Jane Hillston

Joint work with Anastasis Georgoulas and Guido Sanguinetti, School of Informatics, University of Edinburgh

16th August 2017







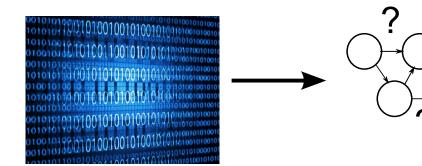
- 1 Introduction
- 2 Probabilistic Programming
- 3 ProPPA
- 4 Inference
- 5 Results
- 6 Conclusions

Outline

- 1 Introduction
- Probabilistic Programming
- 3 ProPPA
- 4 Inference
- 5 Results
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- Formal languages (process algebras, Petri nets, rule-based) provide a convenient interface for describing complex systems.
- High-level abstraction makes writing and manipulating models easier.
- They can capture different kinds of behaviour: deterministic, stochastic, ...
- Formal nature lends itself to automatic, rigorous methods for analysis and verification.
- ... but what if parts of the system are unknown?

Alternative perspective



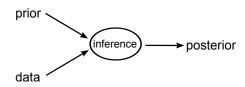
Model creation is data-driven

There are two approaches to model construction:

Machine Learning: extracting a model from the data generated by the system, or refining a model based on system behaviour using statistical techniques.

Mechanistic Modelling: starting from a description or hypothesis, construct a model that algorithmically mimics the behaviour of the system, validated against data.

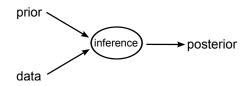
Machine Learning: Bayesian statistics



- Represent belief and uncertainty as probability distributions (prior, posterior).
- Treat parameters and unobserved variables similarly.
- Bayes' Theorem:

$$P(\theta \mid D) = \frac{P(\theta) \cdot P(D \mid \theta)}{P(D)}$$

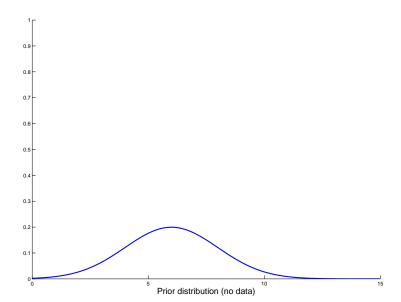
posterior ∝ prior · likelihoo

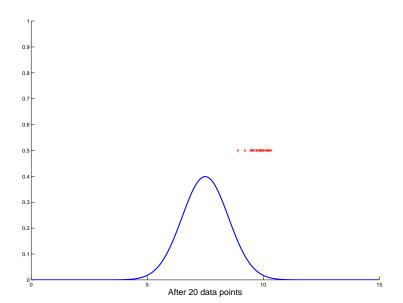


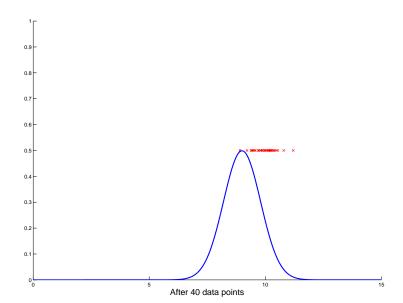
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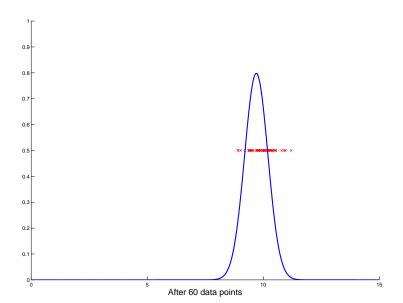
$$P(\theta \mid D) = \frac{P(\theta) \cdot P(D \mid \theta)}{P(D)}$$

posterior \propto prior \cdot likelihood









Mechanistic modelling

Models are constructed reflecting what is known about the components of the biological system and their behaviour.

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These are then compiled into executable models which can be run to deepen understanding of the model.

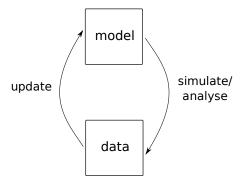
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Several approaches originating in theoretical computer science have been proposed to capture the system behaviour in a high-level way.

These are then compiled into executable models which can be run to deepen understanding of the model.

Executing the model generates data that can be compared with biological data.

Usual process of parameterising a model is iterative and manual.



Data-driven modelling:

- + rigorous handling of parameter uncertainty
- limited or no treatment of stochasticity
- in many cases bespoke solutions are required which can limit the size of system which can be handled

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Mechanistic modelling:

- + general execution "engine" (deterministic or stochastic) can be reused for many models
- + models can be used speculatively to investigate roles of parameters, or alternative hypotheses
 - parameters are assumed to be known and fixed, or costly approaches must be used to seek appropriate parameterisation

What if we could...

- include information about uncertainty in the model?
- automatically use observations to refine this uncertainty?
- do all this in a formal context?

Starting from the existing process algebra (Bio-PEPA), we have developed a new language ProPPA that addresses these issues.

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obabilistic programming

A programming paradigm for describing incomplete knowledge scenarios, and resolving the uncertainty.

 Describe how the data is generated in syntax like a conventional programming language, but leaving some variables uncertain.

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- Specify observations, which impose constraints on acceptable outputs of the program.

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- Run program forwards: Generate data consistent with observations.

A programming paradigm for describing incomplete knowledge scenarios, and resolving the uncertainty.

- Describe how the data is generated in syntax like a conventional programming language, but leaving some variables uncertain.
- Specify observations, which impose constraints on acceptable outputs of the program.
- Run program forwards: Generate data consistent with observations.
- Run program backwards: Find values for the uncertain variables which make the output match the observations.

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Q is the infinitesimal generator matrix characterising the CTMC.

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The language may be used to generate a Markov Process (CTMC).

Q is the infinitesimal generator matrix characterising the CTMC.

Models are typically executed by simulation using Gillespie's Stochastic Simulation Algorithm (SSA) or similar.

A Probabilistic Programming Process Algebra: ProPPA

The objective of ProPPA is to retain the features of the stochastic process algebra:

- simple model description in terms of components
- rigorous semantics giving an executable version of the model...

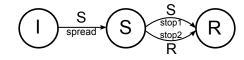
The objective of ProPPA is to retain the features of the stochastic process algebra:

- simple model description in terms of components
- rigorous semantics giving an executable version of the model...

... whilst also incorporating features of a probabilistic programming language:

- recording uncertainty in the parameters
- ability to incorporate observations into models
- access to inference to update uncertainty based on observations

Example Revisited



```
k_s = 0.5;
k_r = 0.1;

kineticLawOf spread : k_s * I * S;
kineticLawOf stop1 : k_r * S * S;
kineticLawOf stop2 : k_r * S * R;

I = (\text{spread},1) \downarrow ;
S = (\text{spread},1) \uparrow + (\text{stop1},1) \downarrow + (\text{stop2},1) \downarrow ;
R = (\text{stop1},1) \uparrow + (\text{stop2},1) \uparrow ;
I[10] \bowtie S[5] \bowtie R[0]
```

Additions

Declaring uncertain parameters:

- $\mathbf{k}_{-}\mathbf{s} = \text{Uniform}(0,1);$
- $k_t = Uniform(0,1);$

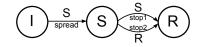
Providing observations:

observe('trace')

Specifying inference approach:

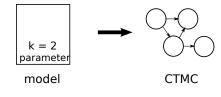
■ infer('ABC')

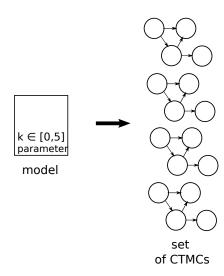
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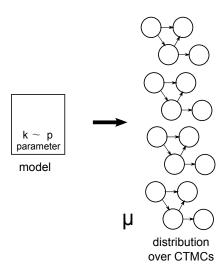


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I = (spread, 1) \downarrow ;
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R = (stop1,1) \uparrow + (stop2,1) \uparrow ;
I[10] \bowtie S[5] \bowtie R[0]
observe('trace')
infer('ABC') //Approximate Bayesian Computation
```

- A Bio-PEPA model can be interpreted as a CTMC; however, CTMCs cannot capture uncertainty in the rates (every transition must have a concrete rate).
- ProPPA models include uncertainty in the parameters, which translates into uncertainty in the transition rates.
- A ProPPA model should be mapped to something like a distribution over CTMCs.







Constraint Markov Chains (CMCs) are a generalization of DTMCs, in which the transition probabilities are not concrete, but can take any value satisfying some constraints.

Constraint Markov Chain

A CMC is a tuple $\langle S, o, A, V, \phi \rangle$, where:

- \blacksquare S is the set of states, of cardinality k.
- $o \in S$ is the initial state.
- A is a set of atomic propositions.
- $V: S \to 2^{2^A}$ gives a set of acceptable labellings for each state.
- $\phi: S \times [0,1]^k \to \{0,1\}$ is the constraint function.

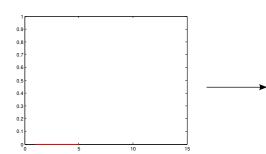
In a CMC, arbitrary constraints are permitted, expressed through the function ϕ : $\phi(s, \vec{p}) = 1$ iff \vec{p} is an acceptable vector of transition probabilities from state s.

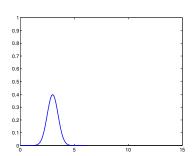
However,

Introduction

- CMCs are defined only for the discrete-time case, and
- this does not say anything about how likely a value is to be chosen, only about whether it is acceptable.

To address these shortcomings, we define **Probabilistic Constraint Markov Chains**.





A Probabilistic Constraint Markov Chain is a tuple (S, o, A, V, ϕ) , where:

- S is the set of states, of cardinality k.
- $o \in S$ is the initial state.
- A is a set of atomic propositions.
- $V: S \to 2^{2^A}$ gives a set of acceptable labellings for each state.
- $\phi: S \times [0,\infty)^k \to [0,\infty)$ is the constraint function.

- This is applicable to continuous-time systems.
- $\phi(s,\cdot)$ is now a probability density function on the transition rates from state s.

The semantics definition follows that of Bio-PEPA, which is defined using two transition relations:

- Capability relation is a transition possible?
- Stochastic relation gives rate of a transition

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The distribution over the parameter values induces a distribution over transition rates.

Rules are expressed as state-to-function transition systems (FuTS).

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This gives rise the underlying PCMC.

Probabilistic Constraint Markov Chains are open to two alternative dynamic interpretations:

- For each trajectory, for each uncertain transition rate, sample once at the start of the run and use that value throughout (Uncertain Markov Chains);
- During each trajectory, each time a transition with an uncertain rate is encountered, sample a value but then discard it and re-sample whenever this transition is visited again (Imprecise Markov Chains).

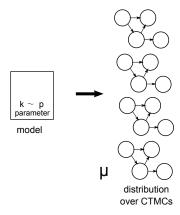
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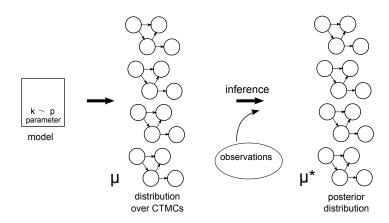
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Our current work is focused on the Uncertain Markov Chain case.

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Inference





$P(\theta \mid D) \propto P(\theta)P(D \mid \theta)$

- The ProPPA semantics does not define a single inference algorithm, allowing for a modular approach.
- Different algorithms can act on different input (time-series vs properties), return different results or in different forms.
- Exact inference is often impossible, as we cannot calculate the likelihood.
- We must use approximate algorithms or approximations of the system.

Transient probabilities can be expressed as:

$$rac{dp_i(t)}{dt} = \sum_{j
eq i} p_j(t) \cdot q_{ji} - p_i(t) \sum_{j
eq i} q_{ij}$$

Inferring likelihood in uncertain CTMCs

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The probability of a single observation (y, t) can then be expressed as

$$p(y,t) = \sum_{i \in S} p_i(t)\pi(y \mid i)$$

where $\pi(y \mid i)$ is the probability of observing y when in state i.

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The likelihood can then be expressed as

$$P(D \mid \theta) = \prod_{i=1}^{N} \sum_{i \in S} p_{(i\mid\theta)}(t_j) \pi(y_j \mid i)$$

For finite state-spaces, the transient probabilities can, in principle, be computed as

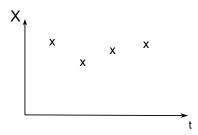
$$\mathbf{p}(t) = \mathbf{p}(0)e^{Qt}.$$

Likelihood is hard to compute:

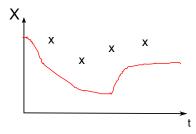
- Computing $e^{\mathbf{Q}t}$ is expensive if the state space is large
- Impossible directly in infinite state-spaces

- Approximate Bayesian Computation is a simple simulation-based solution:
 - Approximates posterior distribution over parameters as a set of samples
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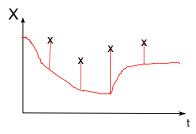


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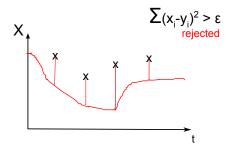
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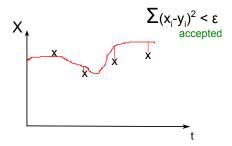


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Introduction

ABC algorithm

- **1** Sample a parameter set from the prior distribution.
- Simulate the system using these parameters.
- 3 Compare the simulation trace obtained with the observations.
- 4 If distance $< \epsilon$, accept, otherwise reject.

J.M.Marin, P.Pudlo, C.P.Robert, R.J.Ryder. Approximate Bayesian computational methods. Statistics and Computing 22(6) (2012)

This results in an approximation to the posterior distribution.

As $\epsilon \to 0$, set of samples converges to true posterior.

We use a more elaborate version based on Markov Chain Monte Carlo sampling.

state-space grows.

Various methods become inefficient or inapplicable as the

How to deal with unbounded systems?

- Multiple simulation runs
- Large population approximations (diffusion, Linear Noise,...)
- Systematic truncation
- Random truncations

■ We want to estimate the value of

$$f=\sum_{n=0}^{\infty}f_n$$

where the f_n 's are computable.

Russian Roulette Truncation

Introduction

■ We want to estimate the value of

$$f=\sum_{n=0}^{\infty}f_n$$

where the f_n 's are computable.

- Choose a single term f_k with probability p_k ; estimate $\hat{f} = \frac{f_k}{p_k}$
- \hat{f} is unbiased... but its variance can be high.

Russian Roulette Truncation

Introduction

■ We want to estimate the value of

$$f=\sum_{n=0}^{\infty}f_n$$

where the f_n 's are computable.

- Truncate the sum randomly: stop at term k with probability q_k .
- Form \hat{f} as a partial sum of the $f_n, n = 1, ..., k$, rescaled appropriately.

$$\hat{f} = \sum_{n=0}^{k} \frac{f_n}{\prod\limits_{j=0}^{k-1} (1 - q_j)}$$

Russian Roulette

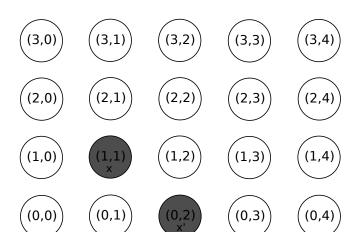
```
\hat{f} \leftarrow f_0
i \leftarrow 1
p \leftarrow 1
loop
       Choose to stop with probability q_i
       if stopping then
             return \hat{f}
       else

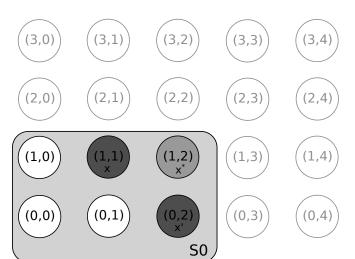
p \leftarrow p \cdot (1 - q_i)

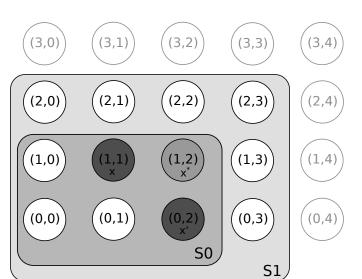
\hat{f} \leftarrow \hat{f} + \frac{f_i}{p}

i \leftarrow i + 1

       end if
end loop
```







Expanding the likelihood

The likelihood can be written as an infinite series:

$$p(x', t' \mid x, t) = \sum_{N=0}^{\infty} p^{(N)}(x', t' \mid x, t)$$
$$= \sum_{N=0}^{\infty} \left[f^{(N)}(x', t' \mid x, t) - f^{(N-1)}(x', t' \mid x, t) \right]$$

where

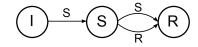
- $x^* = \max\{x, x'\}$
- $p^{(N)}(x', t' \mid x, t)$ is the probability of going from state x at time t to state x' at time t' through a path with maximum state $x^* + N$
- $f^{(N)}$ is the same, except the maximum state cannot exceed $x^* + N$ (but does not have to reach it)

$$\hat{f} = \sum_{n=0}^{k} \frac{f_n}{\prod_{j=0}^{k-1} (1 - q_j)}$$

- In our case, f is a probability that we wish to approximate, p(x', t' | x, t).
- Using \hat{f} instead of f leads to an error; however \hat{f} is unbiased: $E[\hat{f}] = f$.
- \hat{f} is also guaranteed to be positive.
- Pseudo-marginal algorithms can use this and still draw samples from the correct distribution.
- We have developed both Metropolis-Hastings and Gibbs-like sampling algorithms based on this approach.

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Results

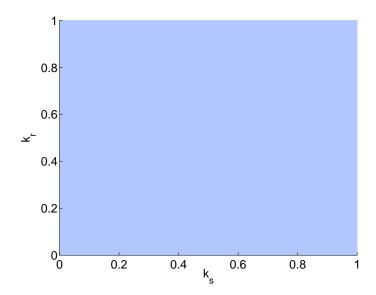


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I = (spread, 1) \downarrow ;
S = (spread, 1) \uparrow + (stop1, 1) \downarrow + (stop2, 1) \downarrow ;
R = (stop1,1) \uparrow + (stop2,1) \uparrow;
I[10] \bowtie S[5] \bowtie R[0]
observe('trace')
infer('ABC') //Approximate Bayesian Computation
```

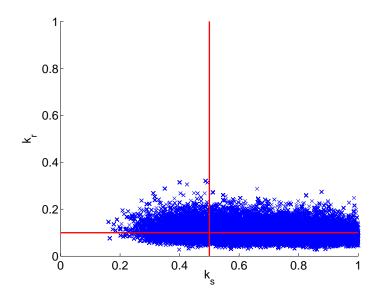
Tested on the rumour-spreading example, giving the two parameters uniform priors.

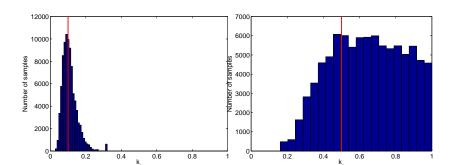
- Approximate Bayesian Computation
- Returns posterior as a set of points (samples)
- Observations: time-series (single simulation)

Results: ABC



Results: ABC





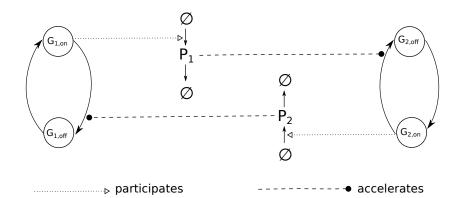
- Two mutually-repressing genes: promoters (unobserved) and their protein products
- Bistable behaviour: switching induced by environmental changes
- Synthesised in E. coli

Gardner, Cantor & Collins, Construction of a genetic toggle switch in Escherichia coli, Nature, 2000

Stochastic variant where switching is induced by noise

Tian & Burrage, Stochastic models for regulatory networks of the genetic toggle switch, PNAS, 2006

Genetic Toggle Switch



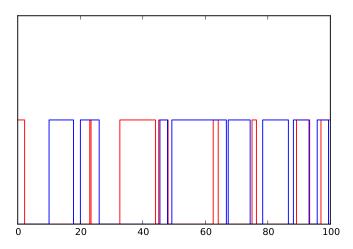
Toggle switch model: species

```
G1 = activ1 \uparrow + deact1 \downarrow + expr1 \oplus;
G2 = activ2 \uparrow + deact2 \downarrow + expr2 \oplus;
P1 = expr1 \uparrow + degr1 \downarrow + deact2 \oplus ;
P2 = expr2 \uparrow + degr2 \downarrow + deact1 \oplus
G1[1] <*> G2[0] <*> P1[20] <*> P2[0]
observe(toggle_obs);
infer(rouletteGibbs);
```

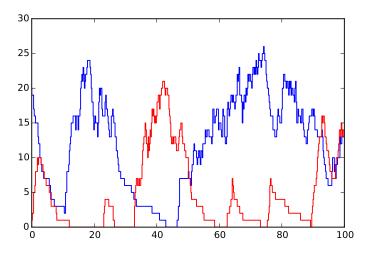
```
\theta_{-1} = \text{Gamma}(3,5); //\text{etc...}
kineticLawOf expr1 : \theta_{-}1 * G1;
kineticLawOf expr2 : \theta_{-}2 * G2;
kineticLawOf degr1 : \theta_{-}3 * P1;
kineticLawOf degr2 : \theta_{-}4 * P2;
kineticLawOf activ1 : \theta_{-}5 * (1 - G1);
kineticLawOf activ2 : \theta_{-}6 * (1 - G2);
kineticLawOf deact1 : \theta_{-}7 * \exp(r * P2) * G1;
kineticLawOf deact2 : \theta_-8 * \exp(r * P1) * G2;
G1 = activ1 \uparrow + deact1 \downarrow + expr1 \oplus;
G2 = activ2 \uparrow + deact2 \downarrow + expr2 \oplus;
P1 = expr1 \uparrow + degr1 \downarrow + deact2 \oplus ;
P2 = expr2 \uparrow + degr2 \downarrow + deact1 \oplus
G1[1] <*> G2[0] <*> P1[20] <*> P2[0]
observe(toggle_obs);
infer(rouletteGibbs);
```

- Simulated observations
- Gamma priors on all parameters (required by algorithm)
- Goal: learn posterior of 8 parameters
- 5000 samples taken using the Gibbs-like random truncation algorithm

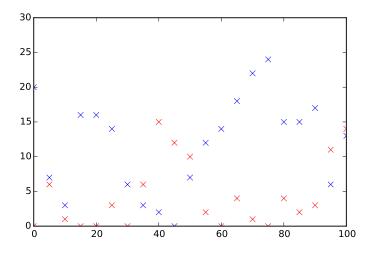
Promoters



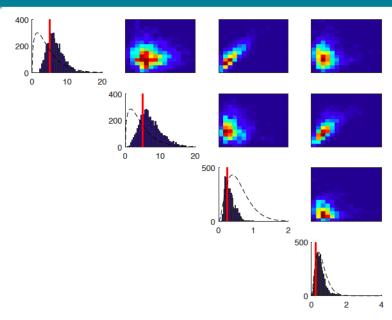
Proteins



Observations used



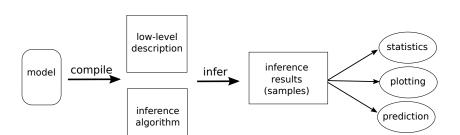
Results



Outille

- Introduction
- Probabilistic Programming
- 3 ProPPA
- 4 Inference
- 5 Results
- 6 Conclusions

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- ProPPA is a process algebra that incorporates uncertainty and observations directly in the model, influenced by probabilistic programming.
- Syntax remains similar to Bio-PEPA.
- Semantics defined in terms of an extension of Constraint Markov Chains.
- Observations can be either time-series or logical properties.
- Parameter inference based on random truncations (Russian Roulette) offers new possibilities for inference.

Anastasis Georgoulas



■ Guido Sanguinetti

