

# Improved bounds for the zeros of the chromatic polynomial

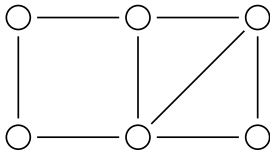
Viresh Patel

Queen Mary, University of London

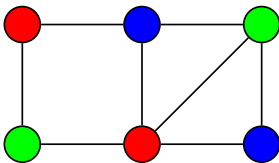
MCMC 2.0 - Shonan Seminar 186

Joint work with  
Matthew Jenssen (KCL) and Guus Regts (Amsterdam)

## Graph colouring



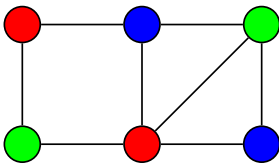
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**$k$  colouring of  $G$**

$f : V(G) \rightarrow \{1, \dots, k\}$  s.t.  $f(u) \neq f(v)$  whenever  $uv \in E(G)$

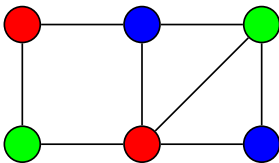
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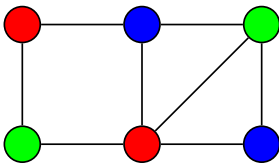
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$C_G(k) = \#$   $k$ -colourings of  $G$

(Birkhoff 1912)

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- Originally introduced to approach four colour problem
- Examples
  - $C_{K_r}(k) = k(k-1)(k-2)\cdots(k-r+1)$
  - $C_{n\text{-vertex tree}}(k) = k(k-1)^{n-1}$

# Computational counting

Want FPTAS / FPRAS for

$$C_G(k) \text{ for } k \in \mathbb{N}$$

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- exactly  $\forall z \in \mathbb{C} \setminus \{0, 1, 2\}$ . Jaeger-Vertigan-Welsh 1990
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**Conjecture**  $\exists$  FPTAS for  $C_G(k)$  provided  $k > \Delta(G)$

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For  $k \in \mathbb{N}$

- FPRAS for  $k \geq 2\Delta$  Jerrum 1994
- FPRAS for  $k > (\frac{11}{6} - \epsilon)\Delta$  Vigoda 2006, CDMPP 2019
- FPTAS for  $k \geq 2\Delta$  Liu-Sinclair-Srivastava 2019

**Question:** (Brenti, Royle, Wagner 1994)

Is there a function  $f(k)$  such that  $C_G(z) \neq 0$  whenever  $\Delta(G) \leq k$  and  $|z| \geq f(k)$ ?

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**Conjecture:**  $C_G(z) \neq 0$  if  $\Re(z) > \Delta(G)$

Sokal 2003

Implies

**Conjecture:**  $\exists$  FPTAS for  $C_G(k)$  provided  $k > \Delta(G)$

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(new short proof)

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Let  $G = (V, E)$  with  $\Delta(G) \leq k$  and  $|V| = n$

$$C_G(z) = a_0 z^n + a_1 z^{n-1} + \cdots + a_{n-1} z + a_n$$

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Theorem (Whitney 1932)

$a_i =$  number of *broken-circuit free sets* of size  $i$  in  $G$

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### Broken circuit free sets (BCF sets)

Fix an ordering of  $E$ . We say  $A \subseteq E$  is broken circuit free if

- $A$  is a forest, and
- each  $e \in E \setminus A$  is not the largest edge in the unique cycle of  $A + e$  (when it exists)

Note: number of BCF sets is independent of edge ordering!

Example

We work with a simple transformation:

$$C_G(z) = a_0 z^n - a_1 z^{n-1} + \cdots + (-1)^n a_n$$

$$B_G(z) = a_0 + a_1 z + \cdots + a_n z^n = z^n C_G(-z^{-1}) = \sum_{F \subseteq E} z^{|F|} \text{BCF}$$

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Enough to show that whenever  $\Delta(G) \leq \Delta$  and  $|z| \leq 1/K\Delta$

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$\forall u \in V$  and some constants  $a \in (0, 1)$  and  $K > 0$  (to be determined).

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Enough to complete the induction ...

**Claim:** For each  $S \subseteq N(u)$ , have  $\left| \sum_{\substack{F+uS \text{ is} \\ \text{BCF tree}}} z^{|F|} \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right| \leq 1 + \varepsilon$

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If  $F$  occurs in second sum but not the first then

- Every non-trivial component of  $F$  hits  $S$
- Some component of  $F$  hits  $N(u)$  twice (or more)

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If  $F$  occurs in second sum but not the first then

- Every non-trivial component of  $F$  hits  $S$
- **Some component of  $F$  hits  $N(u)$  twice (or more)**

$$\left| \sum_{\substack{F+uS \text{ is} \\ \text{BCF tree}}} z^{|F|} \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right| \leq 1 + \left| \sum_{\text{BCF } F \in X} z^{|F|} \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right|$$

If BCF  $F$  occurs in second sum but not the first then

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$$\leq 1 + \sum_{F \in X} |z|^{|F|} \left| \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right|$$

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$$\begin{aligned}
 \left| \sum_{\substack{F+uS \text{ is} \\ \text{BCF tree}}} z^{|F|} \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right| &\leq 1 + \left| \sum_{\text{BCF } F \in X} z^{|F|} \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right| \\
 &\leq 1 + \sum_{F \in X} |z|^{|F|} \left| \frac{B_{G-u-S-V(F)}}{B_{G-u}} \right| \\
 &\leq 1 + \sum_{F \in X} (K\Delta)^{-|F|} (1-a)^{-|F|-|S|} \leq 1 + \varepsilon
 \end{aligned}$$

If BCF  $F$  occurs in second sum but not the first then

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 \end{aligned}$$

For final inequality we introduce and bound

$$T_{G,v_1,v_2}(x) = \sum_{\substack{T \text{ tree:} \\ v_1, v_2 \in V(T)}} x^{|T|} \quad \text{and bound} \quad T_{G,v_1,v_2} \left( \frac{\ln \alpha}{\alpha \Delta} \right) \leq \frac{\alpha \ln \alpha}{\Delta}$$

# Conclusion

$C_G(z) \neq 0$  whenever  $\Delta(G) \leq \Delta$  and  $|z| \geq 5.93\Delta$

- Try to improve on 5.93
- Unclear what the correct constant should be (perhaps complete bipartite graphs are extremal?)
- Can we leverage BCF characterisation for further progress?

# Further Results

Forest generating polynomial of  $G = (V, E)$

$$F_G(z) = \sum_{F \subseteq E \text{ forest}} z^{|F|}$$

- Also called partition function of arboreal gas model
- Our methods extend here, but can go further using a different recursion

$F_G(z) \neq 0$  whenever  $\Delta(G) \leq \Delta$  and  $|z| \leq 1/(2\Delta)$

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- Cannot replace  $1/(2\Delta)$  with  $1/\Delta$  due to  $\Delta$ -multi edge.
- Can we get close to  $1/\Delta$ ?