## University of Edinburgh

## INFR11156: Algorithmic Foundations of Data Science (2019)

## Lecture 5: Best-fit Subspaces and Singular Value Decomposition (1)

## 1 Singular values and singular vectors

Let $a=\left(a_{1}, \ldots, a_{d}\right)$ be a point in $\mathbb{R}^{d}$. We look at the projection of the points $a$ onto the line through the origin in the direction of $v$, see Figure 1 for illustration. Then we have that

$$
a_{1}^{2}+a_{2}^{2}+\ldots+a_{d}^{2}=(\text { length of projection })^{2}+(\text { distance of point to line })^{2},
$$

and therefore

$$
(\text { distance of point to line })^{2}=a_{1}^{2}+a_{2}^{2}+\ldots+a_{d}^{2}-(\text { length of projection })^{2} .
$$

Since $\sum_{j=1}^{d} a_{j}^{2}$ is constant independent of the line, minimising $a$ 's distance to the line is equivalent to maximising its projection onto the line.


Figure 1: The projection of the point $a$ onto the line through the origin in the direction of $v$.
Generalising the case above, we assume that there are $n$ points, each of which is represented in $\mathbb{R}^{d}$. We use matrix $A \in \mathbb{R}^{n \times d}$ to represent these $n$ points. Then, for any fixed direction $v \in \mathbb{R}^{d}$, the length of the projection of the $i$-th point, expressed by $A_{i}$, is $\left|\left\langle A_{i}, v\right\rangle\right|=\left|A_{i} v\right|$, and therefore the best-fit line is the one that maximises $\sum_{i=1}^{n}\left|A_{i} v\right|^{2}=\|A v\|^{2}$. With this in mind, we define the first singular vector $v_{1}$ of $A$ as

$$
v_{1} \triangleq \arg \max _{\|v\|=1}\|A v\| .
$$

The value $\sigma_{1}(A) \triangleq\left\|A v_{1}\right\|$ is called the first singular value of $A$. Notice that $\sigma_{1}^{2}(A)=\sum_{i=1}^{n}\left|A_{i} v_{1}\right|^{2}$ is the sum of the squared lengths of the projections of the points onto the line determined by $v_{1}$.

If the data points were all either on a line or close to a line, then intuitively $v_{1}$ should give us the direction of that line. However, if the data points are not close to that line but lie close to a 2-dimensional subspace, then further work is needed. We will look at the following greedy approach.

We start by finding $v_{1}$ and then find the best 2-dimensional subspace containing $v_{1}$. Notice that, for every 2 -dimensional subspace containing $v_{1}$, the sum of squared lengths of the projections onto the subspace equals the sum of squared projections onto $v_{1}$ plus the sum of squared projections along a vector perpendicular to $v_{1}$ in the subspace. Hence, instead of looking for the best 2-dimensional subspace containing $v_{1}$, we look for a unit vector vector $v$ perpendicular to $v_{1}$ that maximises $\|A v\|^{2}$ among all such unit vectors. This motivates the definition of the second singular vector $v_{2}$, which is the best-fit line perpendicular to $v_{1}$. Formally, we have

$$
v_{2} \triangleq \arg \max _{v \perp v_{1}\|v\|=1}\|A v\| .
$$

The value $\sigma_{2}(A)=\left\|A v_{2}\right\|$ is called the second singular value of $A$. The third singular vector $v_{3}$ and the thrid singular value are defined similarly by

$$
v_{3} \triangleq \arg \max _{\substack{v \perp v_{1}, v_{2} \\\|v\|=1}}\|A v\|
$$

and $\sigma_{3}(A)=\left\|A v_{3}\right\|$.
The greedy algorithm finds the $v_{1}$ that maximises $\|A v\|$ and then the best-fit 2-dimensional subspace containing $v_{1}$, etc. The following theorem shows that this simple greedy algorithm finds the best-fit subspace of every dimension.
Theorem 1 (The Greedy Algorithm Works). Let $A \in \mathbb{R}^{n \times d}$ be a matrix with singular vectors $v_{1}, \ldots, v_{r}$. For any $1 \leq k \leq r$, let $V_{k}$ be the subspace spanned by $v_{1}, \ldots, v_{k}$. Then, for every $k$, $V_{k}$ is the best-fit $k$-dimensional subspace for $A$.
Proof. The proof is by induction. The statement is obviously true for $k=1$. For $k=2$, let $W$ be a best-fit 2-dimensional subspace for $A$. For any orthonormal basis ( $w_{1}, w_{2}$ ) of $W$, $\left\|A w_{1}\right\|^{2}+\left\|A w_{2}\right\|^{2}$ is the sum of squared lengths of the projections of the rows of $A$ onto $W$. We choose an orthonormal basis $\left(w_{1}, w_{2}\right)$ of $W$ as follows:

1. If $v_{1}$ is perpendicular to $W$, any unit vector in $W$ that we choose as $w_{2}$ is perpendicular to $v_{1}$.
2. Otherwise, we choose $w_{2}$ to be the unit vector in $W$ perpendicular to the projection of $v_{1}$ onto $W$. This makes $w_{2}$ perpendicular to $v_{1}$.
Since $v_{1}$ maximises $\|A v\|^{2}$, it holds that $\left\|A w_{1}\right\|^{2} \leq\left\|A v_{1}\right\|^{2}$. Since $v_{2}$ maximises $\|A v\|^{2}$ overall $v$ perpendicular to $v_{1}$, we have that $\left\|A w_{2}\right\|^{2} \leq\left\|A v_{2}\right\|^{2}$. Thus, we have that

$$
\left\|A w_{1}\right\|^{2}+\left\|A w_{2}\right\|^{2} \leq\left\|A v_{1}\right\|^{2}+\left\|A v_{2}\right\|^{2}
$$

Hence, $V_{2}$ is at least as good as $W$ and is a best-fit 2-dimensional subspace.
This proof can be used inductively to prove the case for a general $k$.
The vectors $v_{1}, \ldots, v_{r}$ are called the right-singular vectors. We normalise these vectors and define

$$
u_{i} \triangleq \frac{1}{\sigma_{i}(A)} A v_{i} .
$$

These $u_{i}$ are called the left-singular vectors. It is easy to show that $u_{i}$ similarly maximises $\left\|u^{\top} A\right\|$ over all $u$ perpendicular to $u_{1}, \ldots u_{i-1}$, and these left-singular vectors are also orthogonal.
Lemma 2. Let $A \in \mathbb{R}^{n \times d}$. Then it holds that $\sum_{i=1}^{n} \sum_{j=1}^{d} a_{i j}^{2}=\sum_{i=1}^{r} \sigma_{i}^{2}(A)$.
Proof. Let $v_{1}, \ldots, v_{r}$ be the singular vectors of $A$. Then it holds that

$$
\sum_{j=1}^{n}\left\|A_{j}\right\|^{2}=\sum_{j=1}^{n} \sum_{i=1}^{r}\left(A_{j} v_{i}\right)^{2}=\sum_{i=1}^{r} \sum_{j=1}^{n}\left(A_{j} v_{i}\right)^{2}=\sum_{i=1}^{r}\left\|A v_{i}\right\|^{2}=\sum_{i=1}^{r} \sigma_{i}^{2}(A)
$$

The statement holds by noticing that $\sum_{j=1}^{n}\left\|A_{j}\right\|^{2}=\sum_{i=1}^{n} \sum_{j=1}^{d} a_{i j}^{2}$.

## 2 Singular Value Decomposition

We can think of $A \in \mathbb{R}^{n \times d}$ as a linear transformation taking a vector $v_{1}$ in its row space to a vector $u_{1}=A v_{1}$ in its column space. Many applications require to find an orthogonal basis for the row space and transform it into an orthogonal basis for the column space: $A v_{i}=\sigma_{i} u_{i}$. The heart of the problem is to find $v_{1}, \ldots, v_{r}$ for the row space of $A$ for which

$$
\begin{align*}
A\left[v_{1}, v_{2}, \ldots, v_{r}\right] & =\left[\sigma_{1} u_{1}, \sigma_{2} u_{2}, \ldots, \sigma_{r} u_{r}\right] \\
& =\left[u_{1}, u_{2}, \ldots, u_{r}\right]\left(\begin{array}{lll}
\sigma_{1} & & \\
& \ddots & \\
& & \sigma_{r}
\end{array}\right) . \tag{1}
\end{align*}
$$

Then, it is easy to see that the left and right-singular vectors $u_{i}=\frac{1}{\sigma_{i}} A v_{i}, v_{i}$, and their associated singular values $\sigma_{i}$ satisfy (1). With these vectors $u_{i} \mathrm{~s}, v_{i} \mathrm{~S}$, and the singular values $\sigma_{i} \mathrm{~s}$, we can write $A$ in matrix notation as

$$
A=U D V^{\top}
$$

where $u_{i}$ is the $i$-th column of $U, v_{i}^{\top}$ is the $i$-th row of $V^{\top}$, and $D$ is the diagonal matrix with $\sigma_{i}$ as the $i$-th entry on its diagonal. This factorisation of $A$ in the form of $U D V^{\top}$ is called Singular value decomposition. It is easy to check that

$$
A=\sum_{i=1}^{r} \sigma_{i} u_{i} v_{i}^{\top}
$$

