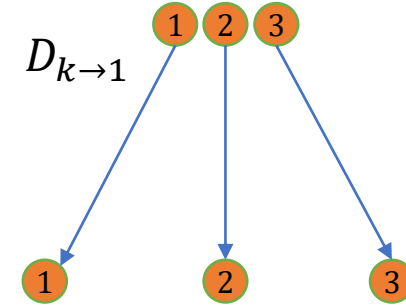
Inductive hypothesis on $\mu(\cdot | i): \binom{[n] \setminus \{i\}}{k-1} \rightarrow \mathbb{R}_{\geq 0}, \mu(S|i) \propto \mu(S)$

$$\mathcal{D}_{\text{KL}}(\nu(\cdot|i)||\mu(\cdot|i)) \geq (1 + \gamma_{k-1}) \mathcal{D}_{\text{KL}}(\nu(\cdot|i) D_{(k-1) \rightarrow (k-\ell-1)} || \mu D_{(k-1) \rightarrow (k-\ell-1)})$$

Local-to-global

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$

$D_{k \rightarrow 1}(S)$: sample $i \in S$ uniformly



$$\nu D_{k \rightarrow 1}(i) = \text{marginal of } i \propto \mathbb{P}_{S \sim \nu}[i \in S] = \nu(i)$$

Local-to-global

Entropy contraction

$$\forall \nu : \mathcal{D}_{\text{KL}}(\nu \| \mu) \geq (1 + \gamma_k) \mathcal{D}_{\text{KL}}(\nu D_{k \rightarrow (k-\ell)} \| \mu D_{k \rightarrow (k-\ell)})$$

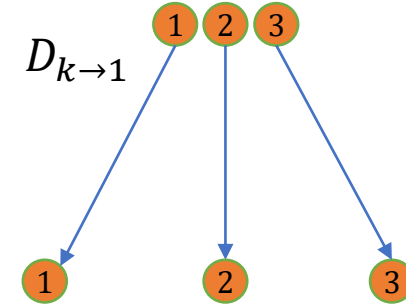
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$$\sum_i \nu(i) (\mathcal{D}_{\text{KL}}(\nu(\cdot | i) \| \mu(\cdot | i)) - (1 + \gamma_{k-1}) \mathcal{D}_{\text{KL}}(\nu(\cdot | i) D_{(k-1) \rightarrow (k-\ell-1)} \| \mu D_{(k-1) \rightarrow (k-\ell-1)})) \geq 0$$

Local-to-global

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$



$D_{k \rightarrow 1}(S)$: sample $i \in S$ uniformly

$\frac{1}{\alpha}$ -entropic independence $\Leftrightarrow \forall \nu: \mathcal{D}_{KL}(\nu || \mu) \geq \alpha k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$

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Local-to-global

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$$\mathcal{D}_{KL}(\nu || \mu) + \gamma_{k-1} \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1}) \geq (1 + \gamma_{k-1}) \mathcal{D}_{KL}(\nu D_{k \rightarrow (k-\ell)} || \mu D_{k \rightarrow (k-\ell)})$$

Overview

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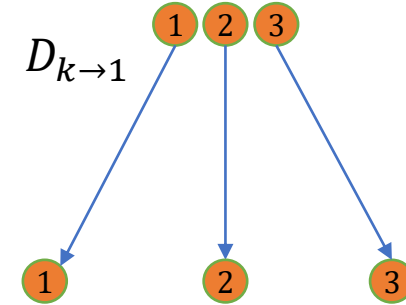
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Entropic independence

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$



$D_{k \rightarrow 1}(S)$: sample $i \in S$ uniformly

$\frac{1}{\alpha}$ -**entropic** independence $\Leftrightarrow \forall \nu: \mathcal{D}_{KL}(\nu || \mu) \geq \alpha k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$

Thm: If $\mu_S = \mu(\cdot | S)$ are $\frac{1}{\alpha}$ -entropic independence $\forall S$

\Rightarrow **entropy** contraction of $D_{k \rightarrow (k-\ell)}$

\Rightarrow optimal bound for **modified Log-Sobolev** constant and **mixing time** of down-up walks

Spectral independence

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$

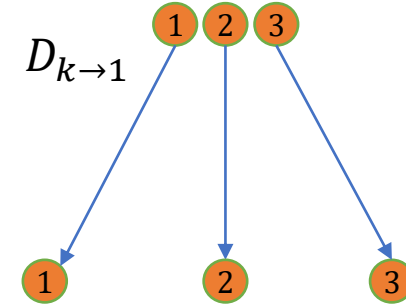
$D_{k \rightarrow 1}(S)$: sample $i \in S$ uniformly

$$\frac{1}{\alpha}\text{-spectral independence} \Leftrightarrow \forall v: \mathcal{D}_{x^2}(v || \mu) \geq \alpha k \mathcal{D}_{x^2}(v D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$$

Thm: If $\mu_S = \mu(\cdot | S)$ are $\frac{1}{\alpha}$ -spectral independence $\forall S$

\Rightarrow **variance** contraction of $D_{k \rightarrow (k-\ell)}$

\Rightarrow optimal bound for **spectral gap/Poincare** constant of down-up walks



What if $\mu(\cdot | S)$ is not entropically independent for some S ?

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Average local-to-global

Thm [Alimohammadi-Anari-Shiragur-V'21, Anari-Liu-V'22]: If $\mu_T = \mu(\cdot | T)$ are $(k - |T|)(1 - \rho(T))$ -spectral independence $\forall T$

$$(1 - \rho(T)) \mathcal{D}_{KL}(\nu || \mu_T) \geq \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu_T D_{k \rightarrow 1})$$

Harmonic mean:

$$\gamma_T = \mathbb{E}_{e_1, \dots, e_{|T|} \sim \text{permutation}(T)} \left[\left(\rho(\emptyset) \rho(\{e_1\}) \dots \rho(\{e_1, \dots, e_{|T|-1}\}) \right)^{-1} \right]^{-1}$$

Then for $\kappa = \min\{\gamma_T | T \in \binom{[n]}{\ell}\}$

$$(1 - \kappa) \mathcal{D}_{KL}(\nu || \mu) \geq \mathcal{D}_{KL}(\nu D_{k \rightarrow (k-\ell)} || \mu D_{k \rightarrow (k-\ell)})$$

Average local to global

Intuition:

For each **fixed** set $T \in \binom{[n]}{\ell}$ and s ,

If for **average** $S \in \binom{T}{s}$, μ_S is entropically independent
then $D_{k \rightarrow (k-\ell)}$ has good entropy contraction

Average local to global

Intuition:

~~For each **fixed** set $T \in \binom{[n]}{\ell}$ and s ,~~

If for **average** $S \in \binom{T}{s}$, μ_S is entropically independent
then $D_{k \rightarrow (k-\ell)}$ has good entropy contraction?

Not true!

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Spin system

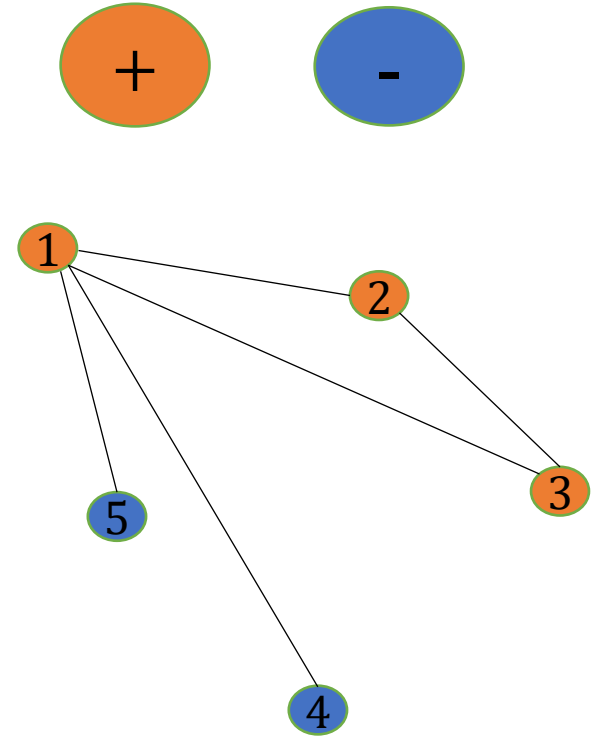
n vertices, each assigned spin $+1$ or -1 .

Distribution over configurations $\sigma \in \{-1, 1\}^n$

Density function:

$$\mu: \{-1, 1\}^n \rightarrow \mathbb{R}_{\geq 0}$$

Gibbs measure: $\mu(\sigma) = \exp(H(\sigma))$



Spin system

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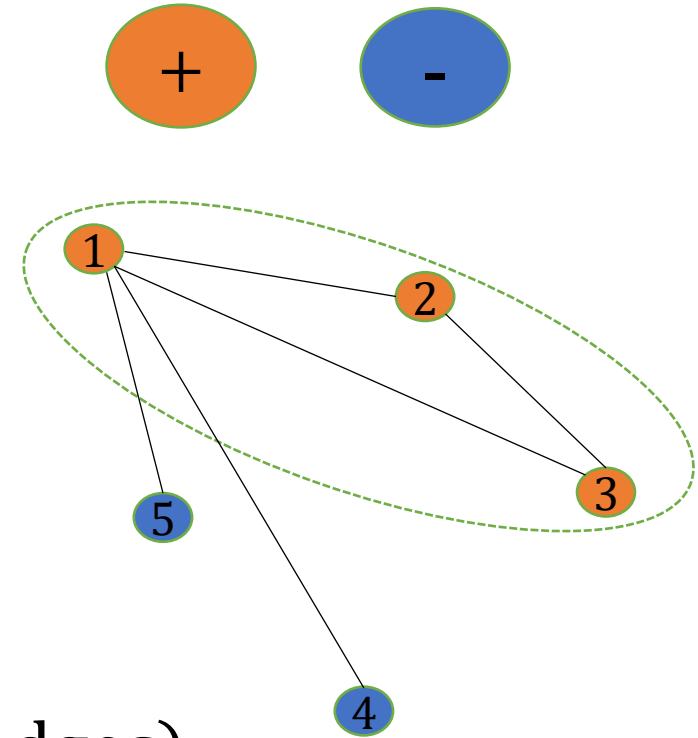
$$\mu: \{-1, 1\}^n \rightarrow \mathbb{R}_{\geq 0}$$

Ising model: pairwise interaction (graph edges)

$$H(\sigma) = \frac{\beta}{\sqrt{n}} \sum_{i \leq j} J_{ij} \sigma_i \sigma_j + \sum_i h_i \sigma_i$$

P-spin: interaction between p vertices (hypergraph edges)

$$H(\sigma) = \sum_{p, i_1, \dots, i_p} \frac{\beta_p}{n^{\frac{p-1}{2}}} J_{i_1, \dots, i_p} \sigma_1 \dots \sigma_p + \sum_i h_i \sigma_i$$



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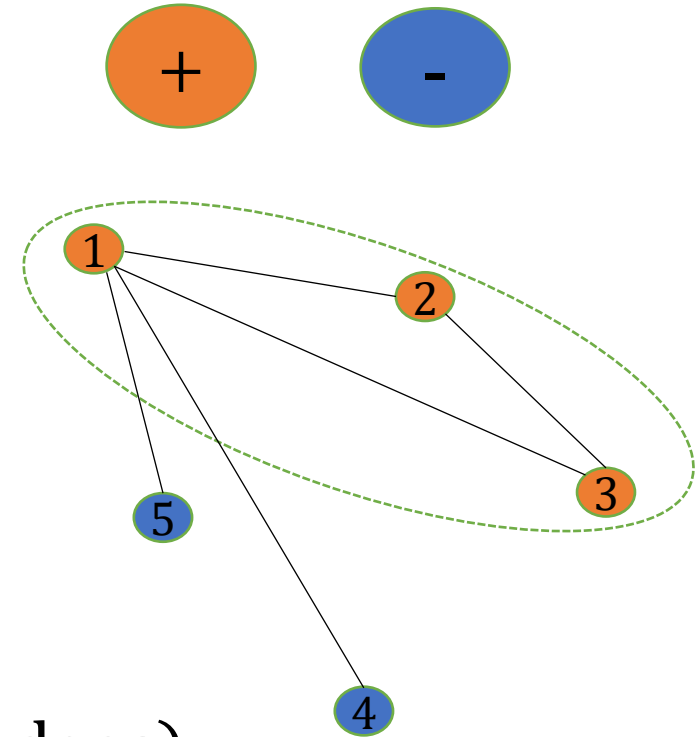
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Gaussians $\mathcal{N}(0,1)$



Glauber dynamics on p-spin

- Warmup, $p = 2$: if $0 \leq \beta \leq \epsilon$, optimal $O_\beta(n \log n)$ mixing of Glauber dynamics
- If $\sum_p \sqrt{p^3 \log p} \beta_p \leq \epsilon$. Let $\beta = \sum_p \sqrt{2^p p^3 \log p} \beta_p$, optimal $O_\beta(n \log n)$ mixing of Glauber dynamics [Anari-Jain-Koehler-Pham-V'23]
 - [Adhikari, Brennecke, Xu, Yau'22] spectral gap $\geq \Omega_\beta\left(\frac{1}{n}\right) \Rightarrow O_\beta(n^2 \log n)$ mixing

Thm [Anari-Jain-Koehler-Pham-V'23]: $\epsilon = \theta(1) \cdot \mu(\sigma) = \exp(H(\sigma))$.

If $\beta = \max_{\sigma \in \{\pm 1\}^n} \|\nabla^2 H(\sigma)\|_{op} \leq \epsilon$ then

μ^{hom} is $(1 + O(\beta))$ -entropically independent

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Conditionals of $\mu^{hom} \equiv$ pinnings of μ : $\mu(\cdot | \sigma_S = \tau)$

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Conditionals of $\mu^{hom} \equiv$ pinnings of μ : $\mu(\cdot | \sigma_S = \tau) = \exp(H_{S,\tau}(\sigma))$

$$\beta_{S,\tau} = \max_{\sigma_{S^c} \in \{\pm 1\}^n} \|\nabla^2 H_{S,\tau}(\sigma_{S^c})\|_{op} \leq \beta$$

$\mu^{hom}(\cdot | \sigma_S = \tau)$ is $(1 + O(\beta))$ -entropically independent

$\Rightarrow \tilde{O}(n^{1+O(\beta)})$ mixing time

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$$p = 2: H_{S,\tau}(\sigma) = \frac{\beta}{\sqrt{n}} \sum_{i,j \in S^c} J_{ij} \sigma_i \sigma_j + \sum_i h'_i \sigma_i$$

$$\beta_{S,\tau} \leq \beta \sqrt{\frac{|S^c|}{n}} \Rightarrow \tilde{O}(n) \text{ mixing time}$$

Conditionals of $\mu^{hom} \equiv$ pinnings of $\mu: \mu(\cdot | \sigma_S = \tau) = \exp(H_{S,\tau}(\sigma))$

$$\beta_{S,\tau} = \max_{\sigma_{S^c} \in \{\pm 1\}^n} \left\| \nabla^2 H_{S,\tau}(\sigma_{S^c}) \right\|_{op} \leq \beta$$

$$p = 3: H_{S,\tau}(\sigma) = \frac{\beta_3}{n} \left(\sum_{i,j,k \in S^c} J_{ijk} \sigma_i \sigma_j + \sum_{i,j \in S^c, k \in S} J_{ijk} \sigma_i \sigma_j \tau_k \right) + \sum_i h'_i \sigma_i$$

Hope: $\beta_{S,\tau} \leq \beta \sqrt{\frac{|S^c|}{n}}$. But, there are bad pinnings! E.g. $S^c = \{1,2\}$

$$H_{S,\tau}(\sigma) = \frac{\beta_3}{n} \sum_{k \in S} J_{12k} \tau_k \sigma_1 \sigma_2$$

Conditionals of $\mu^{hom} \equiv$ pinnings of $\mu: \mu(\cdot | \sigma_S = \tau) = \exp(H_{S,\tau}(\sigma))$

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Hope: $\beta_{S,\tau} \leq \beta \sqrt{\frac{|S^c|}{n}}$. But, there are bad pinnings! E.g. $S^c = \{1,2\}$

$$H_{S,(sign(J_{12k}))_k}(\sigma) = \frac{\beta_3}{n} \sum_{k \in S} |J_{12k}| \sigma_i \sigma_j \text{ thus } \beta_{S,\tau} \approx \beta_3 = \theta(1)$$

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But, “most” pinnings are good!

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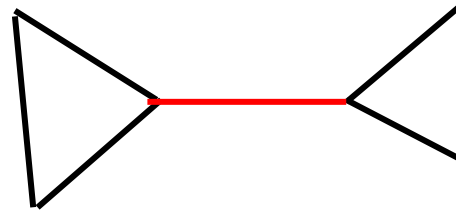
Sampling random spanning trees

Given G , output spanning tree T with probability $\frac{1}{\#spanning-trees}$

To find one spanning tree, need $\Omega(|E|)$ time.

\Rightarrow need $\Omega(|E|)$ time to sample.

[Anari-Liu-OveisGharan-Vinzant-Vuong'21] $O(|E|\log^2|E|)$ using up-down walk



Sampling random spanning trees

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Can we produce sample in sublinear time after preprocessing?

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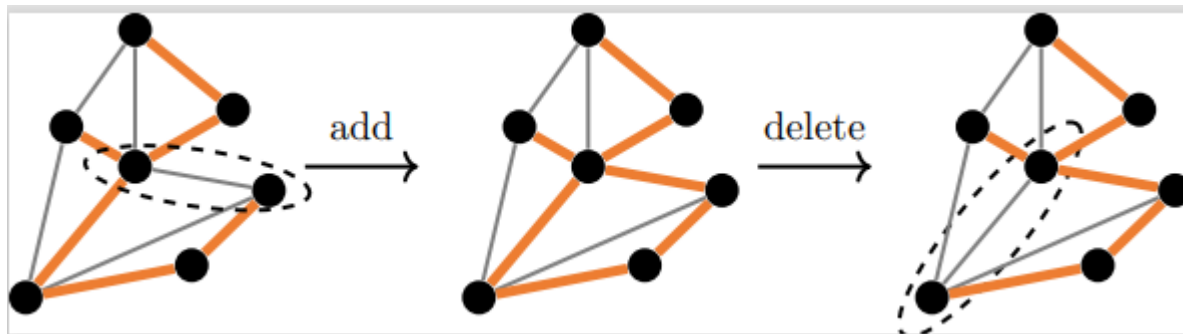
Can we produce sample in sublinear time after preprocessing?

[Anari-Liu-V--FOCS'22]: sample in $O(|V| \log^2 |V|)$ time after $O(|E| \log^2 |V|)$ preprocessing

Up-down walk

Repeat for sufficiently many times. Take tree T

1. Add an edge e
2. Remove an edge f uniformly at random from the unique circle in $T + e$



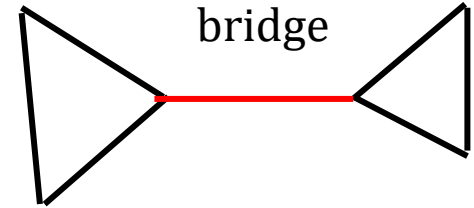
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Up-down walk \equiv down-up walk on the complement $\bar{\mu}: \binom{[n]}{n-k} \rightarrow \mathbb{R}_{\geq 0}$
defined by $\bar{\mu}([n] \setminus S) = \mu(S)$

Up-down walk



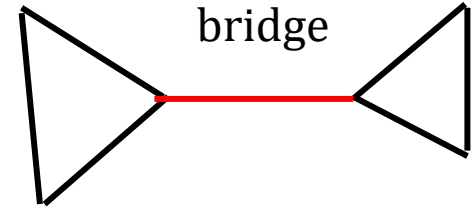
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Key points:

- Can implement 1 and 2 in $O(\log |V|)$ -time using link-cut tree
- If \exists bridge edge, need $\theta(|E| \log |E|)$ time to converge

Up-down walk



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Key points:

- Can implement 1 and 2 in $O(\log |V|)$ -time using link-cut tree
- If \exists bridge, need $\theta(|E| \log |E|)$ time to mix
- If all edges have same marginal, mixes in $O(|V| \log |V|)$ time

Isotropic transformation

Goal: make all edges/elements having the same marginal.

$$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$$

Let $p_e = Pr_{\mu}[e \in T]$. Replace edge e with $t_e = \lceil \frac{np_e}{k} \rceil$ parallel edges e' .



Strongly Rayleigh distributions

$D_{k \rightarrow 1}(S)$: sample $i \in S$ uniformly

μ is $\frac{1}{\alpha}$ -entropic independence $\Leftrightarrow \forall \nu$:

$$\mathcal{D}_{KL}(\nu || \mu) \geq \alpha k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$$

μ **strongly Rayleigh** \Rightarrow 1-entropic independence

$$\mathcal{D}_{KL}(\nu || \mu) \geq k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$$

Examples:

- $U(\{\text{spanning trees}\})$
- Determinantal point processes:

Improved entropic independence under uniform marginals

$\mu: \binom{[n]}{k} \rightarrow \mathbb{R}_{\geq 0}$ strongly Rayleigh. When $p_e \leq \tilde{O}\left(\frac{k}{n}\right) \forall e \in [n]$

$$\mathcal{D}_{KL}(\bar{\nu} || \bar{\mu}) \geq (n - k) \log(n/k) \mathcal{D}_{KL}(\nu D_{(n-k) \rightarrow 1} || \mu D_{(n-k) \rightarrow 1})$$

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1. $\mathcal{D}_{KL}(\bar{\nu} || \bar{\mu}) = \mathcal{D}_{KL}(\nu || \mu) \geq k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$

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1. $\mathcal{D}_{KL}(\bar{\nu} || \bar{\mu}) = \mathcal{D}_{KL}(\nu || \mu) \geq k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1})$

2. $k \mathcal{D}_{KL}(\nu D_{k \rightarrow 1} || \mu D_{k \rightarrow 1}) \geq (n - k) \log\left(\frac{n}{k}\right) \mathcal{D}_{KL}(\bar{\nu} D_{(n-k) \rightarrow 1} || \bar{\mu} D_{(n-k) \rightarrow 1})$

Here we use the uniform marginal assumption.

Improved EI implies improved mixing time

Entropy contraction of $D_{(n-k) \rightarrow 1}$ for $\bar{\mu}$ and its conditionals

\Rightarrow Entropy contraction of $D_{(n-k) \rightarrow (n-k-1)}$

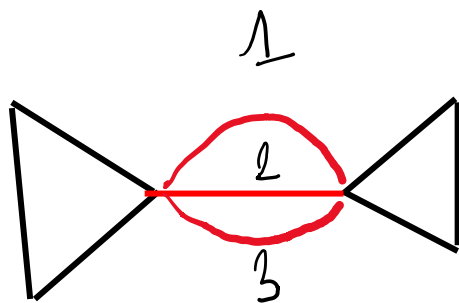
\Rightarrow Mixing time of up-down walk.

$(n - k)$ contraction $\Rightarrow n \log n$ mixing time ☹️

$(n - k) \log\left(\frac{n}{k}\right)$ contraction $\Rightarrow k \log n$ mixing time 😊

But, not all conditionals of $\bar{\mu}$ has improved entropy contraction ☹️
 i.e. exists \bar{S} s.t.

$$\mathcal{D}_{KL}(\bar{\nu}_{\bar{S}} || \bar{\mu}_{\bar{S}}) < (n - k) \log \left(\frac{n}{k} \right) \mathcal{D}_{KL}(\bar{\nu}_{\bar{S}} D_{(n-k) \rightarrow 1} || \bar{\mu}_{\bar{S}} D_{(n-k) \rightarrow 1})$$



$$\bar{S} = \{ \bar{1}, \bar{3} \}$$

Average local to global

For each set $\bar{W} \in \binom{[n]}{n-k-1}$ and s , if for “many” $\bar{S} \in \binom{\bar{W}}{n-s}$

$\bar{\mu}_{\bar{S}}$ has uniform marginal thus improved entropy contraction

then we still get $k \log n$ mixing time ☺

“many” = w/ prob. $1 - 1/n^{10}$ over uniformly chosen \bar{S}

Average local to global

For each set $\bar{W} \in \binom{[n]}{n-k-1}$ and s , if for “many” $\bar{S} \in \binom{\bar{W}}{n-s}$

$\bar{\mu}_{\bar{S}}$ has uniform marginal thus improved entropy contraction

Proof:

Compare marginals of $\bar{\mu}_{\bar{S}}$ and $\bar{\mu}_{\bar{S} \cup \{s'\}}$ for random s'

Since μ is strongly Rayleigh, marginal doesn't change much

Use martingale argument and Bernstein ineq.