## University of Edinburgh INFR11156: Algorithmic Foundations of Data Science (2019) Homework 2

**Problem 1:** Suppose you want to estimate the unkown centre of the Gaussian in  $\mathbb{R}^d$  which has unit variance in each direction. Show that  $O(\log d/\epsilon^2)$  random samples form the Gaussian are sufficient to get an estimate  $\mathbf{m}_s$  of the true centre  $\boldsymbol{\mu}$ , so that with probability at least 99%,

$$\|\boldsymbol{\mu} - \mathbf{m}_s\|_{\infty} \leq \epsilon.$$

How many samples are sufficient to ensure that with probability at least 99%

$$\|\boldsymbol{\mu} - \mathbf{m}_s\|_2 \leq \epsilon?$$

Note that  $\|\mathbf{x}\|_{\infty} := \max_i |x_i|$ .

**Problem 2 (challenging):** This question is to try to design a dimension reduction lemma for  $\ell_1$ , similar to the Johnson-Lindenstrauss (JL) lemma for the Euclidean space. Remember that JL lemma says that we can pick a matrix  $\Phi$ , of dimension  $k \times d$  for large enough k, where each entry is chosen from a Gaussian distribution, such that: for any  $x \in \mathbb{R}^d$ , we have that  $\frac{1}{\sqrt{k}} \|\Phi x\|_2$  is a  $(1 + \epsilon)$  approximation to  $\|x\|_2$  with probability at least  $1 - e^{-\Omega(\epsilon^2 k)}$ .

For  $\ell_1$ , the equivalent of Gaussian distribution is the Cauchy distribution, which has probability distribution function  $p(x) = \frac{1}{\pi(1+x^2)}$ . Namely, the corresponding "stability" property of Cauchy distribution is the following. Consider  $s = \sum_{i=1}^{d} x_i c_i$ , for  $x \in \mathbb{R}^d$  and  $c_i$  each independently chosen from Cauchy distribution. Then s has distribution  $||x||_1 \cdot c$  where c is also distributed as a Cauchy distribution.

It is tempting to construct a dimensionality reducing map for  $\ell_1$  in the same way as what we did for Euclidean space, just by replacing the Gaussian distribution with the Cauchy distribution. In particular, let C be a matrix of size  $k \times d$ , where each entry is chosen independently from the Cauchy distribution.

1. Argue that this approach does *not* work for dimensionality reduction for  $\ell_1$ . Namely, for (say) k = 1000 and  $x = (1, 0, 0, \dots, 0)$ , the estimator  $\frac{1}{k} ||Cx||_1$  is not a 2-approximation to  $||x||_1 = 1$  with probability at least 10%.

In fact, it has been proven that there does *not* exist an equivalent dimensionality reduction for  $\ell_1$  at all. Instead, we will construct a sketch that provides a weaker form of "dimension reduction".

2. The median estimate is defined as the median of the absolute values of k coordinates of the vector Cx. Prove that for any  $x \in \mathbb{R}^d$ , the median estimate is a  $1 + \epsilon$  approximation to  $||x||_1$  with at least  $1 - e^{-\Omega(\epsilon^2 k)}$  probability. You might want to use the following concentration bound, called *Chernoff bound*: for any k independent and identically distributed random variables  $x_1, \ldots, x_k \in \{0, 1\}$ , each with expectation  $\mathbf{E}[x_i] = \mu \in [0, 1]$ , we have that

$$\mathbf{P}\left[\left|\frac{1}{k}\sum_{i}x_{i}-\mu\right| > \epsilon\right] \le e^{-\epsilon^{2}k}.$$

Note that we obtain a "sketch", which is not a regular dimension reduction scheme: namely, the "target" space is not  $\ell_1$ , but "median" (which is not even a metric/distance). Nevertheless, it is a linear map, and is useful for sketching and streaming as we will see in future lectures.