# Reversible random walks on *dynamic* graphs

Nobutaka Shimizu (Tokyo Institute of Technology) \*Takeharu Shiraga

(Chuo University)

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#### Contents

# Nobutaka Shimizu, Takeharu Shiraga, **Reversible random walks on dynamic graphs** *Random Structures & Algorithms*, 2023

#### Main topic. Time-inhomogeneous Markov chain

Notation and Main result

Previous work

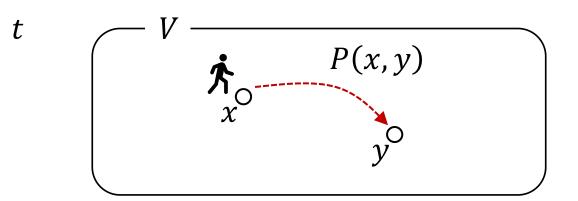
Idea of proof

Other topic

- $\checkmark$  V: Set of *n* vertices
- ✓  $P \in [0,1]^{V \times V}$ : Transition matrix on V

At each discrete time step t = 1, 2, ...,

the walker moves from x to y with probability P(x, y)



Referred to as the **random walk according to P** 

✓ i.e., time-homogeneous Markov chain

- $\checkmark$  V: Set of *n* vertices
- ✓  $P \in [0,1]^{V \times V}$ : Transition matrix on V

#### Random walk according to P.

A sequence of random variables  $X_0, X_1, X_2, \dots$  s.t.

$$\Pr\left(\begin{array}{c|c} X_{t} = x_{t} \\ X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t-1} = x_{t-1} \\ x_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t-1} \\ X_{t-1} = x_{t-1} \\ x_{t-1} \end{array}\right)$$
holds for all  $t = 1, 2, \dots$  and  $(x_{0}, \dots, x_{t}) \in V^{t+1}$ 
(where  $\Pr(X_{0} = x_{0}, \dots, X_{t-1} = x_{t-1}) > 0$ )

✓ i.e., time-homogeneous Markov chain

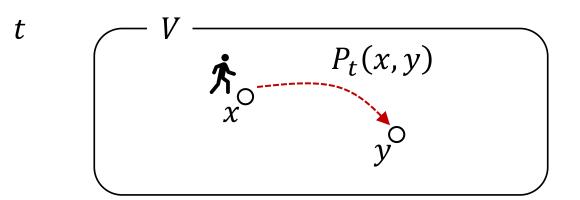
# Notation: Random walk according to $(P_t)_{t \ge 1}$

 $\checkmark$  V: Set of *n* vertices

✓  $(P_t)_{t \ge 1} = (P_1, P_2, ...)$ : <u>Sequence of transition matrices</u> on V

At each discrete time step t = 1, 2, ...,

the walker moves from x to y with probability  $P_t(x, y)$ 



Referred to as the **random walk according to**  $(P_t)_{t \ge 1}$ 

- ✓ i.e., time-**in**homogeneous Markov chain
  - > Transition matrix at time t is  $P_t$

# Notation: Random walk according to $(P_t)_{t \ge 1}$

- ✓ *V*: Set of *n* vertices
- ✓  $(P_t)_{t \ge 1} = (P_1, P_2, ...)$ : <u>Sequence of transition matrices</u> on *V*

#### Random walk according to $(P_t)_{t \ge 1}$ .

A sequence of random variables  $X_0, X_1, X_2, \dots$  s.t.

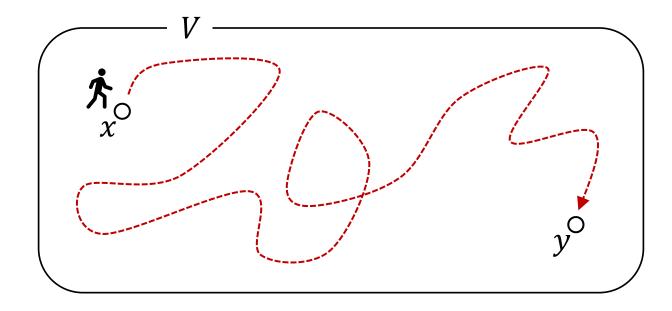
$$\Pr\left(\begin{array}{c|c} X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t} = x_{t} \\ X_{t-1} = x_{t-1} \end{array}\right) = \Pr\left(\begin{array}{c} X_{t-1} = x_{t-1} \\ P_{t}(x_{t-1}, x_{t}) \end{array}\right)$$
holds for all  $t = 1, 2, ...$  and  $(x_{0}, ..., x_{t}) \in V^{t+1}$ 
(where  $\Pr(X_{0} = x_{0}, ..., X_{t-1} = x_{t-1}) > 0$ )

- ✓ i.e., time-inhomogeneous Markov chain
  - > Transition matrix at time t is  $P_t$

# Hitting time

# **Hitting time** $t_{hit} \coloneqq \max_{x,y \in V} \mathbf{E}_x[\min\{t \ge 0 \mid X_t = y\}]$

✓ The expected # of steps for the walker to move from x to y (considering the worst-case pair of vertices x and y)



### Hitting time

# **Hitting time** $t_{hit} \coloneqq \max_{x,y \in V} \mathbf{E}_x[\min\{t \ge 0 \mid X_t = y\}]$

 ✓ The expected # of steps for the walker to move from x to y (considering the worst-case pair of vertices x and y)

Write  $t_{hit}(P)$  as the HT of the RW according to P

i.e., HT of a RW with a time-invariant transition matrix

✓ There is much previous work

Write  $t_{hit}((P_t)_{t\geq 1})$  as the HT of the RW according to  $(P_t)_{t\geq 1}$ 

i.e., HT of a RW with time-varying transition matrices

✓ Not much is known (This work)

### Main result (Hitting time)

✓ We give an upper bound on HT of a RW with <u>time-varying</u> transition matrices in terms of HTs of <u>time-invariant</u> ones:

**<u>Thm.1 (Hitting time)</u>**. Suppose  $(P_t)_{t \ge 1}$  satisfies the following:

- ✓ All  $P_1, P_2, ...$  are irreducible, reversible, and lazy
- ✓ All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$

Then, there is a constant C > 0 s.t.

$$\begin{split} t_{\rm hit} \big( (P_t)_{t \ge 1} \big) &\leq C \max_{t \ge 1} t_{\rm hit} (P_t) \\ &= C \max\{ t_{\rm hit} (P_1), t_{\rm hit} (P_2), t_{\rm hit} (P_3), \dots \}. \end{split}$$

 $t_{\text{hit}}((P_t)_{t\geq 1})$ : HT of the RW according to  $(P_t)_{t\geq 1}$ , i.e., HT of the RW with timevarying transition matrices ( $P_t$  at time t)  $t_{hit}(P_1)$ : HT of the RW according to  $P_1$ , i.e., HT of the RW with the time-invariant transition matrix ( $P_1$  at all times)

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Notation and Main result

**Previous work** 

Idea of proof

Other topic

**<u>Thm.1 (Hitting time)</u>**. Suppose  $(P_t)_{t \ge 1}$  satisfies the following:

- ✓ All  $P_1, P_2, ...$  are irreducible, reversible, and lazy
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 $\begin{aligned} t_{\rm hit} \big( (P_t)_{t \ge 1} \big) &\leq C \max_{t \ge 1} t_{\rm hit} (P_t) \\ &= C \max\{ t_{\rm hit} (P_1), t_{\rm hit} (P_2), t_{\rm hit} (P_3), \dots \}. \end{aligned}$ 

# Previous work: lazy simple RW on a (*static*) graph 11

✓ G: n-vertex graph  $\succ$  V(G): Vertex set of G  $\succ$  E(G): Edge set of G

→ deg(G, x): Degree of vertex  $x \in V(G)$ 

 $\frac{Lazy \ simple \ random \ walk \ on \ (static) \ G}{P(x,y)}:= \begin{cases} \frac{1}{2 \deg(G,x)} & (if \{x,y\} \in E(G)) \\ \frac{1}{2} & (if x = y) \\ 0 & (otherwise) \end{cases}$  At time t, the walker 1/6 1/2 1/6 1/2 1/6 1/2 1/6 1/2 1/6 1/2 1/6 1/6 1/2 1/6 1/

For any connected graph G, let P be the transition matrix ofthe lazy simple RW on G. Then,There exists a tight

$$t_{\rm hit}(P) = O(n^3).$$
 exam

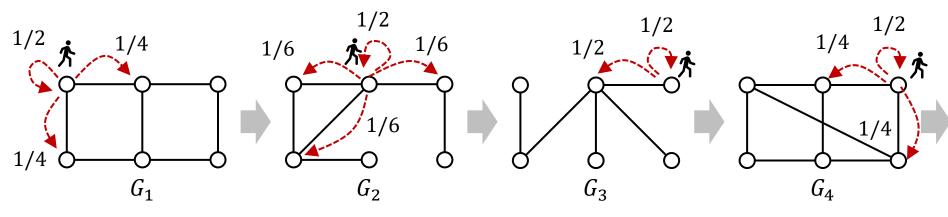
example (Lollipop graph)

[Aleliunas, Karp, Lipton, Lovász, Rackoff. FOCS 79]

# Previous work: lazy simple RW on a *dynamic* graph 12

 $\checkmark$  G<sub>1</sub>, G<sub>2</sub>, ... : Sequence of edge-changing *n*-vertex graphs

 $\frac{Lazy \ simple \ random \ walk \ on \ G_1, G_2, \dots}{\text{moves from } x \ to \ y \ w.p. \ P_t(x, y) \ defined \ as \ follows:}$   $P_t(x, y) \coloneqq \begin{cases} \frac{1}{2 \deg(G_t, x)} & (\text{if } \{x, y\} \in E(G_t)) \\ \frac{1}{2} & (\text{if } x = y) \\ 0 & (\text{otherwise}) \end{cases}$ 

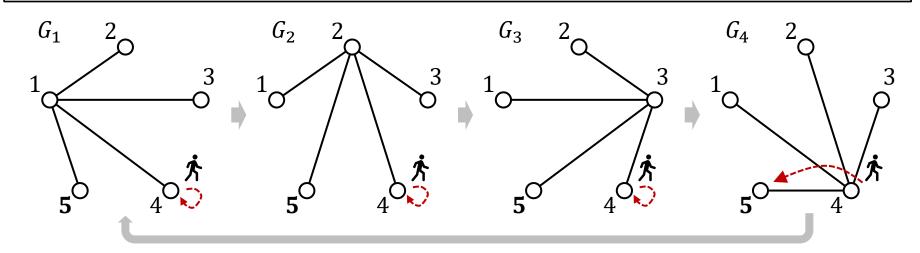


Ex. Lazy simple random walk on  $G_1, G_2, ...$ 

#### Previous work: Exponential lower bound for LSRW 13

<u>Sisyphus wheel</u>. Sequence of star graphs  $G_1, G_2, ...$  with  $V(G_t) = \{1, ..., n\}$ , where the center changes periodically in 1, ..., n - 1

For Sisyphus wheel  $G_1, G_2, ..., \text{ let } P_t$  be the transition matrix of the <u>lazy simple random walk</u> on  $G_t$ . Then,  $t_{\text{hit}}((P_t)_{t\geq 1}) = 2^{\Omega(n)}$ . [Avin, Kouský, Lotler. ICALP 08, RS&A 18]



Ex of n = 5.

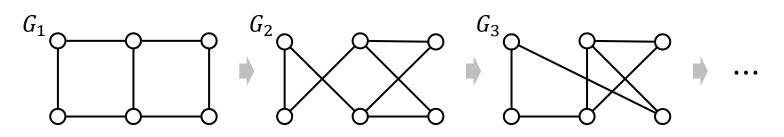
The walker must stay n - 2 consecutive steps to reach the vertex n

# Previous work: Upper bound for lazy simple RW 14

∀sequence of connected graphs  $G_1, G_2, ...$  with an <u>invariant</u> <u>degree distribution</u>, let  $P_t$  be the transition matrix of lazy simple walk on  $G_t$ . Then,

$$t_{\rm hit}\big((P_t)_{t\geq 1}\big)=O\big(n^3\log n\big).$$

[Sauerwald, Zanetti. ICALP 19]



- ✓ In general, there exists a sequence of graphs s.t. HT is exponential (Sisyphus wheel)
- $\checkmark$  If the degree distribution is invariant, HT is polynomial

We can apply Thm.1 for <u>lazy simple RW</u> on  $G_1, G_2, ...$  with <u>time-invariant degree distribution</u>!

- ✓  $P_t$  (Transition matrix of LSRW on  $G_t$ ) is irreducible, reversible and lazy (if  $G_t$  is connected)
- ✓ Stationary distribution of  $P_t$  is  $\frac{\deg(G_t, x)}{2|E(G_t)|}$ 
  - Stationary distribution is invariant if degree dist. is !

**<u>Thm.1 (Hitting time)</u>**. Suppose  $(P_t)_{t \ge 1}$  satisfies the following:

- ✓ All  $P_1, P_2, ...$  are irreducible, reversible, and lazy
- ✓ All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$

Then, there is a constant C > 0 s.t.  $t_{\text{hit}}((P_t)_{t \ge 1}) \le C \max_{t \ge 1} t_{\text{hit}}(P_t)$ =  $C \max\{t_{\text{hit}}(P_1), t_{\text{hit}}(P_2), t_{\text{hit}}(P_3), ...\}$ 

# Application of Theorem 1: Lazy simple RW

**<u>Corollary of Thm.1</u>**.  $\forall$ sequence of connected graphs  $(G_t)_{t\geq 1}$ with an <u>invariant degree distribution</u>, let  $P_t$  be the transition matrix of the <u>lazy simple RW</u> on  $G_t$ . Then,  $t_{\text{hit}}((P_t)_{t\geq 1}) = O(\mathbf{n}^3)$ .

✓ Improves  $O(n^3 \log n)$  bound of the previous work!

[Sauerwald, Zanetti. ICALP 19]

#### Remark.

HT of LSRW on  $G_1$ 

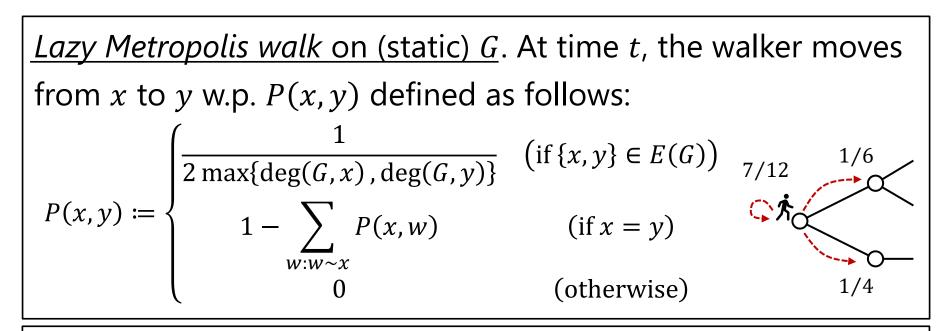
✓ 
$$t_{\text{hit}}(P_1) = O(n^3), t_{\text{hit}}(P_2) = O(n^3), ...$$

[Aleliunas et al. 79] (previous work on static graphs)

<u>Thm.1.</u> Suppose all  $P_1, P_2, ...$  are irreducible, reversible, lazy, and have the same stationary distribution  $\pi$ . Then,  $t_{\text{hit}}((P_t)_{t\geq 1}) \leq C \max\{t_{\text{hit}}(P_1), t_{\text{hit}}(P_2), ...\}$ .

HT of LSRW on  $G_2$ 

#### Other example: lazy Metropolis walk



For any <u>connected</u> graph *G*, let *P* be the transition matrix of the <u>lazy Metropolis walk</u> on *G*. Then,  $t_{hit}(P) = O(n^2)$ . [Nonaka, Ono, Sadakane, Yamashita. Theoretical Compt. Sci. 10]

- ✓ Using local degree information achieves  $O(n^2)$  hitting time
- ✓ There are no previous studies about dynamic cases

# Application of Theorem 1: Lazy Metropolis walk 18

**<u>Cor. of Thm.1</u>**.  $\forall$  sequence of connected graphs  $(G_t)_{t\geq 1}$ , let  $P_t$  be the transition matrix of the <u>lazy Metropolis W</u> on  $G_t$ . Then,  $t_{\text{hit}}((P_t)_{t\geq 1}) = O(n^2)$ .

**Remark.** Same bound as the static graph! i.e., LMW is **robust** for edge-changes

- ✓ Stationary distribution of LMW is the uniform distribution  $(\pi(x) = 1/n \text{ for any graph})$ 
  - Stationary distribution is invariant for any graphs!

✓ 
$$t_{\text{hit}}(P_1) = O(n^2), t_{\text{hit}}(P_2) = O(n^2), ...$$

HT of LMW on  $G_1$  HT of LMW on  $G_2$ 

[Nonaka et al. 10] (previous work on static graphs)

<u>Thm.1</u>. Suppose all  $P_1, P_2, ...$  are irreducible, reversible, lazy, and have the same stationary distribution  $\pi$ . Then,  $t_{\text{hit}}((P_t)_{t\geq 1}) \leq C \max\{t_{\text{hit}}(P_1), t_{\text{hit}}(P_2), ...\}$ .

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Notation and Main result

**Previous works** 

Idea of proof

Other topic

 $\checkmark \tau_y \coloneqq \min\{t \ge 0 \mid X_t = y\}$ : **Hitting time to y** (random variable)

**Remark.** Hitting time  $t_{hit} = \max_{x,y \in V} \mathbf{E}_x[\tau_y]$  from definition

**<u>Hitting time lemma</u>.** Suppose  $(P_t)_{t\geq 1}$  satisfies the following:  $\checkmark$  All  $P_1, P_2, ...$  are irreducible and reversible  $\checkmark$  All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$ Then, for any  $w \in V$  and  $T \geq 0$ ,

$$\Pr_{\pi}(\tau_w > T) \le \left(1 - \frac{1}{\max_{t \ge 1} t_{\text{hit}}(P_t)}\right)^T$$

✓ For the walker <u>starting from the stationary distribution</u>, the hitting time to a vertex decreases exponentially

$$\checkmark \tau_y \coloneqq \min\{t \ge 0 \mid X_t = y\}$$
: **Hitting time to y** (random variable)

**<u>Hitting time lemma.</u>** Suppose all  $P_1, P_2, ...$  are irreducible, reversible, and have the same stationary distribution  $\pi$ . Then, for any  $w \in V$  and  $T \ge 0$ ,  $\Pr_{\pi}(\tau_w > T) \le \left(1 - \frac{1}{\max_{t>1} t_{\text{hit}}(P_t)}\right)^T.$ 

➢ HTL implies that "Hitting time from stationary E<sub>π</sub>[τ<sub>w</sub>]" is bounded by max t<sub>hit</sub>(P<sub>t</sub>):  $E_{\pi}[τ_w] \leq \sum_{T=0}^{\infty} \left(1 - \frac{1}{\max_{t \geq 1} t_{hit}(P_t)}\right)^T = \max_{t \geq 1} t_{hit}(P_t).$ 

# Proof of hitting time lemma (1/3)

✓  $D_w \in \{0,1\}^{V \times V}$ : diagonal matrix where  $D_w(x,x) = \mathbf{1}_{x \neq w}$ ➤ **Key observation**.  $\tau_w$  can be expressed in terms of  $D_w$ :

$$\Pr_{x}(\tau_{w} > T, X_{T} = y) = \Pr_{x}\left(\bigwedge_{t=0}^{T} \{X_{t} \neq w\}, X_{T} = y\right)$$

$$= \left(\prod_{t=1}^{T} D_{w}P_{t}D_{w}\right)(x, y).$$
"Transitions that exclude reaching w"
$$D_{w} = \left(\bigvee_{t=1}^{W} 1 \bigcup_{t=1}^{W} w D_{w}PD_{w} = \left(\bigcup_{t=1}^{W} 0 \bigcup_{t=1}^{W} w D_{w}PD_{w}\right)(x, y)\right)$$
Identity matrix
$$P \text{ except that}$$
its w-th row and column are set to 0

#### Proof of Hitting time lemma (2/3)

**Courant-Fischer-Weyl Min-max theorem** 

✓  $\rho(A)$ : the spectral radius of A

# Proof of Hitting time lemma (3/3)

✓ The following lemma, a basic consequence of the *Perron-Frobenius theorem*, concludes the proof:

**Lem.** Suppose *P* is irreducible & reversible. Then,  $\forall w \in V$ , the spectral radius  $\rho(D_w P D_w)$  of  $D_w P D_w$  is bounded by  $1 - \frac{1}{t_{hit}(P)}$ . [Aldous, Fill 02]

$$\Pr_{\pi}(\tau_{w} > T) \le \prod_{t=1}^{T} \rho(D_{w}P_{t}D_{w}) \le \prod_{t=1}^{T} \left(1 - \frac{1}{t_{\text{hit}}(P_{t})}\right).$$

Hitting time lemma. Suppose all P<sub>1</sub>, P<sub>2</sub>, ... are irreducible, and have the same

stationary distribution 
$$\pi$$
. Then,  $\Pr_{\pi}(\tau_w > T) \le \left(1 - \frac{1}{\max_{t \ge 1} t_{\text{hit}}(P_t)}\right)^T$ 

Proof overview:

Time taken for a walker

 $t_{\rm hit}((P_t)_{t\geq 1}) \leq$ 

to converge  $\pi$  (from the worst initial pos.)

 $+ \max_{t \ge 1} t_{\text{hit}}(P_t)$ 

Hitting time lemma:

 $\mathbf{E}_{\pi}(\tau_w) \leq \max_{t>1} t_{\mathrm{hit}}(P_t) \,.$ 

✓ Mixing time bound for time-inhomogeneous Markov chain

✓ For time-homogeneous MC, the following is well-known:

<u>**Thm.</u>** Suppose *P* is irreducible, reversible and lazy. Then,  $t_{\text{mix}}^{(\infty)}(P) \leq Ct_{\text{hit}}(P).$ [Levin, Peres, Wilmer. 08]</u>

#### Key tool 2: Mixing time $\leq$ Hitting time

**<u>Thm.2 (Mixing time)</u>**. Suppose  $(P_t)_{t \ge 1}$  satisfies the following:

- ✓ All  $P_1, P_2, ...$  are irreducible, reversible and lazy
- ✓ All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$

Then, there is a constant C > 0 s.t.

$$t_{\min}^{(\infty)}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\operatorname{hit}}(P_t).$$

$$(\ell_{\infty}(\pi))$$
 -) Mixing time  $t_{\min}^{(\infty)}((P_t)_{t\geq 1})$ .

$$\begin{split} t_{\min}^{(\infty)} \big( (P_t)_{t \ge 1} \big) \\ &\coloneqq \min \left\{ t \ge 0 : \max_{s \ge 0, x, y \in V} \left| \frac{(P_{s+1} P_{s+2} \cdots P_{s+t})(x, y)}{\pi(y)} - 1 \right| \le \frac{1}{2} \right\} \end{split}$$

#### Remark.

# ✓ The following bound of $t_{\min}^{(\infty)}((P_t)_{t \ge 1})$ has been shown:

Suppose all  $P_1, P_2, ...$  are irreducible, aperiodic, reversible, and have the same stationary distribution  $\pi$ . Then,  $\exists$  constant C s.t.

$$t_{\min}^{(\infty)}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} \left(\frac{\log \pi_{\min}^{-1}}{1-\lambda_{\star}(P_t)}\right).$$

[Saloff-Coste, Zúñiga. Stochastic Processes and their Applications 07]

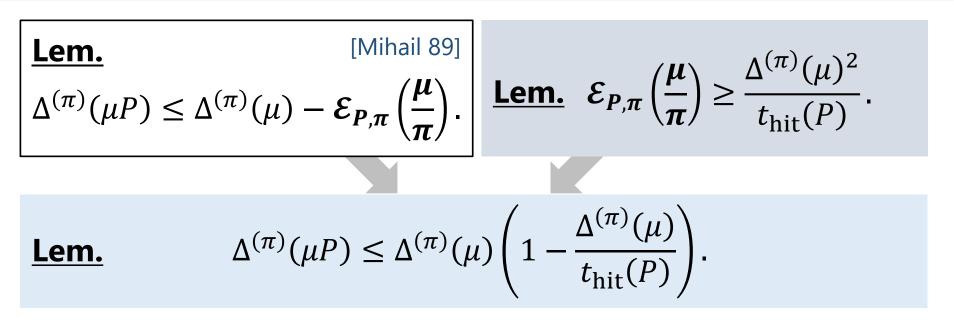
 $\lambda_{\star}(P)$ : 2<sup>nd</sup> largest eigenvalue in absolute value of P

- ✓ For some cases (e.g., LSRW on expander graphs), this gives a better bound than our bound of  $\max_{t \ge 1} t_{hit}(P_t)$
- ✓ However, **there exists bad examples** (e.g, LSRW on cycles) where this bound gets  $\max_{t \ge 1} t_{hit}(P_t) \log n$

#### Proof outline for Theorem 2 (Mixing time)

For a probability vector  $\mu \in [0,1]^V$ , let  $\Delta^{(\pi)}(\mu) \coloneqq \left\| \frac{\mu}{\pi} - \mathbf{1} \right\|_{2,\pi}^2 = \sum_{x,y} \pi(x) \left( \frac{\mu(x)}{\pi(x)} - 1 \right)^2.$  $\Delta^{(\pi)}(\mu P) \leq \Delta^{(\pi)}(\mu) \left(1 - \frac{\Delta^{(\pi)}(\mu)}{t_{\text{hit}}(P)}\right).$ Lem.  $\mu P$  gets closer to  $\pi$  than  $\mu$  in terms of  $t_{\rm hit}(P)$ (dist. after one step) Applying repeatedly For  $T \ge C \max_{t\ge 1} t_{\text{hit}}(P_t)$ ,  $\Delta^{(\pi)}(\mu P_1 \cdots P_T) \le \frac{1}{2}$ . Lem.  $\ell_2(\pi)$ -norm  $\rightarrow \ell_{\infty}(\pi)$ -norm **<u>Thm.2.</u>** For  $T \ge C \max_{t \ge 1} t_{\text{hit}}(P_t)$ ,  $\max_{s \ge 0, x, y \in V} \left| \frac{(P_{s+1} \cdots P_{s+T})(x, y)}{\pi(y)} - 1 \right| \le \frac{1}{2}$ .

#### Key technical lemma



For  $f \in \mathbb{R}^{V}$ , let  $\mathcal{E}_{P,\pi}(f) \coloneqq \mathcal{E}_{P,\pi}(f,f) = \frac{1}{2} \sum_{x,y \in V} \pi(x)P(x,y)(f(x) - f(y))^{2}$  $= \langle f, f \rangle_{\pi} - \langle f, Pf \rangle_{\pi}.$  (Dirichlet form)

 $\langle f, g \rangle_{\pi} \coloneqq \sum_{x \in V} \pi(x) f(x) g(x)$ :  $\pi$ -inner product

#### Key technical lemma (Proof, 1/3)

**Lem.** For any irreducible & reversible *P*,  $\mathcal{E}_{P,\pi}\left(\frac{\mu}{\pi}\right) \geq \frac{\Delta^{(\pi)}(\mu)^2}{t_{\text{hit}}(P)}.$ 

$$\checkmark \text{ Let } g(x) \coloneqq \left\| \frac{\mu}{\pi} \right\|_{\infty} - \frac{\mu(x)}{\pi(x)}. \qquad \succ \left\| \frac{\mu}{\pi} \right\|_{\infty} = \max_{x \in V} \frac{\mu(x)}{\pi(x)}$$

✓ The proof consists of the following three statements:

**1.** 
$$\mathcal{E}_{P,\pi}\left(\frac{\mu}{\pi}\right) = \langle g,g \rangle_{\pi} - \langle Pg,g \rangle_{\pi}$$
  
**2.**  $\langle Pg,g \rangle_{\pi} \leq \left(1 - \frac{1}{t_{\mathrm{hit}}(P)}\right) \langle g,g \rangle_{\pi}.$ 

**3.**  $\langle g, g \rangle_{\pi} \geq \Delta^{(2,\pi)}(\mu)^2$ 

#### Key technical lemma (Proof, 2/3)

$$\checkmark g(x) \coloneqq \left\|\frac{\mu}{\pi}\right\|_{\infty} - \frac{\mu(x)}{\pi(x)}.$$

$$I. \text{ Since } \frac{\mu(x)}{\pi(x)} - \frac{\mu(y)}{\pi(y)} = g(y) - g(x),$$

$$\varepsilon_{P,\pi}\left(\frac{\mu}{\pi}\right) = \varepsilon_{P,\pi}(g) = \langle g, g \rangle_{\pi} - \langle Pg, g \rangle_{\pi}.$$

**2.** Let  $w \in V$  be a vertex s.t.  $\frac{\mu(w)}{\pi(w)} = \left\|\frac{\mu}{\pi}\right\|_{\infty}$ . Then, g(w) = 0 and  $\langle Pg, g \rangle_{\pi} = \langle D_w P D_w g, g \rangle_{\pi} \le \rho (D_w P D_w) \langle g, g \rangle_{\pi}$  $\le \left(1 - \frac{1}{t_{\text{hit}}(P)}\right) \langle g, g \rangle_{\pi}.$ 

 $D_w$ : Identity matrix except that its (w, w)-entry is 0

$$\underline{\text{Lem.}} \rho(D_w P D_w) \le 1 - \frac{1}{t_{\text{hit}}(P)}.$$

#### Key technical lemma (Proof, 3/3)

$$\checkmark g(x) \coloneqq \left\|\frac{\mu}{\pi}\right\|_{\infty} - \frac{\mu(x)}{\pi(x)}.$$

3. From

$$\begin{split} \Delta^{(\pi)}(\mu) &= \sum_{x \in V} \pi(x) \left( \frac{\mu(x)}{\pi(x)} - 1 \right)^2 = \sum_{x \in V} \pi(x) \left( \frac{\mu(x)}{\pi(x)} \right)^2 - 1 \le \left\| \frac{\mu}{\pi} \right\|_{\infty} - 1, \\ \langle g, g \rangle_{\pi} &= \sum_{x \in V} \pi(x) \left( \left\| \frac{\mu}{\pi} \right\|_{\infty} - \frac{\mu(x)}{\pi(x)} \right)^2 \\ &= \left\| \frac{\mu}{\pi} \right\|_{\infty}^2 + \sum_{x \in V} \pi(x) \left( \frac{\mu(x)}{\pi(x)} \right)^2 - 2 \left\| \frac{\mu}{\pi} \right\|_{\infty} \\ &\ge \left( \left\| \frac{\mu}{\pi} \right\|_{\infty} - 1 \right)^2 \ge \Delta^{(\pi)}(\mu)^2. \qquad \sum_{x \in V} \pi(x) \left( \frac{\mu(x)}{\pi(x)} \right)^2 \ge 1 \end{split}$$

#### Contents

Notation and Main result

Previous work

Idea of proof

#### **Other topic**

✓ We also studied other parameters of a random walk according to  $(P_t)_{t \ge 1}$ :

#### Cover time

- $\succ$  Hitting and cover times of k-independent walkers
- Coalescing time

#### Cover time

**Cover time** 
$$t_{cov} := \max_{x \in V} \mathbf{E}_x [\min\{t \ge 0 \mid \{X_0, X_1, \dots, X_t\} = V\}]$$

 ✓ The expected # of steps for the walker to <u>visit all vertices</u> (from the worst initial position)

Write  $t_{cov}(P)$  as the CT of the RW according to P

i.e., CT of a RW with a time-invariant transition matrix

✓ There is much previous work

Write  $t_{cov}((P_t)_{t\geq 1})$  as the CT of the RW according to  $(P_t)_{t\geq 1}$ 

i.e., CT of a RW with <u>time-varying</u> transition matrices
✓ Not much is known

 $\checkmark$  There is much previous work for time-invariant P, e.g.,

For any <u>connected</u> graph G, let P be the transition matrix of the lazy simple RW on G. Then,

 $t_{\rm cov}(P) = O(n^3).$ 

There exists a tight

example (Lollipop graph)

[Aleliunas, Karp, Lipton, Lovász, Rackoff. FOCS 79]

For any <u>connected</u> graph G, let P be the transition matrix of the lazy Metropolis walk on G. Then, There exists a tight

$$t_{\rm cov}(P) = O(n^2 \log n).$$

example (glitter star)

[Nonaka, Ono, Sadakane, Yamashita. Theoretical Compt. Sci. 10]

For any irreducible P,  $t_{cov}(P) \leq t_{hit}(P) \log n$ .

[Matthews. Annals of Proability 88]

#### **Result (Cover time)**

**Thm.3 (Cover time).** Suppose  $(P_t)_{t \ge 1}$  satisfies the following:  $\checkmark$  All  $P_1, P_2, ...$  are irreducible, reversible, and lazy  $\checkmark$  All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$ Then, there is a constant C > 0 s.t.

 $t_{\rm cov}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\rm hit}(P_t) \log n.$ 

- ✓ Multiplying  $\max_{t \ge 1} t_{hit}(P_t)$  by  $O(\log n)$  is sufficient to cover all vertices (even for the time-inhomogeneous case)
- ✓ Theorem 3 gives tight bounds for some cases
  - Lazy Metropolis walk

# Application of Theorem 3: Lazy Metropolis walk 37

**<u>Cor. of Thm.3</u>**.  $\forall$ sequence of connected graphs  $(G_t)_{t \ge 1}$ , let  $P_t$  be the transition matrix of the <u>lazy Metropolis W</u> on  $G_t$ . Then,  $t_{cov}((P_t)_{t \ge 1}) = O(n^2 \log n).$ 

**Remark.** Same bound as the static graph! i.e., LMW is **robust** for edge-changes

- ✓ Stationary distribution of LMW is the uniform distribution  $(\pi(x) = 1/n \text{ for any graph})$ 
  - Stationary distribution is invariant for any graphs!

✓ 
$$t_{\text{hit}}(P_1) = O(n^2), t_{\text{hit}}(P_2) = O(n^2), ...$$

HT of LMW on  $G_1$  HT of LMW on  $G_2$ 

[Nonaka et al. 10] (previous work on static graphs)

**<u>Thm.3</u>**. Suppose all  $P_1, P_2, ...$  are irreducible, reversible, lazy, and have the same stationary distribution  $\pi$ . Then,  $t_{cov}((P_t)_{t\geq 1}) \leq C \max_{t>1} t_{hit}(P_t) \log n$ .

We can assume the initial position  $\sim \pi$  from Thm. 2

$$\checkmark \tau_{\rm cov} \coloneqq \min\{t \ge 0 \mid \{X_0, \dots, X_t\} = V\}$$

$$\checkmark \tau_y \coloneqq \min\{t \ge 0 \mid X_t = y\}$$

$$\frac{\text{Thm 2 (Mixing time).}}{t_{\min}^{(\infty)}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\text{hit}}(P_t).}$$

For  $T = \max_{t \ge 1} t_{hit}(P_t) \log n$ , Union bound + HTL implies  $\Pr_{\pi}(\tau_{cov} > Ti) = \Pr_{\pi}\left(\bigcup_{w \in V} \{\tau_w > Ti\}\right) \le n\left(1 - \frac{1}{\max_{t \ge 1} t_{hit}(P_t)}\right)^{Ti}$   $\le n^{-(i-1)}.$ Hitting time lemma.

$$\Pr_{\pi}(\tau_w > T) \le \left(1 - \frac{1}{\max_{t \ge 1} t_{\text{hit}}(P_t)}\right)^T.$$
  
Hence,  $\mathbf{E}_{\pi}[\tau_{\text{cov}}] = O(T).$ 

**Corollary of Thm.3.**  $\forall$ sequence of connected graphs  $(G_t)_{t\geq 1}$ with an invariant degree distribution, let  $P_t$  be the transition matrix of the lazy simple RW on  $G_t$ . Then,  $t_{cov}((P_t)_{t\geq 1}) = O(n^3 \log n).$ 

Q. Is it tight?

Time-invariant case:

Is there a bad sequence of graphs with an invariant degree dist. s.t. t<sub>cov</sub>((P<sub>t</sub>)<sub>t≥1</sub>) = Ω(n<sup>3</sup> log n)?

> There exists a tight example (Lollipop graph)

For any <u>connected</u> graph G, let P be the transition matrix of the <u>lazy simple</u> <u>RW</u> on G. Then,  $t_{cov}(P) = O(n^3)$ . [Aleliunas, Karp, Lipton, Lovász, Rackoff. FOCS 79]

# Open problem

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**<u>Thm.3.</u>** Suppose all  $P_1, P_2, ...$  are irreducible, reversible, lazy, and have the same stationary distribution  $\pi$ . Then,

 $t_{\text{cov}}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\text{hit}}(P_t) \log n$ .

**<u>Conjecture</u>**. Suppose all  $P_1, P_2, ...$  are irreducible, reversible, lazy, and have the same stationary distribution  $\pi$ . Then,  $t_{cov}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{cov}(P_t) \ (?)$ 

✓ Is it true? Or a counter-example exists?

> Do good tools like the Hitting time lemma exist?

e.g.,  $\Pr_{\pi}(\tau_{\rm cov} > T) \leq \cdots$ 

#### Conclusion

✓ We give an upper bound on HT of a RW with <u>time-varying</u> transition matrices, in terms of HTs of <u>time-invariant</u> ones:

**<u>Thm.1 (Hitting time)</u>**. Suppose  $(P_t)_{t \ge 1}$  satisfies the following:

- ✓ All  $P_1$ ,  $P_2$ , ... are irreducible, reversible, and lazy
- ✓ All  $P_1, P_2, ...$  have the same stationary distribution  $\pi$

Then, there is a constant C > 0 s.t.

$$t_{\text{hit}}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\text{hit}}(P_t) = C \max\{t_{\text{hit}}(P_1), t_{\text{hit}}(P_2), t_{\text{hit}}(P_3), \dots\}.$$

✓ We also give upper bounds for the mixing and cover times in terms of  $\max_{t \ge 1} t_{\text{hit}}(P_t)$ 

# Summary: Thank you for your attention!

#### Lazy simple RW

Graph	<i>t</i> <sub>hit</sub>	t <sub>cov</sub>	
$\forall$ connected G (static)	$O(n^3)$		[Aleliunas et al. 79]
$\exists$ connected $G$ (static)	$\Omegaig(n^3ig)$ (Lollipop)		[Feige. 95]
$\exists$ seq. of connected graphs $G_1, G_2,$	$2^{\Omega(n)}$ (Sisyphus wheel)		[Avin, Kouský, Lotler. 08]
$\forall$ seq. of connected graphs $G_1, G_2,$ with <u>time-invariant degree dist</u> .	$O(n^3 \log n)$	$O(n^3 \log^2 n)$	[Sauerwald, Zanneti. 19]
	$O(n^{3})$	$O(n^3 \log n)$	[Shimizu, S. 23]

#### Lazy Metropolis walk

Graph	t <sub>hit</sub>	t <sub>cov</sub>	
$\forall$ connected G (static)	$O(n^2)$	$O(n^2 \log n)$	[Nonaka et al. 10]
$\exists$ connected $G$ (static)	$\Omegaig(n^2ig)$ (e.g., path)	$\Omegaig(n^2\log nig)$ (glitter star)	
$\forall$ seq. of connected graphs $G_1, G_2,$	$O(n^2)$	$O(n^2 \log n)$	[Shimizu, S. 23]
$t_{\text{hit}}((\boldsymbol{P}_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\text{hit}}(\boldsymbol{P}_t)$	$t_{\rm cov}((P_t)_{t\geq 1}) \leq C \max_{t\geq 1} t_{\rm hit}(P_t) \log n.$		