

POLYNOM\_3

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*C polynom_3_begin          ***** POLYNOM_3 *****
*C omral_com               =====
                           ORDERED-MONOID RING A-LISTS
                           =====
                           omral = o(rdered) m(onooid) a-l(ist)
*C omral_com_1            =====
                           DEFINITION OF OMRAL TYPE
                           =====

*D omralist_df            omral(<g:g:*>;<r:r:*>)== omralist{<g>; <r>}
*A omralist              omral(g;r) == oal(g↓oSet;r↓+gp)
*T omralist_wf 5.0 sec.
├ ∀g:OCMon. ∀r:CRng. omral(g;r) ∈ DSet
|
BY (Unfold 'omralist' 0 ...)
*M omralist_ml          note_reduction_strength 'omralist' '8';;
*D omral_dom_df
      dom{<g:g:*>, <r:r:*>}<ps:ps:*>== omral_dom{<g>; <r>; <ps>}
      dom(<ps:ps:*>)== omral_dom{<g>; <r>; <ps>}
*A omral_dom            dom(ps) == dom(ps)
*T omral_dom_wf 6.2 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ps:(|g| × |r|) List. dom(ps) ∈ MSet{g↓oSet}
|
BY (Unfold 'omral_dom' 0 ...)
*T omral_dom_wf2 8.8 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ps:|omral(g;r)|. dom(ps) ∈ FSet{g↓oSet}
|
BY (Unfold 'omral_dom' 0 ...)
*M omral_dom_eval
      let omral_dom_nilC =
        MacroC 'omral_dom_nilC'
          (EvalC 'omral_dom'
            [dom([])]
            IdC
            [0{g↓oSet}])
      ;;
      let omral_dom_cons_prC =
        MacroC 'omral_cons_prC'
          (EvalC 'omral_dom'
            [dom(<k, v>::ps)]
            (UnfoldC 'omral_dom'
              [mset_inj{g↓oSet}(k) + dom(ps)]
            ))
      ;;
      add_AbReduce_conv 'omral_dom'
        (omral_dom_nilC ORELSEC omral_dom_cons_prC)
      ;;
*T omralist_car_properties 5.7 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ws:|omral(g;r)|. ↑sd_ordered(map(λx.x.1;ws)) ∧ ¬↑(0 ∈b map(λx.x.2;ws))
|
BY ProveSpecializedLemma 'oalist_car_properties' 2 [[parm{i}]; [g↓oSet]; [r↓+gp]]
  ] ((AbReduceIfC (\e t.not is_term 'oalist' (subterm t 1))
    ANDTHENC TryC (Folds
      C
        'omralist omral_plus grp_lt grp_leq grp_blt'))))
*M oal_to_omral
      % Lifting Theorems from oalists to omralists %
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let omral_opids =
  ‘‘omralist omral_plus omral_dom grp_lt grp_leq grp_blt
    omral_zero omral_minus omral_inj‘‘
;;
let OmRalC =
  ForceReduceC ‘5‘ ANDTHENC TryC (FoldsC omral_opids) ;;
let OmRalCStr =
  "ForceReduceC ‘5‘ ANDTHENC TryC (FoldsC ‘“"
  J
  concatenate_strings
    (map (\id.tok_to_string id J " ") omral_opids)
  J
  “‘)”
;;
let mk_omral_thm old_name new_name new_pos =
  add_specialized_theorem
    old_name
    [‘g’,[OCMon];‘r’,[CRng]] % New outer context %
    [‘[parm{i}]’;‘g↓oset’;‘r↓+gp’] % Bindings for outer context of old thm %
    OmRalC
    OmRalCStr
    new_name
    new_pos
    ; refresh()
;;
*T rng_before_imp_before_all 5.7 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀ps:|omral(g;r)|.
|   ↑before(k;map(λz.z.1;ps)) ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). x <_b k)
|
BY ProveSpecializedLemma ‘before_imp_before_all‘ 2 [‘[parm{i}]’;‘g↓oset’;‘r↓+gp’]
] ((AbReduceIfC (\e t.not is_term ‘oalist‘ (subterm t 1))
  ANDTHENC FoldsC
  ‘
  ‘omralist omral_plus grp_lt grp_leq grp_blt‘))
*T rng_before_all_imp_before 5.8 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀ps:(|g| × |r|) List.
|   ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). x <_b k) ⇒ ↑before(k;map(λz.z.1;ps))
|
BY ProveSpecializedLemma ‘before_all_imp_before‘ 2 [‘[parm{i}]’;‘g↓oset’;‘r↓+gp’]
] ((AbReduceIfC (\e t.not is_term ‘oalist‘ (subterm t 1))
  ANDTHENC FoldsC
  ‘
  ‘omralist omral_plus grp_lt grp_leq grp_blt‘))
*T omralist_cases 5.5 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀Q:|omral(g;r)| → ℙ.
|   Q[[]]
|   ⇒ (∀ws:|omral(g;r)|. ∀x:|g|. ∀y:|r|.
|     ↑before(x;map(λx.x.1;ws)) ⇒ ¬(y = 0) ⇒ Q[<x, y>::ws])
|   ⇒ {∀ws:|omral(g;r)|. Q[ws]}
|
BY (D 0 THENM D 0 ...a)
| THEN AssertLemma ‘oalist_cases_a‘ []
| THEN (With ‘g↓oset’ (D (-1)) THENM With ‘r↓+gp’ (D (-1)) ...a)
|
1. g: OCMon
2. r: CRng
3. ∀Q:|oal(g↓oset;r↓+gp)| → ℙ

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Q[[]]
⇒ (∀ws:|oal(g↓oset;r↓+gp)|. ∀x:|(g↓oset)|. ∀y:|r↓+gp|.
  ↑before(x;map(λx.x.1;ws)) ⇒ ¬(y = e) ⇒ Q[<x, y>::ws])
⇒ {∀ws:|oal(g↓oset;r↓+gp)|. Q[ws]}
├ ∀Q:|omral(g;r)| → ℙ
|   Q[[]]
|   ⇒ (∀ws:|omral(g;r)|. ∀x:|g|. ∀y:|r|.
|       ↑before(x;map(λx.x.1;ws)) ⇒ ¬(y = 0) ⇒ Q[<x, y>::ws])
|   ⇒ {∀ws:|omral(g;r)|. Q[ws]}
|
BY % This is too fiddly. Need better way of doing this %
  AbReduceIf (\e t.not is_term 'oalist' (subterm t 1)) (-1)
  THEN Unfold 'omralist' 0 THEN Trivial
*T omralist_ind_a 7.0 sec.
├ ∀g:OCMon. ∀r:CRng. ∀Q:|omral(g;r)| → ℙ.
|   Q[[]]
|   ⇒ (∀ws:|omral(g;r)|
|       Q[ws] ⇒ (∀x:|g|. ∀y:|r|. ↑before(x;map(λx.x.1;ws)) ⇒ ¬(y = 0) ⇒ Q[<x, y>::ws]))
|   ⇒ {∀ws:|omral(g;r)|. Q[ws]}
|
BY ProveSpecializedLemma 'oalist_ind_a' 2 [⌈parm{i}⌉;⌈g↓oset⌉;⌈r↓+gp⌉]
  ] ((AbReduceIfC (\e t.not is_term 'oalist' (subterm t 1))
    ANDTHENC TryC (Folds
      C
        'omralist omral_plus grp_lt grp_leq grp_blt'))))
*T omral_lookups_same_a 6.3 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. (∀u:|g|. ps[u] = qs[u]) ⇒ ps = qs
|
BY ProveSpecializedLemma 'lookups_same_a' 2 [⌈parm{i}⌉;⌈g↓oset⌉;⌈r↓+gp⌉]
  ] ((AbReduceIfC (\e t.not is_term 'oalist' (subterm t 1))
    ANDTHENC TryC (Folds
      C
        'omralist omral_plus grp_lt grp_leq grp_blt'))))
*T rng_lookup_before_start 5.3 sec.
├ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀ps:|omral(g;r)|. ↑before(k;map(λz.z.1;ps)) ⇒ ps[k] = 0
|
BY (D 0 THENM D 0
|   THENM AssertLemma 'lookup_before_start' []
|   THENM With ⌈g↓oset⌉ (D (-1)) THENM With ⌈r↓+gp⌉ (D (-1)) ...a)
|
1. g: OCMon
2. r: CRng
3. ∀k:|(g↓oset)|. ∀ps:|oal(g↓oset;r↓+gp)|. ↑before(k;map(λz.z.1;ps)) ⇒ ps[k] = e
├ ∀k:|g|. ∀ps:|omral(g;r)|. ↑before(k;map(λz.z.1;ps)) ⇒ ps[k] = 0
|
BY AbReduceIf (\e t.not is_term 'oalist' (subterm t 1)) 3
|   THEN Fold 'omralist' 3
|
3. ∀k:|g|. ∀ps:|omral(g;r)|. ↑before(k;map(λz.z.1;ps)) ⇒ ps[k] = 0
|
BY Trivial
*T lookup_omral_eq_zero 6.2 sec.
├ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀ps:|omral(g;r)|. ¬↑(k ∈b dom(ps)) ⇒ ps[k] = 0
|
BY ProveSpecializedLemma 'lookup_oal_eq_id' 2 [⌈parm{i}⌉;⌈g↓oset⌉;⌈r↓+gp⌉]
  ] ((AbReduceIfC (\e t.not is_term 'oalist' (subterm t 1))
    ANDTHENC TryC (Folds

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C
  ‘‘omralist omral_plus omral_dom grp_lt grp_leq grp_blt‘‘))
*C omral_plus_com
  =====
  OMRAL PLUS FUNCTION
  =====
  Lifting of oal merge function
*D omral_plus_df
  Parens ::Prec(inop)::
    <ps:ps:L> ++<g:g:L>,<r:r:L> <qs:qs:L>
    == omral_plus{<g>; <r>; <ps>; <qs>}
  Parens ::Prec(inop)::
    <ps:ps:L> ++ <qs:qs:L>
    == omral_plus{<g>; <r>; <ps>; <qs>}
*A omral_plus      ps ++ qs == ps ++ qs
*T omral_plus_wf  9.4 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:(|g| × |r|) List. ps ++ qs ∈ (|g| × |r|) List
|
BY (Unfold ‘omral_plus‘ 0 THEN RepD ...a)
|
1. g: OCMon
2. r: CRng
3. ps: (|g| × |r|) List
4. qs: (|g| × |r|) List
⊢ ps ++ qs ∈ (|g| × |r|) List
|
BY % MemCD picks second wf lemma which is not wanted here %
  (BLemma ‘oal_merge_wf‘ ...)
*T omral_plus_sd_ordered  6.2 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:(|g| × |r|) List.
|   ↑sd_ordered(map(λx.x.1;ps))
|   ⇒ ↑sd_ordered(map(λx.x.1;qs))
|   ⇒ ↑sd_ordered(map(λx.x.1;ps ++ qs))
|
BY ProveSpecializedLemma ‘oal_merge_sd_ordered‘ 2 [⌈parm{i}⌋;⌈g↓oset⌋;⌈r↓+gp⌋]
] ((AbReduceIfC (λe t.not is_term ‘oalist‘ (subterm t 1))
  ANDTHENC TryC (Folds
    C
      ‘‘omralist omral_plus grp_lt grp_leq grp_blt‘‘))
*T omral_plus_non_zero_vals  7.2 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:(|g| × |r|) List.
|   ¬↑(0 ∈b map(λx.x.2;ps)) ⇒ ¬↑(0 ∈b map(λx.x.2;qs)) ⇒ ¬↑(0 ∈b map(λx.x.2;ps ++ qs))
|
BY ProveSpecializedLemma ‘oal_merge_non_id_vals‘ 2 [⌈parm{i}⌋;⌈g↓oset⌋;⌈r↓+gp⌋]
] ((AbReduceIfC (λe t.not is_term ‘oalist‘ (subterm t 1))
  ANDTHENC TryC (Folds
    C
      ‘‘omralist omral_plus grp_lt grp_leq grp_blt‘‘))
*T omral_plus_wf2  5.6 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. ps ++ qs ∈ |omral(g;r)|
|
BY ProveSpecializedLemma ‘oal_merge_wf2‘ 2 [⌈parm{i}⌋;⌈g↓oset⌋;⌈r↓+gp⌋]
] ((AbReduceIfC (λe t.not is_term ‘oalist‘ (subterm t 1))
  ANDTHENC TryC (Folds
    C
      ‘‘omralist omral_plus omral_dom grp_lt grp_leq grp_blt‘‘))
*T omral_plus_dom  6.1 sec.

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├  $\forall g:OCMon. \forall r:CRng. \forall ps,qs:|omral(g;r)|. \uparrow(\text{dom}(ps ++ qs) \subseteq_b \text{dom}(ps) \cup \text{dom}(qs))$ 
|
BY ProveSpecializedLemma ‘oal_dom_merge’ 2 [[parm{i}]];[g↓oset];[r↓+gp]
] ((AbReduceIfC (\e t.not is_term ‘oalist’ (subterm t 1))
  ANDTHENC TryC (Folds
    C
      ‘‘omralist omral_plus omral_dom grp_lt grp_leq grp_blt’’)))
*T lookup_omral_plus 7.0 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall k:|g|. \forall ps,qs:|omral(g;r)|. (ps ++ qs)[k] = ps[k] +r qs[k]$ 
|
BY ProveSpecializedLemma ‘lookup_merge’ 2 [[parm{i}]];[g↓oset];[r↓+gp]
] ((AbReduceIfC (\e t.not is_term ‘oalist’ (subterm t 1))
  ANDTHENC TryC (Folds
    C
      ‘‘omralist omral_plus grp_lt grp_leq grp_blt’’)))
*T omral_plus_comm 6.3 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall ps,qs:|omral(g;r)|. ps ++ qs = qs ++ ps$ 
|
BY ProveSpecializedLemma ‘oal_merge_comm’ 2 [[parm{i}]];[g↓oset];[r↓+gp]
] (ForceReduceC ‘5’ ANDTHENC TryC (FoldsC ‘‘omralist omral_plus omral_dom grp_lt
  grp_leq grp_blt omral_zero omral_minus omral_inj ‘’))
*T omral_plus_assoc 7.2 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall ps,qs,rs:|omral(g;r)|. ps ++ (qs ++ rs) = (ps ++ qs) ++ rs$ 
|
BY ProveSpecializedLemma ‘oal_merge_assoc’ 2 [[parm{i}]];[g↓oset];[r↓+gp]
] (ForceReduceC ‘5’ ANDTHENC TryC (FoldsC ‘‘omralist omral_plus omral_dom grp_lt
  grp_leq grp_blt omral_zero omral_minus omral_inj ‘’))
*C omral_zmi_com
=====
OMRAL ZERO, MINUS AND INJECTION FUNCTIONS
=====
All lifted from oal development.
*D omral_zero_df 00<g:g:*,<r:r:*>== omral_zero{<g>; <r>}
*A omral_zero 00g,r == 00
*T omral_zero_wf 4.7 sec.
├  $\forall g:OCMon. \forall r:CRng. 00g,r \in |omral(g;r)|$ 
|
BY (Unfolds ‘‘omral_zero omralist’’ 0 ...)
*D omral_minus_df
  Parens ::Prec(preop)::
    --<g:g:L>,<r:r:L> <ps:ps:L>
    == omral_minus{<g>; <r>; <ps>}
  Parens ::Prec(preop):: --<ps:ps:L>== omral_minus{<g>; <r>; <ps>}
*A omral_minus --ps == --ps
*T omral_minus_wf 17.0 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall ps:|omral(g;r)|. --ps \in |omral(g;r)|$ 
|
BY (Unfold ‘omral_minus’ 0 ...)
*D omral_inj_df
  inj{<g:g:*,<r:r:*>}<k:k:*,<v:v:*>== omral_inj{<g>; <r>; <k>; <v>}
  inj{<k:k:*,<v:v:*>}== omral_inj{<g>; <r>; <k>; <v>}
*A omral_inj inj(k,v) == inj(k,v)
*T omral_inj_wf 10.7 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall k:|g|. \forall v:|r|. inj(k,v) \in |omral(g;r)|$ 
|
BY (Unfold ‘omral_inj’ 0 ...)
*T omral_dom_inj 6.8 sec.

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┆ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀v:|r|.
|   dom(inj(k,v)) = if v =b 0 then 0{g↓oset} else mset_inj{g↓oset}(k) fi
|
BY ProveSpecializedLemma 'oal_dom_inj' 2 [[parm{i}];[g↓oset];[r↓+gp]
] (ForceReduceC '5' ANDTHENC TryC (Foldsc "'omralist omral_plus omral_dom grp_lt
grp_leq grp_blt omral_zero omral_minus omral_inj '"))
*T lookup_omral_inj 6.7 sec.
┆ ∀g:OCMon. ∀r:CRng. ∀k,k':|g|. ∀v:|r|. inj(k,v)[k'] = when k =b k'. v
|
BY ProveSpecializedLemma 'lookup_oal_inj' 2 [[parm{i}];[g↓oset];[r↓+gp]
] (ForceReduceC '5' ANDTHENC TryC (Foldsc "'omralist omral_plus omral_dom grp_lt
grp_leq grp_blt omral_zero omral_minus omral_inj '"))
*T comb_for_omral_inj_wf 1.4 sec.
┆ (λg,r,k,v,z.inj(k,v)) ∈ g:OCMon → r:CRng → k:|g| → v:|r| → ↓True → |omral(g;r)|
|
BY ProveOpCombTyping 'omral_inj_wf'
*T omral_fact 7.0 sec.
┆ ∀g:OCMon. ∀r:CRng. ∀ps:|omral(g;r)|.
|   ps = msFor{oal_mon(g↓oset;r↓+gp)} k' ∈ dom(ps). inj(k',ps[k'])
|
BY ProveSpecializedLemma 'oalist_fact' 2 [[parm{i}];[g↓oset];[r↓+gp]
] (ForceReduceC '5' ANDTHENC TryC (Foldsc "'omralist omral_plus omral_dom grp_lt
grp_leq grp_blt omral_zero omral_minus omral_inj '"))
*T omral_fact_a 4.9 sec.
┆ ∀g:OCMon. ∀r:CRng. ∀ps:|omral(g;r)|.
|   ps = msFor{omral_alg(g;r)↓grp} k' ∈ dom(ps). inj(k',ps[k'])
|
BY % Uggh ! %
| RWH (MacroC 'x'
|   (EvalC "'mset_for mon_for'"
|     ANDTHENC UnfoldsC "'omral_plus omral_zero'" ) [msFor{omral_alg(g;r)↓grp} x ∈ a
|                                                         f[x]1
|   (EvalC "'mset_for mon_for'"
|     [msFor{oal_mon(g↓oset;r↓+gp)} x ∈ a. f[x]1] 0
|
┆ ∀g:OCMon. ∀r:CRng. ∀ps:|omral(g;r)|.
|   ps = msFor{oal_mon(g↓oset;r↓+gp)} k' ∈ dom(ps). inj(k',ps[k'])
|
BY Lemma 'omral_fact'
*C omral_scale_com
=====
OMRAL SCALING FUNCTION
=====
Scales keys and values of an omralist.
*D omral_scale_df
Parens ::Prec(preop)::
  <<k:k:*,<v:v:*>>*<g:mon:L>,<r:rng:L> <ps:ps:E>
  == omral_scale{<g>; <r>; <k>; <v>; <ps>}
Parens ::Prec(preop)::
  <<k:k:*,<v:v:*>>* <ps:ps:E>
  == omral_scale{<g>; <r>; <k>; <v>; <ps>}
*M omral_scale_ml
  <k,v>* ps
  ==r case ps of
    [] => []
    p::ps' => if (v * p.2) =b 0
              then <k,v>* ps'

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else <k * p.1, v * p.2>::(<k,v>* ps')
fi
esac
*M omral_scale_eval
  let omral_scale_nilC =
    FwdMacroC 'omral_scale_nilC'
    (RecEvalC 'omral_scale' ' ' [ <k,v>* [] ] ;;
  let omral_scale_cons_prC =
    FwdMacroC 'omral_scale_cons_prC'
    (RecEvalC 'omral_scale' ' ' [ <k,v>* (<k', v'>::ps) ] ;;
  add_AbReduce_conv 'omral_scale'
    (omral_scale_nilC ORELSEC omral_scale_cons_prC) ;;
*T omral_scale_wf 4.2 sec.
⊢ ∀g:GrpSig. ∀r:RngSig. ∀k:|g|. ∀v:|r|. ∀ps:(|g| × |r|) List. <k,v>* ps ∈ (|g| × |r|) List
|
BY (RepD
| THENM New ['p';'ps\'] (OnVar 'ps' ListInd)
| THENM AbReduce 0 ...a)
| \
| 1. g: GrpSig
| 2. r: RngSig
| 3. k: |g|
| 4. v: |r|
| 5. ps: (|g| × |r|) List
| ⊢ [] ∈ (|g| × |r|) List
| |
1 BY Auto
\
1. g: GrpSig
2. r: RngSig
3. k: |g|
4. v: |r|
5. ps: (|g| × |r|) List
6. p: |g| × |r|
7. ps': (|g| × |r|) List
8. <k,v>* ps' ∈ (|g| × |r|) List
⊢ <k,v>* (p::ps') ∈ (|g| × |r|) List
|
BY % Have to be careful here. Exists a later wf lemma
| which would also apply. The price to pay for having
| no half-way decent library object dependency tracking. %
|
|
| (D 6 THENM AbReduce 0
| THENM MemCD ...a)
| \
| 6. p1: |g|
| 7. p2: |r|
| 8. ps': (|g| × |r|) List
| 9. <k,v>* ps' ∈ (|g| × |r|) List
| ⊢ (v * p2) =b 0 ∈ ℤ
| |
1 BY Auto
| \
| 6. p1: |g|
| 7. p2: |r|
| 8. ps': (|g| × |r|) List

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| 9. <k,v>* ps' ∈ (|g| × |r|) List
| ⊢ <k,v>* ps' ∈ (|g| × |r|) List
| |
1 BY Trivial
  \
  6. p1: |g|
  7. p2: |r|
  8. ps': (|g| × |r|) List
  9. <k,v>* ps' ∈ (|g| × |r|) List
  ⊢ (<k * p1, v * p2>::(<k,v>* ps')) ∈ (|g| × |r|) List
  |
  BY (MemCD ...a)
  | \
  | ⊢ <k * p1, v * p2> ∈ |g| × |r|
  | |
  1 BY Auto
    \
    ⊢ <k,v>* ps' ∈ (|g| × |r|) List
    |
    BY Auto
*T omral_scale_dom_pred 27.9 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀Q:|g| → ℤ. ∀k:|g|. ∀v:|r|. ∀ps:(|g| × |r|) List.
|   ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). Q[k * x]) ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* ps). Q[x])
|
BY (CDToVarThen 'ps' ListIndA ...a)
| \
| 1. g: OCMon
| 2. r: CRng
| 3. Q: |g| → ℤ
| 4. k: |g|
| 5. v: |r|
| ⊢ ↑(∀bx(:|g|) ∈ map(λz.z.1;[]). Q[k * x]) ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* []). Q[x])
| |
1 BY (Reduce 0 ...)
  \
  1. g: OCMon
  2. r: CRng
  3. Q: |g| → ℤ
  4. k: |g|
  5. v: |r|
  6. p: |g| × |r|
  7. ps: (|g| × |r|) List
  8. ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). Q[k * x]) ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* ps). Q[x])
  ⊢ ↑(∀bx(:|g|) ∈ map(λz.z.1;p::ps). Q[k * x])
  | ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* (p::ps)). Q[x])
  |
  BY New ['kp';'vp'] (D 6) THEN Reduce 0
  | THEN (D 0 THENM RW bool_to_propC (-1)
  |   THENM D (-1) ...a)
  |
  6. kp: |g|
  7. vp: |r|
  8. ps: (|g| × |r|) List
  9. ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). Q[k * x]) ⇒ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* ps). Q[x])
  10. ↑Q[k * kp]
  11. ↑(∀bx(:|g|) ∈ map(λz.z.1;ps). Q[k * x])
  ⊢ ↑(∀bx(:|g|) ∈ map(λz.z.1;if (v * vp) =b 0

```



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|       then <k,v>* ps
|       else <k * kp, v * vp>::(<k,v>* ps)
|       fi )
|       Q[x])
|
BY (SplitOnConclITE THENM Reduce 0 ...a)
| \
| 12. v * vp = 0
|  ⊢ ↑(∀bx(:|g|) ∈ map(λz.z.1;<k,v>* ps). Q[x])
|  |
| 1 BY (HypBackchain ...)
|  \
|  12. ¬(v * vp = 0)
|  ⊢ ↑(Q[k * kp] ∧b (∀bx(:|g|) ∈ map(λz.z.1;<k,v>* ps). Q[x]))
|  |
|  BY (RW bool_to_propC 0
|      THENM HypBackchain ...)
*T omral_dom_scale 93.5 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀v:|r|. ∀ps:|omral(g;r)|.
|   ↑(dom(<k,v>* ps) ⊆b fs-map(λk'.k' * k, dom(ps)))
|
BY (RepD THENM BLemma 'mem_bsubset' THENM RepD ...a)
|
1. g: OCMon
2. r: CRng
3. k: |g|
4. v: |r|
5. ps: |omral(g;r)|
6. x: |(g↓oset)|
7. ↑(x ∈b dom(<k,v>* ps))
⊢ ↑(x ∈b fs-map(λk'.k' * k, dom(ps)))
|
BY (Negate 0 THENM Negate 7 ...a)
|
7. ¬↑(x ∈b fs-map(λk'.k' * k, dom(ps)))
⊢ ¬↑(x ∈b dom(<k,v>* ps))
|
BY (Unfold 'fset_map' 7
|   THENM RWH (LemmaC 'fset_of_mset_mem') 7 ...a)
|
7. ¬↑(x ∈b msmset{g↓oset,g↓oset}(λk'.k' * k;dom(ps)))
|
BY RepUnfolds 'omral_dom oal_dom' 7
| THENM (RWH (LemmaC 'mset_map_char') 7 ...a)
|
7. ¬↑(x ∈b mk_mset(map(λk'.k' * k;map(λz.z.1;ps))))
|
BY (RWH (LemmaWithC ['C', '|g|'] 'map_map') 7 ...a)
|
7. ¬↑(x ∈b mk_mset(map((λk'.k' * k) o (λz.z.1);ps)))
|
BY Unfold 'compose' 7 THEN Reduce 7
|
7. ¬↑(x ∈b mk_mset(map(λx.x.1 * k;ps)))
|
BY % blow away mset stuff %
|

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| Reduce 6 THEN RenameVar 'k\' 6
| THEN OnCls [0;7] (RepUnfolds 'omral_dom oal_dom mk_mset mset_mem mem')
| THEN OnCls [0;7] (Fold 'bexists')
|
6. k': |g|
7.  $\neg \uparrow(\exists_b x(:|(g \downarrow \text{oset})|) \in \text{map}(\lambda x.x.1 * k; \text{ps}). x =_b k')$ 
 $\vdash \neg \uparrow(\exists_b x(:|(g \downarrow \text{oset})|) \in \text{map}(\lambda z.z.1; \langle k, v \rangle * \text{ps}). x =_b k')$ 
|
BY (OnMCls [0;7] (RW (SweepDnC
| (LemmaC 'bnot_thru_exists'
| ORELSEC RevLemmaC 'assert_of_bnot')))) ...a)
|
7.  $\uparrow(\forall_b x(:|(g \downarrow \text{oset})|) \in \text{map}(\lambda x.x.1 * k; \text{ps}). \neg_b(x =_b k'))$ 
 $\vdash \uparrow(\forall_b x(:|(g \downarrow \text{oset})|) \in \text{map}(\lambda z.z.1; \langle k, v \rangle * \text{ps}). \neg_b(x =_b k'))$ 
|
BY (OnCls [0;7] Reduce THENM BLemma 'omral_scale_dom_pred' ...a)
|
7.  $\uparrow(\forall_b x(:|g|) \in \text{map}(\lambda x.x.1 * k; \text{ps}). \neg_b(x =_b k'))$ 
 $\vdash \uparrow(\forall_b x(:|g|) \in \text{map}(\lambda z.z.1; \text{ps}). \neg_b((k * x) =_b k'))$ 
|
BY % Push both map funs onto mon_for arg %
| (OnMCls [0;7]
| (\i.Unfold 'ball' i
| THENM RWH (LemmaC 'mon_for_map') i
| THENM Fold 'ball' i) ...a)
|
7.  $\uparrow(\forall_b x(:|(g \downarrow \text{oset} \times r \downarrow + \text{gp} \downarrow \text{set})|) \in \text{ps}. \neg_b(((\lambda x.x.1 * k) x) =_b k'))$ 
 $\vdash \uparrow(\forall_b x(:|(g \downarrow \text{oset} \times r \downarrow + \text{gp} \downarrow \text{set})|) \in \text{ps}. \neg_b((k * ((\lambda z.z.1) x)) =_b k'))$ 
|
BY OnCls [0;7] Reduce
|
7.  $\uparrow(\forall_b x(:|g| \times |r|) \in \text{ps}. \neg_b((x.1 * k) =_b k'))$ 
 $\vdash \uparrow(\forall_b x(:|g| \times |r|) \in \text{ps}. \neg_b((k * x.1) =_b k'))$ 
|
BY (RWH (LemmaC 'abmonoid_comm') 0 ...)
*T omral_scale_dom_bound 33.5 sec.
 $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall \text{bound}, k:|g|. \forall v:|r|. \forall \text{ps}:(|g| \times |r|) \text{List}.$ 
|  $\uparrow(\forall_b x(:|g|) \in \text{map}(\lambda z.z.1; \text{ps}). x <_b \text{bound})$ 
|  $\Rightarrow \uparrow(\forall_b x(:|g|) \in \text{map}(\lambda z.z.1; \langle k, v \rangle * \text{ps}). x <_b k * \text{bound})$ 
|
BY (RepD THENM BLemma 'omral_scale_dom_pred' ...a)
|
1. g: OCMon
2. r: CRng
3. bound: |g|
4. k: |g|
5. v: |r|
6. ps: (|g| × |r|) List
7.  $\uparrow(\forall_b x(:|g|) \in \text{map}(\lambda z.z.1; \text{ps}). x <_b \text{bound})$ 
 $\vdash \uparrow(\forall_b x(:|g|) \in \text{map}(\lambda z.z.1; \text{ps}). k * x <_b k * \text{bound})$ 
|
BY % This is so ugly! (ball_char wants to match a set_car,
| not a grp_car) %
| (Assert [|g| = |(g ↓ oset)|] THENA Reduce 0 ...)
| THEN (OnMCls [0;7] (RewriteWith [-1] 'ball_char') ...a)
| THEN Thin (-1)
|

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7.  $\forall x: |(g \downarrow \text{oset})|. \uparrow(x \in_b \text{map}(\lambda z.z.1; \text{ps})) \Rightarrow \uparrow(x <_b \text{bound})$ 
 $\vdash \forall x: |(g \downarrow \text{oset})|. \uparrow(x \in_b \text{map}(\lambda z.z.1; \text{ps})) \Rightarrow \uparrow(k * x <_b k * \text{bound})$ 
|
BY (RepD THENM RW bool_to_propC 0 ...a)
|
8.  $x: |(g \downarrow \text{oset})|$ 
9.  $\uparrow(x \in_b \text{map}(\lambda z.z.1; \text{ps}))$ 
 $\vdash k * x < k * \text{bound}$ 
|
BY (BLemma 'grp_op_preserves_lt'
    THENM RW (RevLemmaC 'assert_of_grp_blt') 0
    THENM BHyp 7 ...)
*C omral_scale_sd_ordered_com
    The proof here needs some cleaning up.
    Probably, worth pulling out the second
    induction and generalizing it a bit.
*T omral_scale_sd_ordered 109.9 sec.
 $\vdash \forall g: \text{OCMon}. \forall r: \text{CRng}. \forall k: |g|. \forall v: |r|. \forall \text{ps}: (|g| \times |r|) \text{List}.$ 
 $|\ \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \text{ps})) \Rightarrow \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \langle k, v \rangle * \text{ps}))$ 
|
BY (CDToVarThen 'ps' (\i.New ['q'; 'qs'] (ListInd i))
    ...a)
|\
| 1.  $g: \text{OCMon}$ 
| 2.  $r: \text{CRng}$ 
| 3.  $k: |g|$ 
| 4.  $v: |r|$ 
| 5.  $\text{ps}: (|g| \times |r|) \text{List}$ 
|  $\vdash \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; [])) \Rightarrow \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \langle k, v \rangle * []))$ 
| |
1 BY AbReduce 0
| |
|  $\vdash \text{True} \Rightarrow \text{True}$ 
| |
1 BY Auto
\
| 1.  $g: \text{OCMon}$ 
| 2.  $r: \text{CRng}$ 
| 3.  $k: |g|$ 
| 4.  $v: |r|$ 
| 5.  $\text{ps}: (|g| \times |r|) \text{List}$ 
| 6.  $q: |g| \times |r|$ 
| 7.  $\text{qs}: (|g| \times |r|) \text{List}$ 
| 8.  $\uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \text{qs})) \Rightarrow \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \langle k, v \rangle * \text{qs}))$ 
 $\vdash \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; q::\text{qs})) \Rightarrow \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \langle k, v \rangle * (q::\text{qs})))$ 
|
BY New ['kq'; 'vq'] (OnVar 'q' D) THEN AbReduce 0
| THEN (D 0 THENM RW bool_to_propC (-1) THENM D (-1) ...a)
| THEN (SplitOnConclITE ...a)
|\
| 6.  $kq: |g|$ 
| 7.  $vq: |r|$ 
| 8.  $\text{qs}: (|g| \times |r|) \text{List}$ 
| 9.  $\uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \text{qs})) \Rightarrow \uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \langle k, v \rangle * \text{qs}))$ 
| 10.  $\uparrow \text{before}(kq; \text{map}(\lambda z.z.1; \text{qs}))$ 
| 11.  $\uparrow \text{sd\_ordered}(\text{map}(\lambda z.z.1; \text{qs}))$ 
| 12.  $v * vq = 0$ 

```

```

| ⊢ ↑sd_ordered(map(λz.z.1;<k,v>* qs))
| |
1 BY (BHyp 9 ...)
\
6. kq: |g|
7. vq: |r|
8. qs: (|g| × |r|) List
9. ↑sd_ordered(map(λz.z.1;qs)) ⇒ ↑sd_ordered(map(λz.z.1;<k,v>* qs))
10. ↑before(kq;map(λz.z.1;qs))
11. ↑sd_ordered(map(λz.z.1;qs))
12. ¬(v * vq = 0)
⊢ ↑sd_ordered(map(λz.z.1;<k * kq, v * vq>::(<k,v>* qs)))
|
BY % Can't use AbReduce as middle step because it destroys pattern
|   RevLemmaC to match against %
|
| (RWH (LemmaC 'sd_ordered_char') 0
|   THENM RWH map_cons_unrollC 0
|   THENM RWH mon_htfor_consC 0
|   THENM RWH (RevLemmaC 'sd_ordered_char') 0 ...a)
|\
| 13. w: |(g↓oset)|
| ⊢ (λz.z.1) <k * kq, v * vq> ∈ |(g↓oset)|
| |
1 BY (AbReduce 0 ...)
\
⊢ ↑((∀bw(:|(g↓oset)|) ∈ map(λz.z.1;<k,v>* qs). w <_b (λz.z.1) <k * kq, v * vq>)
| * sd_ordered(map(λz.z.1;<k,v>* qs)))
|
BY AbReduce 0 THENM (RW bool_to_propC 0 THENM D 0 ...a)
|\
| ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* qs). w <_b k * kq)
| |
1 BY (Assert [↑sd_ordered(map(λz.z.1;<kq, vq>::qs))]
| |   THENA (AbReduce 0 THEN RW bool_to_propC 0 THENM Auto) ...a)
| |   THEN (RWH (LemmaC 'sd_ordered_char') (-1) ...a)
| |
| 13. ↑(HTFor{<ℕ, ∧_b>} v::vs ∈ map(λz.z.1;<kq, vq>::qs). ∀bw(:|(g↓oset)|) ∈ vs. w <_b v)
| |
1 BY AbReduce (-1) THENM (RW bool_to_propC (-1) THENM RepD ...a)
| |
| 13. ↑(∀bw(:|g|) ∈ map(λz.z.1;qs). w <_b kq)
| 14. ↑(HTFor{<ℕ, ∧_b>} v::vs ∈ map(λz.z.1;qs). ∀bw(:|g|) ∈ vs. w <_b v)
| |
1 BY % The last few steps were pretty ugly. They should be cleaned up %
| |   OnHyps [14;12;11;10;9] Thin
| |   THEN (OnVar 'qs' (MoveDepHypsToConclFor ListInd) ...a)
| |\
| | 9. ↑(∀bw(:|g|) ∈ map(λz.z.1;[]). w <_b kq)
| | ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* []). w <_b k * kq)
| | |
1 2 BY (AbReduce 0 ...)
| |\
| | ⊢ map(λz.z.1;[]) ∈ |g| List
| | |
1 2 BY % Sigh. The tribulations of implicit polymorphism
| |   and a weak type inf routine %

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| | (AbReduce 0 ...)
| \
| 9. u: |g| × |r|
| 10. v1: (|g| × |r|) List
| 11. ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
|     ⇒ {↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* v1). w <b k * kq)}
| 12. ↑(∀bw(:|g|) ∈ map(λz.z.1;u::v1). w <b kq)
| ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* (u::v1)). w <b k * kq)
|
|
1 BY D 9 THEN AbReduce 0 THEN (SplitOnConclITE ...a)
| \
| | 9. u1: |g|
| | 10. u2: |r|
| | 11. v1: (|g| × |r|) List
| | 12. ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
| |     ⇒ {↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* v1). w <b k * kq)}
| | 13. ↑(∀bw(:|g|) ∈ map(λz.z.1;<u1, u2>::v1). w <b kq)
| | 14. v * u2 = 0
| | ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* v1). w <b k * kq)
| |
|
1 2 BY (BHyp 12 ...)
| |
| | ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
| |
|
1 2 BY (AbReduce 13 THENM RW bool_to_propC 13 ...)
| \
| 9. u1: |g|
| 10. u2: |r|
| 11. v1: (|g| × |r|) List
| 12. ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
|     ⇒ {↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* v1). w <b k * kq)}
| 13. ↑(∀bw(:|g|) ∈ map(λz.z.1;<u1, u2>::v1). w <b kq)
| 14. ¬(v * u2 = 0)
| ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k * u1, v * u2>::(<k,v>* v1)). w <b k * kq)
|
|
1 BY (AbReduce 0 THENM RW bool_to_propC 0
|   THENM D 0 ...a)
| \
| | ⊢ k * u1 <g↓oset k * kq
| |
|
1 2 BY (AbReduce 13 THENM RW bool_to_propC 13 THENM RepD ...a)
| |
| | 13. u1 <g↓oset kq
| | 14. ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
| | 15. ¬(v * u2 = 0)
| |
| |
1 2 BY OnCls [0;13] (Fold 'grp_lt')
|   THEN (BLemma 'grp_op_preserves_lt' ...)
| \
| | ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;<k,v>* v1). w <b k * kq)
| |
|
1 BY (BHyp 12 ...)
| |
| | ⊢ ↑(∀bw(:|g|) ∈ map(λz.z.1;v1). w <b kq)
| |
|
1 BY (AbReduce 13 THENM RW bool_to_propC 13 ...)
| \

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```

    ⊢ ↑sd_ordered(map(λz.z.1;<k,v>* qs))
    |
    BY (BHyp 9 ...)
*T omral_scale_non_zero_vals 43.7 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀v:|r|. ∀ps:(|g| × |r|) List.
|   ¬↑(0 ∈b map(λx.x.2;ps)) ⇒ ¬↑(0 ∈b map(λx.x.2;<k,v>* ps))
|
BY (CDToVarThen 'ps' (\i.New ['q';'qs'] (ListInd i))
|   ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| 3. k: |g|
| 4. v: |r|
| 5. ps: (|g| × |r|) List
| ⊢ ¬↑(0 ∈b map(λx.x.2;[])) ⇒ ¬↑(0 ∈b map(λx.x.2;<k,v>* []))
| |
1 BY (AbReduce 0 ...)
\
  1. g: OCMon
  2. r: CRng
  3. k: |g|
  4. v: |r|
  5. ps: (|g| × |r|) List
  6. q: |g| × |r|
  7. qs: (|g| × |r|) List
  8. ¬↑(0 ∈b map(λx.x.2;qs)) ⇒ ¬↑(0 ∈b map(λx.x.2;<k,v>* qs))
  ⊢ ¬↑(0 ∈b map(λx.x.2;q::qs)) ⇒ ¬↑(0 ∈b map(λx.x.2;<k,v>* (q::qs)))
  |
  BY (New ['kq';'vq'] (OnVar 'q' D)
  |   THENM AbReduce 0 THENM RW bool_to_propC 0
  |   THENM D 0 THENM RWH (LemmaC 'not_over_or') (-1) ...a)
  |
  6. kq: |g|
  7. vq: |r|
  8. qs: (|g| × |r|) List
  9. ¬↑(0 ∈b map(λx.x.2;qs)) ⇒ ¬↑(0 ∈b map(λx.x.2;<k,v>* qs))
  10. ¬(vq = 0) ∧ ¬↑(0 ∈b map(λx.x.2;qs))
  ⊢ ¬↑(0 ∈b map(λx.x.2;if (v * vq) =b 0 then <k,v>* qs else <k * kq, v * vq>::(<k,v>* qs) fi ))
  |
  BY (SplitOnConclITE THENM AbReduce 0 ...a)
  |\
  | 11. v * vq = 0
  | ⊢ ¬↑(0 ∈b map(λx.x.2;<k,v>* qs))
  | |
  1 BY (BHyp 9 ...)
  \
    11. ¬(v * vq = 0)
    ⊢ ¬↑(((v * vq) =b 0) ∨b(0 ∈b map(λx.x.2;<k,v>* qs)))
    |
    BY (RW bool_to_propC 0
        THENM ProveProp ...)
*T omral_scale_wf2 22.0 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀v:|r|. ∀ps:|omral(g;r)|. <k,v>* ps ∈ |omral(g;r)|
|
BY (RepD ...a)
|

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```

1. g: OCMon
2. r: CRng
3. k: |g|
4. v: |r|
5. ps: |omral(g;r)|
├ <k,v>* ps ∈ |omral(g;r)|
|
BY (OnCls [0;5] (Unfold 'omralist'
| THENM AddCarProperties 5
| THENM AbReduce 0 THENM MemTypeCD ...a)
| \
| 5. ps: |oal(g↓oset;r↓+gp)|
| 6. ↑sd_ordered(map(λx.x.1;ps)) ∧ ¬↑(e ∈b map(λx.x.2;ps))
| ─ <k,v>* ps ∈ (|g| × |r|) List
| |
1 BY Auto
| \
| 5. ps: |oal(g↓oset;r↓+gp)|
| 6. ↑sd_ordered(map(λx.x.1;ps))
| 7. ¬↑(e ∈b map(λx.x.2;ps))
| ─ ↑sd_ordered(map(λx.x.1;<k,v>* ps))
| |
1 BY (BLemma 'omral_scale_sd_ordered' ...)
| \
| 5. ps: |oal(g↓oset;r↓+gp)|
| 6. ↑sd_ordered(map(λx.x.1;ps))
| 7. ¬↑(e ∈b map(λx.x.2;ps))
| ─ ¬↑(0 ∈b map(λx.x.2;<k,v>* ps))
|
BY AbReduce 7 THEN (BLemma 'omral_scale_non_zero_vals' ...)
*T lookup_omral_scale_a 89.5 sec.
├ ∀g:OCMon. ∀r:CRng. ∀k,k':|g|. ∀v:|r|. ∀ps:|omral(g;r)|. (<k,v>* ps)[k * k'] = v * ps[k']
|
BY (RepeatMFor 5 (D 0) ...a)
|
1. g: OCMon
2. r: CRng
3. k: |g|
4. k': |g|
5. v: |r|
├ ∀ps:|omral(g;r)|. (<k,v>* ps)[k * k'] = v * ps[k']
|
BY (Unfold 'omralist' 0
| THEN BLemma 'oalist_ind' ...a)
| \
| ─ (<k,v>* [])[k * k'] = v * [][k']
| |
1 BY AbReduce 0
| |
| ─ 0 = v * 0
| |
1 BY (RW RngNormC 0 ...)
| \
| ─ ∀ps:|oal(g↓oset;r↓+gp)|
|   (<k,v>* ps)[k * k'] = v * ps[k']
|   ⇒ (∀x:|(g↓oset)|. ∀y:|r↓+gp|.
|     ↑before(x,map(λx.x.1;ps))

```

```

|           ⇒ ¬(y = e)
|           ⇒ (<k,v>* (<x, y>::ps))[k * k'] = v * (<x, y>::ps)[k']
|
BY AbReduceIf (\e t.not is_term 'oalist' (subterm t 1)) 0
| THEN Fold 'omralist' 0 THEN RenameBVars ['x','kp';'y','kv'] 0
| THEN (RepD ...a)
|
6. ps: |omral(g;r)|
7. (<k,v>* ps)[k * k'] = v * ps[k']
8. kp: |g|
9. kv: |r|
10. ↑before(kp;map(λkp.kp.1;ps))
11. ¬(kv = 0)
├ if (v * kv) =b 0 then <k,v>* ps else <k * kp, v * kv>::(<k,v>* ps) fi [k * k']
| = v * if kp =b k' then kv else ps[k'] fi
|
BY (SplitOnConclITEs ...a)
| \
| 12. v * kv = 0
| 13. kp = k'
| ─ (<k,v>* ps)[k * k'] = v * kv
| |
1 BY (RewriteWith [12;13] ['rng_lookup_before_start'] 0 ...a)
| | \
| | ─ ↑before(k * k';map(λz.z.1;<k,v>* ps))
| | |
1 2 BY (BLemma 'rng_before_all_imp_before'
| | | THENM BLemma 'omral_scale_dom_bound'
| | | THENM BLemma 'rng_before_imp_before_all' ...a)
| | |
| | ─ ↑before(k';map(λz.z.1;ps))
| | |
1 2 BY (RWH (RevHypC 13) 0 ...)
| | \
| | ─ 0 = 0
| | |
1 BY Auto
| | \
| | 12. v * kv = 0
| | 13. ¬(kp = k')
| | ─ (<k,v>* ps)[k * k'] = v * ps[k']
| | |
1 BY Trivial
| | \
| | 12. ¬(v * kv = 0)
| | 13. kp = k'
| | ─ (<k * kp, v * kv>::(<k,v>* ps))[k * k'] = v * kv
| | |
1 BY (Reduce 0 THEN SplitOnConclITE ...a)
| | \
| | 14. k * kp = k * k'
| | ─ v * kv = v * kv
| | |
1 2 BY Auto
| | \
| | 14. ¬(k * kp = k * k')
| | ─ (<k,v>* ps)[k * k'] = v * kv

```



```

| |
1 BY (D 14 THENM RWH (HypC 13) 0 ...)
| \
| 12.  $\neg(v * kv = 0)$ 
| 13.  $\neg(kp = k')$ 
|  $\vdash (\langle k * kp, v * kv \rangle :: (\langle k, v \rangle * ps)) [k * k'] = v * ps[k']$ 
| |
| BY (Reduce 0 THEN SplitOnConclITE ...a)
| \
| | 14.  $k * kp = k * k'$ 
| |  $\vdash v * kv = v * ps[k']$ 
| | |
| 1 BY % Contradiction %
| | (FLemma 'ocmon_cancel' [14] ...)
| \
| 14.  $\neg(k * kp = k * k')$ 
|  $\vdash (\langle k, v \rangle * ps) [k * k'] = v * ps[k']$ 
| |
| BY Trivial
*T lookup_omral_scale_b 98.9 sec.
 $\vdash \forall g:OCMon. \forall r:CRng. \forall k, k':|g|. \forall v:|r|. \forall ps:(|g| \times |r|) List.$ 
|  $\neg(\exists d:|g|. \uparrow(d \in_b \text{dom}(ps)) \wedge k * d = k') \Rightarrow (\langle k, v \rangle * ps) [k'] = 0$ 
| |
| BY (CDToVarThen 'ps' ListIndA
| THENM Reduce 0 ...a)
| \
| 1. g: OCMon
| 2. r: CRng
| 3. k: |g|
| 4. k': |g|
| 5. v: |r|
|  $\vdash \neg(\exists d:|g|. \text{False} \wedge k * d = k') \Rightarrow 0 = 0$ 
| |
| 1 BY Auto
| \
| 1. g: OCMon
| 2. r: CRng
| 3. k: |g| | |
| 4. k': |g|
| 5. v: |r|
| 6. p: |g|  $\times$  |r|
| 7. ps: (|g|  $\times$  |r|) List
| 8.  $\neg(\exists d:|g|. \uparrow(d \in_b \text{dom}(ps)) \wedge k * d = k') \Rightarrow (\langle k, v \rangle * ps) [k'] = 0$ 
|  $\vdash \neg(\exists d:|g|. \uparrow(d \in_b \text{dom}(p::ps)) \wedge k * d = k') \Rightarrow (\langle k, v \rangle * (p::ps)) [k'] = 0$ 
| |
| BY New ['kp'; 'vp'] (D 6) THEN Reduce 0
| THEN (D 0 THENM SplitOnConclITE ...a)
| \
| 6. kp: |g|
| 7. vp: |r|
| 8. ps: (|g|  $\times$  |r|) List
| 9.  $\neg(\exists d:|g|. \uparrow(d \in_b \text{dom}(ps)) \wedge k * d = k') \Rightarrow (\langle k, v \rangle * ps) [k'] = 0$ 
| 10.  $\neg(\exists d:|g|. \uparrow((kp =_b d) \vee_b (d \in_b \text{dom}(ps))) \wedge k * d = k')$ 
| 11.  $v * vp = 0$ 
|  $\vdash (\langle k, v \rangle * ps) [k'] = 0$ 
| |
| 1 BY (BHyp 9 THENM D 0 THENM D 10 THENM ExRepD ...a)

```

```

| |
| 10. v * vp = 0
| 11. d: |g|
| 12. ↑(d ∈b dom(ps))
| 13. k * d = k'
| ⊢ ∃d:|g|. ↑((kp =b d) ∨b(d ∈b dom(ps))) ∧ k * d = k'
| |
1 BY (With 「d」 (D 0) ...)
| |
| ⊢ ↑((kp =b d) ∨b(d ∈b dom(ps)))
| |
1 BY (RW bool_to_propC 0 THENM Sel 2 (D 0) ...)
\
6. kp: |g|
7. vp: |r|
8. ps: (|g| × |r|) List
9. ¬(∃d:|g|. ↑(d ∈b dom(ps)) ∧ k * d = k') ⇒ (<k,v>* ps)[k'] = 0
10. ¬(∃d:|g|. ↑((kp =b d) ∨b(d ∈b dom(ps))) ∧ k * d = k')
11. ¬(v * vp = 0)
⊢ (<k * kp, v * vp>::(<k,v>* ps))[k'] = 0
|
BY (Reduce 0 THEN SplitOnConclITE ...a)
|\
| 12. k * kp = k'
| ⊢ v * vp = 0
| |
1 BY % hyp 12 contradicts hyp 10 %
| | (D 10 THENM With 「kp」 (D 0) ...)
| |
| 10. ¬(v * vp = 0)
| 11. k * kp = k'
| ⊢ ↑((kp =b kp) ∨b(kp ∈b dom(ps)))
| |
1 BY (RW bool_to_propC 0 THENM Sel 1 (D 0) ...)
\
12. ¬(k * kp = k')
⊢ (<k,v>* ps)[k'] = 0
|
BY (BHyp 9 THENM D 0 THENM D 10 THENM ExRepD ...)
|
10. ¬(v * vp = 0)
11. ¬(k * kp = k')
12. d: |g|
13. ↑(d ∈b dom(ps))
14. k * d = k'
⊢ ∃d:|g|. ↑((kp =b d) ∨b(d ∈b dom(ps))) ∧ k * d = k'
|
BY (With 「d」 (D 0) ...)
|
⊢ ↑((kp =b d) ∨b(d ∈b dom(ps)))
|
BY (RW bool_to_propC 0 THENM Sel 2 (D 0) ...)
*T lookup_omral_scale_c 75.7 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀z,k:|g|. ∀v:|r|. ∀ps:|omral(g;r)|.
| (<k,v>* ps)[z] = msFor{r|+gp} y ∈ dom(ps). when (k * y) =b z. v * ps[y]
|
BY (RepD ...a)

```

```

|
1. g: OCMon
2. r: CRng
3. z: |g|
4. k: |g|
5. v: |r|
6. ps: |omral(g;r)|
⊢ (<k,v>* ps)[z] = msFor{r↓+gp} y ∈ dom(ps). when (k * y) =b z. v * ps[y]
|
BY % This predicate is constructively decidable, but no need to prove so,
|   since here there is no constructive content. %
|   (AddXM 1 THENM
|     Decide [∃d:|g|. ↑(d ∈b dom(ps)) ∧ k * d = z] ...a)
| \
| 1. XM{i'}
| 2. g: OCMon
| 3. r: CRng
| 4. z: |g|
| 5. k: |g|
| 6. v: |r|
| 7. ps: |omral(g;r)|
| 8. ∃d:|g|. ↑(d ∈b dom(ps)) ∧ k * d = z
| |
1 BY ExRepD
| |
| 8. d: |g|
| 9. ↑(d ∈b dom(ps))
| 10. k * d = z
| |
1 BY Unfold 'rng_when' 0 THEN
| | (RW (SweepUpC
| | (RevHypC 10
| | ORELSEC LemmaC 'lookup_omral_scale_a'
| | ORELSEC LemmaWithC ['u',↑d] 'fset_for_when_unique')) 0 ...a)
| | \
| | ⊢ ↑((k * d) =b (k * d))
| | |
1 2 BY (RW bool_to_propC 0 ...)
| | \
| | 11. v1: |(g↓oset)|
| | 12. ↑((k * v1) =b (k * d))
| | 13. ↑(v1 ∈b dom(ps))
| | ⊢ v1 = d
| | |
1 2 BY (RW bool_to_propC 12
| | THENM FLemma 'ocmon_cancel' [12]
| | THENM Reduce 0 ...)
| | \
| | ⊢ v * ps[d] = v * ps[d]
| | |
1 BY Auto
| \
| 1. XM{i'}
| 2. g: OCMon
| 3. r: CRng
| 4. z: |g|
| 5. k: |g|

```

```

6. v: |r|
7. ps: |omral(g;r)|
8.  $\neg(\exists d:|g|. \uparrow(d \in_b \text{dom}(ps)) \wedge k * d = z)$ 
|
BY (Unfold 'rng_when' 0 THEN
| RewriteWith [] 'lookup_omral_scale_b mset_for_when_none' 0 ...a)
|\
| 9. x: |(g\oset)|
| 10.  $\uparrow(x \in_b \text{dom}(ps))$ 
|  $\vdash \neg\uparrow((k * x) =_b z)$ 
| |
1 BY (D 0 THENM RW bool_to_propC (-1)
| THENM D 8 THENM With 'x' (D 0) ...)
\
|  $\vdash 0 = e$ 
|
BY (Reduce 0 ...)
*T lookup_omral_scale_d 0.2 sec.
 $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall z,k:|g|. \forall v:|r|. \forall ps:|omral(g;r)|.$ 
|  $(\langle k,v \rangle * ps)[z] = (\sum y \in \text{dom}(ps). \text{when}(k * y) =_b z. v * ps[y])$ 
|
BY Unfold 'rng_mssum' 0
| THENM AssertLemma 'lookup_omral_scale_c' []
|
1.  $\forall g:\text{OCMon}. \forall r:\text{CRng}. \forall z,k:|g|. \forall v:|r|. \forall ps:|omral(g;r)|.$ 
|  $(\langle k,v \rangle * ps)[z] = \text{msFor}\{r\}+\text{gp}\} y \in \text{dom}(ps). \text{when}(k * y) =_b z. v * ps[y]$ 
 $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall z,k:|g|. \forall v:|r|. \forall ps:|omral(g;r)|.$ 
|  $(\langle k,v \rangle * ps)[z] = \text{msFor}\{r\}+\text{gp}\} y \in \text{dom}(ps). \text{when}(k * y) =_b z. v * ps[y]$ 
|
BY Unfold 'oset_of_ocmon' 1 THEN Trivial
*C omral_times_com
=====
OMRAL TIMES FUNCTION
=====
*D omral_times_df
Parens ::Prec(inop)::
<ps:ps:L> **<g:g:L>,<r:r:L> <qs:qs:L>
== omral_times{<g>; <r>; <ps>; <qs>}
Parens ::Prec(inop)::
<ps:ps:L> ** <qs:qs:L>
== omral_times{<g>; <r>; <ps>; <qs>}
*M omral_times_ml
ps ** qs==r case ps of [] => [] | p::ps' => <p.1,p.2>* qs ++ (ps' ** qs) esac
*M omral_times_eval
let omral_times_nilC =
FwdMacroC 'omral_times_nilC'
(RecEvalC 'omral_times' '[] ** qs' );
let omral_times_cons_prC =
FwdMacroC 'omral_times_cons_prC'
(RecEvalC 'omral_times' '(<k, v>::ps) ** qs' );
;;
add_AbReduce_conv 'omral_times'
(omral_times_nilC ORELSEC omral_times_cons_prC)
;;
*T omral_times_wf 4.2 sec.
 $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall ps,qs:(|g| \times |r|) \text{List}. ps ** qs \in (|g| \times |r|) \text{List}$ 
|

```

```

BY (RepD THENM New ['p';'ps\''] (OnVar 'ps' ListInd)
| THEN AbReduce 0
| ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| 3. ps: (|g| × |r|) List
| 4. qs: (|g| × |r|) List
| ⊢ [] ∈ (|g| × |r|) List
| |
1 BY Auto
\
| 1. g: OCMon
| 2. r: CRng
| 3. ps: (|g| × |r|) List
| 4. qs: (|g| × |r|) List
| 5. p: |g| × |r|
| 6. ps': (|g| × |r|) List
| 7. ps' ** qs ∈ (|g| × |r|) List
| ⊢ (p::ps') ** qs ∈ (|g| × |r|) List
|
BY New ['k';'v'] (D 5) THEN AbReduce 0
|
| 5. k: |g|
| 6. v: |r|
| 7. ps': (|g| × |r|) List
| 8. ps' ** qs ∈ (|g| × |r|) List
| ⊢ <k,v>* qs ++ (ps' ** qs) ∈ (|g| × |r|) List
|
BY MemCD THEN Auto
*T omral_times_sd_ordered 10.7 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:(|g| × |r|) List.
| ↑sd_ordered(map(λz.z.1;qs)) ⇒ ↑sd_ordered(map(λz.z.1;ps ** qs))
|
BY (RepD THENM OnVar 'ps' ListIndA ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| 3. qs: (|g| × |r|) List
| 4. ↑sd_ordered(map(λz.z.1;qs))
| ⊢ ↑sd_ordered(map(λz.z.1;[] ** qs))
| |
1 BY (Reduce 0 ...)
\
| 1. g: OCMon
| 2. r: CRng
| 3. qs: (|g| × |r|) List
| 4. ↑sd_ordered(map(λz.z.1;qs))
| 5. p: |g| × |r|
| 6. ps: (|g| × |r|) List
| 7. ↑sd_ordered(map(λz.z.1;ps ** qs))
| ⊢ ↑sd_ordered(map(λz.z.1;(p::ps) ** qs))
|
BY New ['kp';'vp'] (OnVar 'p' D)
| THEN Reduce 0 THEN (RepD ...a)
|
| 5. kp: |g|

```

```

6. vp: |r|
7. ps: (|g| × |r|) List
8. ↑sd_ordered(map(λz.z.1;ps ** qs))
├ ↑sd_ordered(map(λz.z.1;<kp,vp>* qs ++ (ps ** qs)))
|
BY (Backchain ‘‘omral_plus_sd_ordered omral_scale_sd_ordered‘‘ ...)
*T omral_times_non_zero_vals 11.2 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ps,qs:(|g| × |r|) List.
|   ¬↑(0 ∈b map(λx.x.2;qs)) ⇒ ¬↑(0 ∈b map(λx.x.2;ps ** qs))
|
BY (RepD THENM OnVar ‘ps’ ListIndA ...a)
| \
| 1. g: OCMon
| 2. r: CRng
| 3. qs: (|g| × |r|) List
| 4. ¬↑(0 ∈b map(λx.x.2;qs))
| 5. ─┬─ ¬↑(0 ∈b map(λx.x.2;[] ** qs))
|     |
| 1 BY (Reduce 0 THEN D 0 ...)
|   \
|   1. g: OCMon
|   2. r: CRng
|   3. qs: (|g| × |r|) List
|   4. ¬↑(0 ∈b map(λx.x.2;qs))
|   5. p: |g| × |r|
|   6. ps:(|g| × |r|) List
|   7. ¬↑(0 ∈b map(λx.x.2;ps ** qs))
|   8. ─┬─ ¬↑(0 ∈b map(λx.x.2;(p::ps) ** qs))
|     |
|   BY New [‘kp’;‘vp’] (OnVar ‘p’ D)
|     | THEN Reduce 0
|     |
|     5. kp: |g|
|     6. vp: |r|
|     7. ps: (|g| × |r|) List
|     8. ¬↑(0 ∈b map(λx.x.2;ps ** qs))
|     9. ─┬─ ¬↑(0 ∈b map(λx.x.2;<kp,vp>* qs ++ (ps ** qs)))
|       |
|       BY (Backchain ‘‘
|         omral_scale_non_zero_vals
|         omral_plus_non_zero_vals‘‘ ...)
*T omral_times_wf2 19.4 sec.
├ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. ps ** qs ∈ |omral(g;r)|
|
BY (RepD ...a)
|
| 1. g: OCMon
| 2. r: CRng
| 3. ps: |omral(g;r)|
| 4. qs: |omral(g;r)|
├ ps ** qs ∈ |omral(g;r)|
|
BY OnHyps [4;3] AddCarProperties
| THEN Reduce 0 THEN (MemTypeCD ...)
| \
| 4. ↑sd_ordered(map(λx.x.1;ps))
| 5. ¬↑(0 ∈b map(λx.x.2;ps))

```

```

| 6. qs: |omral(g;r)|
| 7. ↑sd_ordered(map(λx.x.1;qs))
| 8. ¬↑(0 ∈b map(λx.x.2;qs))
| ⊢ ↑sd_ordered(map(λx.x.1;ps ** qs))
| |
1 BY (BLemma 'omral_times_sd_ordered' ...)
\
  4. ↑sd_ordered(map(λx.x.1;ps))
  5. ¬↑(0 ∈b map(λx.x.2;ps))
  6. qs: |omral(g;r)|
  7. ↑sd_ordered(map(λx.x.1;qs))
  8. ¬↑(0 ∈b map(λx.x.2;qs))
  ⊢ ¬↑(0 ∈b map(λx.x.2;ps ** qs))
  |
  BY (BLemma 'omral_times_non_zero_vals' ...)
*T lookup_omral_times 98.9 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. ∀z:|g|.
|   (ps ** qs)[z]
|   = msFor{r↓+gp} x ∈ dom(ps). msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. ps[x] * qs[y]
|
BY (RepD THENM OnVar 'ps' MoveToConcl
|   THENM BLemma 'omralist_ind_a' ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| 3. qs: |omral(g;r)|
| 4. z: |g|
| ⊢ ([] ** qs)[z]
| | = msFor{r↓+gp} x ∈ dom([]). msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. [][x] * qs[y]
| |
1 BY Reduce 0
| |
| ⊢ 0 = 0
| |
1 BY Auto
\
  1. g: OCMon
  2. r: CRng
  3. qs: |omral(g;r)|
  4. z: |g|
  ⊢ ∀ps:|omral(g;r)|
  |   (ps ** qs)[z]
  |   = msFor{r↓+gp} x ∈ dom(ps). msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. ps[x] * qs[y]
  |   ⇒ (∀x:|g|. ∀y:|r|.
  |       ↑before(x,map(λx.x.1;ps))
  |       ⇒ ¬(y = 0)
  |       ⇒ ((⟨x, y⟩::ps) ** qs)[z]
  |           = msFor{r↓+gp} x@0 ∈ dom(⟨x, y⟩::ps)
  |             msFor{r↓+gp} y@0 ∈ dom(qs)
  |             when (x@0 * y@0) =b z. (⟨x, y⟩::ps)[x@0] * qs[y@0])
  |
BY (RenameBVars ['x','kp';'x@0','x';'y','vp';'y@0','y'] 0
|   THENM RepD ...a)
|
5. ps: |omral(g;r)|
6. (ps ** qs)[z]
   = msFor{r↓+gp} kp ∈ dom(ps)

```

```

      msFor{r↓+gp} vp ∈ dom(qs). when (kp * vp) =b z. ps[kp] * qs[vp]
7. kp: |g|
8. vp: |r|
9. ↑before(kp;map(λkp.kp.1;ps))
10. ¬(vp = 0)
├ ((<kp, vp>::ps) ** qs)[z]
| = msFor{r↓+gp} x ∈ dom(<kp, vp>::ps)
|   msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. (<kp, vp>::ps)[x] * qs[y]
|
BY Reduce 0
|
├ (<kp, vp>* qs ++ (ps ** qs))[z]
| = (msFor{r↓+gp} y ∈ dom(qs)
|   when (kp * y) =b z. if kp =b kp then vp else ps[kp] fi * qs[y])
| +r (msFor{r↓+gp} x ∈ dom(ps)
|   msFor{r↓+gp} y ∈ dom(qs)
|   when (x * y) =b z. if kp =b x then vp else ps[x] fi * qs[y])
|
BY % The ifthenelse reduction should be done
| more intelligently. %
| (RWN 1 (LemmaC 'ite_rw_true') 0
| THENM RWH (LemmaC 'ite_rw_false') 0
| THENM RWH (LemmaC 'lookup_omral_plus') 0 ...a)
| \
| 11. x: |(g↓oset)|
| 12. ↑(x ∈b dom(qs))
| ─ ↑(kp =b kp)
| |
1 BY (RW bool_to_propC 0 ...)
| \
| 11. x: |(g↓oset)|
| 12. ↑(x ∈b dom(ps))
| 13. x1: |(g↓oset)|
| 14. ↑(x1 ∈b dom(qs))
| ─ ¬↑(kp =b x)
| |
1 BY (FLemma 'rng_before_imp_before_all' [9] ...a)
| |
| 15. ↑(∀bx(:|g|) ∈ map(λkp.kp.1;ps). x <b kp)
| |
1 BY % Aaargh! %
| | Assert ↑(∀bx(:|(g↓oset)|) ∈ map(λkp.kp.1;ps). x <b kp)1
| | THENA (Reduce 0 THEN Auto)
| | THEN (RWH (LemmaC 'ball_char') (-1)
| |   THENM RW bool_to_propC (-1) ...a)
| |
| 16. ∀x:|(g↓oset)|. ↑(x ∈b map(λkp.kp.1;ps)) ⇒ x < kp
| |
1 BY (RepUnfolds 'mset_mem omral_dom oal_dom mk_mset' 12
|   THEN RW bool_to_propC 0
|   THENM InstHyp [↑x1] 16
|   THENM RelRST ...)
| \
├ (<kp, vp>* qs)[z] +r (ps ** qs)[z]
| = (msFor{r↓+gp} y ∈ dom(qs). when (kp * y) =b z. vp * qs[y])
| +r (msFor{r↓+gp} x ∈ dom(ps)
|   msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. ps[x] * qs[y])

```



```

|
BY EqCD
|\
| ⊢ +r = +r
| |
1 BY Auto
|\
| ⊢ (<kp, vp>* qs)[z] = msFor{r↓+gp} y ∈ dom(qs). when (kp * y) =b z. vp * qs[y]
| |
1 BY (BLemma 'lookup_omral_scale_c' ...)
\
| ⊢ (ps ** qs)[z]
|   = msFor{r↓+gp} x ∈ dom(ps). msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b z. ps[x] * qs[y]
|
BY Trivial
*T lookup_omral_times_a 0.0 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. ∀z:|g|.
|   (ps ** qs)[z] = (∑x ∈ dom(ps). ∑y ∈ dom(qs). when (x * y) =b z. ps[x] * qs[y])
|
BY Unfold 'rng_mssum' 0 THEN Lemma 'lookup_omral_times'
*T mset_on_grp_eq 1.8 sec.
⊢ ∀g:OCMon. MSet{g↓set} = MSet{g↓oset}
|
BY (D 0 ...a)
|
1. g: OCMon
⊢ MSet{g↓set} = MSet{g↓oset}
|
BY (Assert 「MSet{g↓set} ∈ U」 ...a)
|
2. MSet{g↓set} ∈ U
|
BY Repeat (D 1)
  THENM OnCls [0;-1] (Eval 'mset')
  THENM Trivial
*T mset_inc 1.6 sec.
⊢ ∀g:OCMon. MSet{g↓set} ⊆ MSet{g↓oset}
|
BY (Unfold 'subtype' 0
|   THENM RepD ...a)
|
1. g: OCMon
2. x: MSet{g↓set}
⊢ x ∈ MSet{g↓oset}
|
BY OnCls [0;2] (Unfold 'mset')
|
2. x: as,bs:(|g↓set| List)//(as ≡(|g↓set|) bs)
⊢ x ∈ as,bs:(|g↓oset| List)//(as ≡(|g↓oset|) bs)
|
BY Repeat (D 1)
|
1. car: U
2. g1: eq:(car → car → B)
   × le:(car → car → B)
   × op:(car → car → car)
   × id:car

```

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      × (car → car)
3. IsMonoid(|<car, g1>|;*;e) ∧ IsEqFun(|<car, g1>|;=b)
4. Comm(|<car, g1>|;* )
5. Linorder(|<car, g1>|;x,y.↑(x ≤b y))
   ∧ Cancel(|<car, g1>|;|<car, g1>|;* )
   ∧ (∀z:|<car, g1>|. monot(|<car, g1>|;x,y.↑(x ≤b y);λw.z * w))
6. x: as,bs:(|<car, g1>↓set)| List)//(as ≡(|<car, g1>↓set)|) bs)
⊢ x ∈ as,bs:(|<car, g1>↓set)| List)//(as ≡(|<car, g1>↓set)|) bs)
|
BY OnCls [0;-1] Reduce
|
6. x: as,bs:(car List)//(as ≡(car) bs)
⊢ x ∈ as,bs:(car List)//(as ≡(car) bs)
|
BY Trivial
*T mset_inc_a 3.9 sec.
⊢ ∀g:OCMon. MSet{g↓set} ⊆ MSet{g↓set}
|
BY (Unfold 'subtype' 0
| THENM RepD ...a)
|
1. g: OCMon
2. x: MSet{g↓set}
⊢ x ∈ MSet{g↓set}
|
BY OnCls [0;2] (Unfold 'mset')
|
2. x: as,bs:(|g↓set)| List)//(as ≡(|g↓set)|) bs)
⊢ x ∈ as,bs:(|g↓set)| List)//(as ≡(|g↓set)|) bs)
|
BY Repeat (D 1)
|
1. car: U
2. g1: eq:(car → car → B)
   × le:(car → car → B)
   × op:(car → car → car)
   × id:car
   × (car → car)
3. IsMonoid(|<car, g1>|;*;e) ∧ IsEqFun(|<car, g1>|;=b)
4. Comm(|<car, g1>|;* )
5. Linorder(|<car, g1>|;x,y.↑(x ≤b y))
   ∧ Cancel(|<car, g1>|;|<car, g1>|;* )
   ∧ (∀z:|<car, g1>|. monot(|<car, g1>|;x,y.↑(x ≤b y);λw.z * w))
6. x: as,bs:(|<car, g1>↓set)| List)//(as ≡(|<car, g1>↓set)|) bs)
⊢ x ∈ as,bs:(|<car, g1>↓set)| List)//(as ≡(|<car, g1>↓set)|) bs)
|
BY OnCls [0;-1] Reduce
|
6. x: as,bs:(car List)//(as ≡(car) bs)
⊢ x ∈ as,bs:(car List)//(as ≡(car) bs)
|
BY Trivial
*T omral_times_dom 292.1 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps,qs:|omral(g;r)|. ↑(dom(ps ** qs) ⊆b dom(ps) × dom(qs))
|
BY (RepD ...a)
|

```

```

1. g: OCMon
2. r: CRng
3. ps: |omral(g;r)|
4. qs: |omral(g;r)|
 $\vdash \uparrow(\text{dom}(ps ** qs) \subseteq_b \text{dom}(ps) \times \text{dom}(qs))$ 
|
BY (BLemma 'mem_bsubset' THENM RepD ...a)
|
5. x: |(g↓set)|
6.  $\uparrow(x \in_b \text{dom}(ps ** qs))$ 
 $\vdash \uparrow(x \in_b \text{dom}(ps) \times \text{dom}(qs))$ 
|
BY (OnVar 'ps' MoveToConcl
| THEN BLemma 'omralist_ind_a'
| THENM RepD ...a)
| \
| 3. qs: |omral(g;r)|
| 4. x: |(g↓set)|
| 5.  $\uparrow(x \in_b \text{dom}([] ** qs))$ 
|  $\vdash \uparrow(x \in_b \text{dom}([]) \times \text{dom}(qs))$ 
| |
1 BY Reduce 5
| |
| 5.  $\uparrow(x \in_b 0\{g\downarrow\text{oset}\})$ 
| |
1 BY (RWH (LemmaC 'mset_mem_char') 5 ...a)
| THEN Eval ''oset_of_ocmon'' 5 THEN Trivial
\
3. qs: |omral(g;r)|
4. x: |(g↓set)|
5. ps: |omral(g;r)|
6.  $\uparrow(x \in_b \text{dom}(ps ** qs)) \Rightarrow \uparrow(x \in_b \text{dom}(ps) \times \text{dom}(qs))$ 
7. x1: |g|
8. y: |r|
9.  $\uparrow\text{before}(x1;\text{map}(\lambda x.x.1;ps))$ 
10.  $\neg(y = 0)$ 
11.  $\uparrow(x \in_b \text{dom}(\langle x1, y \rangle :: ps) ** qs)$ 
 $\vdash \uparrow(x \in_b \text{dom}(\langle x1, y \rangle :: ps) \times \text{dom}(qs))$ 
|
BY Reduce 11
|
11.  $\uparrow(x \in_b \text{dom}(\langle x1, y \rangle * qs ++ (ps ** qs)))$ 
|
BY % A good example of monotone reasoning %
|
| (RWH (LemmaC 'omral_plus_dom') 11 ...a)
|
11.  $\uparrow(x \in_b \text{dom}(\langle x1, y \rangle * qs) \cup \text{dom}(ps ** qs))$ 
|
BY (Fold 'oset_of_ocmon' 11
| THENM RWH (LemmaC 'fset_mem_union') 11
| THENM RW bool_to_propC 11
| THENM D 11 ...a)
| \
| 11.  $\uparrow(x \in_b \text{dom}(\langle x1, y \rangle * qs))$ 
| |
1 BY (RWH (LemmaC 'omral_dom_scale') 11 ...a)

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| |
| 11.  $\uparrow(x \in_b \text{fs-map}(\lambda k'.k' * x1, \text{dom}(qs)))$ 
| |
1 BY Unfold 'fset_map' 11
| | THENM (RWH (LemmaC 'fset_of_mset_mem') 11 ...a)
| |
| 11.  $\uparrow(x \in_b \text{mmap}\{g\downarrow\text{oset}, g\downarrow\text{oset}\}(\lambda k'.k' * x1; \text{dom}(qs)))$ 
| |
1 BY (RWH (LemmaC 'mset_mem_char') 11 ...a)
| |
| 11.  $\uparrow(\exists_b\{g\downarrow\text{oset}\} y \in \text{mmap}\{g\downarrow\text{oset}, g\downarrow\text{oset}\}(\lambda k'.k' * x1; \text{dom}(qs)). y =_b x)$ 
| |
1 BY (FLemma 'bmsexists_char_a' [11] ...a)
| |
| 12.  $\downarrow(\exists y:|(g\downarrow\text{oset})|. \uparrow(y \in_b \text{mmap}\{g\downarrow\text{oset}, g\downarrow\text{oset}\}(\lambda k'.k' * x1; \text{dom}(qs))) \wedge \uparrow(y =_b x))$ 
| |
1 BY (D 12 THENM CUnhide
| | THENM ExRepD ...a)
| |
| 12.  $y1: |(g\downarrow\text{oset})|$ 
| 13.  $\uparrow(y1 \in_b \text{mmap}\{g\downarrow\text{oset}, g\downarrow\text{oset}\}(\lambda k'.k' * x1; \text{dom}(qs)))$ 
| 14.  $\uparrow(y1 =_b x)$ 
| |
1 BY (RW bool_to_propC 14 ...a)
| |
| 14.  $y1 = x$ 
| |
1 BY (RWH (LemmaC 'mset_mem_char') 0 ...a)
| |
|  $\vdash \uparrow(\exists_b\{g\downarrow\text{set}\} y \in \text{dom}(\langle x1, y \rangle :: ps) \times \text{dom}(qs). y =_b x)$ 
| |
1 BY (Reduce 0 THEN RenameBVars ['y', 'y2'] 0
| | THENM RW (SweepDnC (RevLemmaC 'bmsexists_char_rw'
| | ORELSEC LemmaC 'assert_of_mon_eq')) 0
| | THENM Reduce 0 ...a)
| |
|  $\vdash \exists y2:|g|. \uparrow(y2 \in_b (\text{mset\_inj}\{g\downarrow\text{oset}\}(x1) + \text{dom}(ps)) \times \text{dom}(qs)) \wedge y2 = x$ 
| |
1 BY (Reduce 14 THEN Inst ['y1'] 0 ...)
| |
| 14.  $y1 = x$ 
|  $\vdash \uparrow(y1 \in_b (\text{mset\_inj}\{g\downarrow\text{oset}\}(x1) + \text{dom}(ps)) \times \text{dom}(qs))$ 
| |
1 BY Unfold 'mset_prod' 0 THEN Fold 'oset_of_ocmon' 0
| | THEN Reduce 0
| |
|  $\vdash \uparrow(y1$ 
| |  $\in_b (\text{msFor}\{\langle \text{MSet}\{g\downarrow\text{oset}\}, \cup, 0 \rangle\} v \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{oset}\}(x1 * v))$ 
| |  $\cup (\text{msFor}\{\langle \text{MSet}\{g\downarrow\text{oset}\}, \cup, 0 \rangle\} u \in \text{dom}(ps)$ 
| |  $\text{msFor}\{\langle \text{MSet}\{g\downarrow\text{oset}\}, \cup, 0 \rangle\} v \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{oset}\}(u * v))$ 
| |
1 BY Unfold 'oset_of_ocmon' 0 THEN Fold 'mset_prod' 0
| |
|  $\vdash \uparrow(y1$ 
| |  $\in_b (\text{msFor}\{\langle \text{MSet}\{g\downarrow\text{set}\}, \cup, 0 \rangle\} v \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{set}\}(x1 * v)) \cup (\text{dom}(ps) \times \text{dom}(qs)))$ 
| |
1 BY (RWH (LemmaC 'fset_mem_union') 0

```

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| | THENM RW bool_to_propC 0 THENM Sel 1 (D 0) ...a)
| |
| |  $\uparrow (y1 \in_b \text{msFor}\{\langle \text{MSet}\{g\downarrow\text{set}\}, \cup, 0 \rangle\} v \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{set}\}(x1 * v))$ 
| |
1 BY Unfold 'mset_map' 13 THEN Reduce 13
| |
| | 13.  $\uparrow (y1 \in_b \text{msFor}\{\text{mset\_mon}\{g\downarrow\text{oset}\}\} x \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{oset}\}(x * x1))$ 
| |
1 BY (Unfold 'oset_of_ocmon' 13
| | THEN OnMCls [0;13] (\i.
| |   RWH (LambdaC [ $\lambda z. y1 \in_b z$ ]) i
| |   THENM RWH (LemmaWithC ['n', [ $\langle \mathbb{B}, \vee_b \rangle$ ]] 'dist_hom_over_mset_for') i
| |   THENM Reduce i) ...a)
| | \
| | 13.  $\uparrow (y1 \in_b \text{msFor}\{\text{mset\_mon}\{g\downarrow\text{set}\}\} x \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{set}\}(x * x1))$ 
| |  $\uparrow (\lambda z. y1 \in_b z) \in \text{MonHom}(\langle \text{MSet}\{g\downarrow\text{set}\}, \cup, 0 \rangle, \langle \mathbb{B}, \vee_b \rangle)$ 
| |
1 2 BY (MemTypeCD ...)
| |
| |  $\uparrow \text{IsMonHom}\{\langle \text{MSet}\{g\downarrow\text{set}\}, \cup, 0 \rangle, \langle \mathbb{B}, \vee_b \rangle\}(\lambda z. y1 \in_b z)$ 
| |
1 2 BY (BLemma 'mset_union_bor_mon_hom' ...)
| | \
| | 13.  $\uparrow ((\lambda z. y1 \in_b z) (\text{msFor}\{\text{mset\_mon}\{g\downarrow\text{set}\}\} x \in \text{dom}(qs). \text{mset\_inj}\{g\downarrow\text{set}\}(x * x1)))$ 
| |  $\uparrow (\lambda z. y1 \in_b z) \in \text{MonHom}(\text{mset\_mon}\{g\downarrow\text{set}\}, \langle \mathbb{B}, \vee_b \rangle)$ 
| |
1 2 BY (MemTypeCD ...)
| |
| |  $\uparrow \text{IsMonHom}\{\text{mset\_mon}\{g\downarrow\text{set}\}, \langle \mathbb{B}, \vee_b \rangle\}(\lambda z. y1 \in_b z)$ 
| |
1 2 BY (BLemma 'mset_sum_bor_mon_hom' ...)
| | \
| | 13.  $\uparrow (\exists_b\{g\downarrow\text{set}\} x \in \text{dom}(qs). y1 \in_b \text{mset\_inj}\{g\downarrow\text{set}\}(x * x1))$ 
| |  $\uparrow (\exists_b\{g\downarrow\text{set}\} v \in \text{dom}(qs). y1 \in_b \text{mset\_inj}\{g\downarrow\text{set}\}(x1 * v))$ 
| |
1 BY (RWH (LemmaC 'abmonoid_comm') 0 ...)
| | \
| | 11.  $\uparrow (x \in_b \text{dom}(ps ** qs))$ 
| |
| | BY % Apply IH %
| | FHyp 6 [11]
| |
| | 12.  $\uparrow (x \in_b \text{dom}(ps) \times \text{dom}(qs))$ 
| |
| | BY % would be more elegant to prove some more monotonicity
| | lemmas and use them here %
| |
| | (OnMCls [0;-1] (RWH (LemmaC 'mset_prod_mem')) ...a)
| |
| | 12.  $\uparrow (\exists_b\{g\downarrow\text{set}\} v \in \text{dom}(ps). \exists_b\{g\downarrow\text{set}\} w \in \text{dom}(qs). x =_b (v * w))$ 
| |  $\uparrow (\exists_b\{g\downarrow\text{set}\} v \in \text{dom}(\langle x1, y \rangle :: ps). \exists_b\{g\downarrow\text{set}\} w \in \text{dom}(qs). x =_b (v * w))$ 
| |
| | BY % Hyp must be transformed first, so squashes can be elimmed %
| | (RW (SweepDnC (LemmaC 'bmsexists_char_a_rw'
| |   ORELSEC LemmaC 'assert_of_mon_eq')) (-1)
| |   THENM Reduce (-1) ...a)
| |

```

```

12.  $\downarrow(\exists v:|g|. \uparrow(v \in_b \text{dom}(ps)) \wedge \downarrow(\exists w:|g|. \uparrow(w \in_b \text{dom}(qs)) \wedge x = v * w))$ 
|
BY (D 12 THENM CUnhide THENM ExRepD ...a)
|
12.  $v: |g|$ 
13.  $\uparrow(v \in_b \text{dom}(ps))$ 
14.  $\downarrow(\exists w:|g|. \uparrow(w \in_b \text{dom}(qs)) \wedge x = v * w)$ 
|
BY (D 14 THENM CUnhide
| THENM ExRepD ...a)
|
14.  $w: |g|$ 
15.  $\uparrow(w \in_b \text{dom}(qs))$ 
16.  $x = v * w$ 
|
BY % bug in rewrite code causes crash when renaming left out %
| RenameBVars ['v','v1';'w','w1'] 0 THEN
| (RW (SweepDnC (RevLemmaC 'bmsexists_char_rw'
| ORELSEC LemmaC 'assert_of_mon_eq')) 0
| THENM Reduce 0 ...a)
|
 $\vdash \exists v1:|g|$ 
|  $\uparrow(v1 \in_b \text{mset\_inj}\{g\downarrow\text{oset}\}(x1) + \text{dom}(ps)) \wedge (\exists w1:|g|. \uparrow(w1 \in_b \text{dom}(qs)) \wedge x = v1 * w1)$ 
|
BY (Inst ['v1';'w1'] 0 ...)
|
 $\vdash \uparrow(v \in_b \text{mset\_inj}\{g\downarrow\text{oset}\}(x1) + \text{dom}(ps))$ 
|
BY Fold 'oset_of_ocmon' 0
    THEN Reduce 0 THEN (RW bool_to_propC 0
    THENM Sel 2 (D 0) ...)
*T omral_times_assoc 648.7 sec.
 $\vdash \forall g:\text{OCMon}. \forall a:\text{CRng}. \text{Assoc}(|\text{omral}(g;a)|;\lambda ps,qs.ps ** qs)$ 
|
BY Force '5' (Eval ''assoc'' 0)
| THEN RenameBVars ['x','ps';'y','qs';'z','rs'] 0
| THEN (RepD ...a)
|
1.  $g: \text{OCMon}$ 
2.  $a: \text{CRng}$ 
3.  $ps: |\text{omral}(g;a)|$ 
4.  $qs: |\text{omral}(g;a)|$ 
5.  $rs: |\text{omral}(g;a)|$ 
 $\vdash ps ** (qs ** rs) = (ps ** qs) ** rs$ 
|
BY (BLemma 'omral_lookups_same_a'
| THENM D 0 ...a)
|
6.  $u: |g|$ 
 $\vdash (ps ** (qs ** rs))[u] = ((ps ** qs) ** rs)[u]$ 
|
BY (RWH (LemmaC 'lookup_omral_times') 0 ...a)
|
 $\vdash \text{msFor}\{a\downarrow+gp\} x \in \text{dom}(ps)$ 
|  $\text{msFor}\{a\downarrow+gp\} y \in \text{dom}(qs ** rs). \text{when}(x * y) =_b u. ps[x] * (qs ** rs)[y]$ 
|  $= \text{msFor}\{a\downarrow+gp\} x \in \text{dom}(ps ** qs)$ 
|  $\text{msFor}\{a\downarrow+gp\} y \in \text{dom}(rs). \text{when}(x * y) =_b u. (ps ** qs)[x] * rs[y]$ 

```

```

|
BY % Expand domains to simplify later 'when' cancellation %
|
| (RWN 2 (LemmaWithC ['q', [dom(qs) × dom(rs)]]
|   'mset_for_dom_shift') 0
|   THENM
|   RWN 3 (LemmaWithC ['q', [dom(ps) × dom(qs)]]
|     'mset_for_dom_shift') 0 ...a)
|\
| 7. x: |(g↓oaset)|
| 8. ↑(x ∈b dom(ps))
| ⊢ ↑(dom(qs ** rs) ⊆b dom(qs) × dom(rs))
| |
1 BY (BLemma 'omral_times_dom' ...)
|\
| 7. x: |(g↓oaset)|
| 8. ↑(x ∈b dom(ps))
| 9. x1: |(g↓oaset)|
| 10. ↑(x1 ∈b (dom(qs) × dom(rs)) - dom(qs ** rs))
| ⊢ when (x * x1) =b u. ps[x] * (qs ** rs)[x1] = e
| |
1 BY (Reduce 0 THEN
| |   RWN 2 (LemmaC 'lookup_omral_eq_zero') 0
| |   THENM RW RngNormC 0
| |   THENM RWH (LemmaC 'rng_when_of_zero') 0 ...)
| |
| ⊢ ¬↑(x1 ∈b dom(qs ** rs))
| |
1 BY % Note for when automating rewrite condition solving:
|   This is example of rewrite condition two levels deep %
|   (RWH (LemmaC 'mset_mem_diff') 10
|     THENM RW bool_to_propC 10 ...)
|\
| ⊢ ↑(dom(ps ** qs) ⊆b dom(ps) × dom(qs))
| |
1 BY (BLemma 'omral_times_dom' ...)
|\
| 7. x: |(g↓oaset)|
| 8. ↑(x ∈b (dom(ps) × dom(qs)) - dom(ps ** qs))
| ⊢ msFor{a↓+gp} y ∈ dom(rs). when (x * y) =b u. (ps ** qs)[x] * rs[y] = e
| |
1 BY (Reduce 0 THEN
| |   RWN 1 (LemmaC 'lookup_omral_eq_zero') 0
| |   THENM RW RngNormC 0 ...a)
| |\
| | 9. x1: |(g↓oaset)|
| | 10. ↑(x1 ∈b dom(rs))
| | ⊢ ¬↑(x ∈b dom(ps ** qs))
| | |
1 2 BY (RWH (LemmaC 'mset_mem_diff') 8
| |   THENM RW bool_to_propC 8 ...)
| \
| ⊢ msFor{a↓+gp} y ∈ dom(rs). when (x * y) =b u. 0 = 0
| |
1 BY (RWH (LemmaC 'rng_when_of_zero') 0
|   THENM RWH (MacroC 'x' IdC [0] ReduceC [e]) 0
|   THENM RWH (LemmaC 'mset_for_of_id') 0 ...)

```

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\
| msFor{a↓+gp} x ∈ dom(ps)
|   msFor{a↓+gp} y ∈ dom(qs) × dom(rs). when (x * y) =b u. ps[x] * (qs ** rs)[y]
| = msFor{a↓+gp} x ∈ dom(ps) × dom(qs)
|   msFor{a↓+gp} y ∈ dom(rs). when (x * y) =b u. (ps ** qs)[x] * rs[y]
|
| BY % Expand lookups of inner multiplications %
|   (RWH (LemmaC 'lookup_omral_times') 0 ...a)
|
| msFor{a↓+gp} x ∈ dom(ps)
|   msFor{a↓+gp} y ∈ dom(qs) × dom(rs)
|     when (x * y) =b u.
|       ps[x]
|         * (msFor{a↓+gp} x ∈ dom(qs)
|           msFor{a↓+gp} y1 ∈ dom(rs). when (x * y1) =b y. qs[x] * rs[y1])
| = msFor{a↓+gp} x ∈ dom(ps) × dom(qs)
|   msFor{a↓+gp} y ∈ dom(rs)
|     when (x * y) =b u.
|       (msFor{a↓+gp} x1 ∈ dom(ps)
|         msFor{a↓+gp} y ∈ dom(qs). when (x1 * y) =b x. ps[x1] * qs[y])
|         * rs[y]
|
| BY Fold 'rng_mssum' 0
|
| (Σx ∈ dom(ps).
|   Σy ∈ dom(qs) × dom(rs).
|     when (x * y) =b u.
|       ps[x] * (Σx ∈ dom(qs). Σy1 ∈ dom(rs). when (x * y1) =b y. qs[x] * rs[y1]))
| = (Σx ∈ dom(ps) × dom(qs).
|   Σy ∈ dom(rs).
|     when (x * y) =b u.
|       (Σx1 ∈ dom(ps). Σy ∈ dom(qs). when (x1 * y) =b x. ps[x1] * qs[y]) * rs[y])
|
| BY % Float up the Sigma's and when's %
|   (RWH "rng_times_mssum_l
|     rng_times_mssum_r
|     rng_mssum_when_swap<
|     rng_times_when_l
|     rng_times_when_r" 0 ...a)
|
| (Σx ∈ dom(ps).
|   Σy ∈ dom(qs) × dom(rs).
|     Σx1 ∈ dom(qs).
|       Σy1 ∈ dom(rs). when (x * y) =b u. when (x1 * y1) =b y. ps[x] * (qs[x1] * rs[y1]))
| = (Σx ∈ dom(ps) × dom(qs).
|   Σy ∈ dom(rs).
|     Σx1 ∈ dom(ps).
|       Σy1 ∈ dom(qs). when (x * y) =b u. when (x1 * y1) =b x. (ps[x1] * qs[y1]) * rs[y])
|
| BY % Bring together the Sigma's and when's that cancel %
|   (RW (NthsC [2;3]
|     (HereDnC (LemmaC 'rng_mssum_swap'
|       ORELSEC LemmaC 'rng_when_swap')))) 0 ...a)
|
| (Σx ∈ dom(ps).
|   Σx1 ∈ dom(qs).
|   Σy1 ∈ dom(rs).

```



```

|       $\Sigma y \in \text{dom}(qs) \times \text{dom}(rs).$ 
|      when  $(x1 * y1) =_b y.$  when  $(x * y) =_b u.$   $ps[x] * (qs[x1] * rs[y1])$ )
| =  $(\Sigma y \in \text{dom}(rs).$ 
|    $\Sigma x1 \in \text{dom}(ps).$ 
|    $\Sigma y1 \in \text{dom}(qs).$ 
|    $\Sigma x \in \text{dom}(ps) \times \text{dom}(qs).$ 
|   when  $(x1 * y1) =_b x.$  when  $(x * y) =_b u.$   $(ps[x1] * qs[y1]) * rs[y]$ )
|
BY % Setup and do cancellation %
| (
|   RWNs [1;3] (LemmaC 'grp_eq_sym') 0
|   THENM RWNs [1;3] oset_of_ocmonC 0
|   THENM RWH (LemmaC 'fset_for_when_eq') 0 ...a)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| 9. x1: |(g↓oset)|
| 10. ↑(x1 ∈b dom(qs))
| 11. x2: |(g↓oset)|
| 12. ↑(x2 ∈b dom(rs))
| ⊢ ↑(x1 * x2 ∈b dom(qs) × dom(rs))
| |
1 BY (BLemma 'prod_in_mset_prod' ...)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(rs))
| 9. x1: |(g↓oset)|
| 10. ↑(x1 ∈b dom(ps))
| 11. x2: |(g↓oset)|
| 12. ↑(x2 ∈b dom(qs))
| ⊢ ↑(x1 * x2 ∈b dom(ps) × dom(qs))
| |
1 BY (BLemma 'prod_in_mset_prod' ...)
\
| ⊢  $(\Sigma x \in \text{dom}(ps).$ 
|    $\Sigma x1 \in \text{dom}(qs).$   $\Sigma y1 \in \text{dom}(rs).$  when  $(x * (x1 * y1)) =_b u.$   $ps[x] * (qs[x1] * rs[y1])$ )
| =  $(\Sigma y \in \text{dom}(rs).$ 
|    $\Sigma x1 \in \text{dom}(ps).$   $\Sigma y1 \in \text{dom}(qs).$  when  $((x1 * y1) * y) =_b u.$   $(ps[x1] * qs[y1]) * rs[y]$ )
|
BY % Final normalization of l and r sides to same %
  (RWN 3 (HereDnC (LemmaC 'rng_mssum_swap')) 0
  THENM RW MonNormC 0
  THENM RW RngNormC 0 ...)
*T omral_times_assoc_a 0.4 sec.
⊢  $\forall g:\text{OCMon}. \forall a:\text{CRng}. \forall ps,qs,rs:|\text{omral}(g;a)|. ps ** (qs ** rs) = (ps ** qs) ** rs$ 
|
BY AssertLemma 'omral_times_assoc' []
  THENM Force '6' (Eval ''assoc'' (-1))
  THEN Trivial
#T omral_times_assoc_b 351.6 sec.
⊢  $\forall g:\text{OCMon}. \forall a:\text{CRng}. \forall ps,qs,rs:|\text{omral}(g;a)|. ps ** (qs ** rs) = (ps ** qs) ** rs$ 
|
BY (RepD
|   THENM BLemma 'omral_lookups_same_a'
|   THENM RepD ...a)
|
1. g: OCMon

```

```

2. a: CRng
3. ps: |omral(g;a)|
4. qs: |omral(g;a)|
5. rs: |omral(g;a)|
6. u: |g|
| (ps ** (qs ** rs))[u] = ((ps ** qs) ** rs)[u]
|
BY % discard one side to avoid clutter %
| SplitRel 「0」
| \
| | (ps ** (qs ** rs))[u] = 0
| |
1 BY (RWO "lookup_omral_times_a" 0 ...a)
| |
| | (Σx ∈ dom(ps). Σy ∈ dom(qs ** rs). when (x * y) =b u. ps[x] * (qs ** rs)[y]) = 0
| |
1 BY (RWO "omral_times_dom" 0 ...a)
| | \
| | | 7. x: |(g↓oset)|
| | | 8. ↑(x ∈b dom(ps))
| | | 9. x1: |(g↓oset)|
| | | 10. ↑(x1 ∈b (dom(qs) × dom(rs)) - dom(qs ** rs))
| | | (ps[x] * (qs ** rs)[x1]) = 0
| | |
1 2 BY (RWN 2 (LemmaC 'lookup_omral_eq_zero') 0
| | | THENM RW RngNormC 0
| | | THENM RWH (LemmaC 'rng_when_of_zero') 0 ...)
| | |
| | | ¬↑(x1 ∈b dom(qs ** rs))
| | |
1 2 BY (RWH (LemmaC 'mset_mem_diff') 10
| | | THENM RW bool_to_propC 10 ...)
| | | \
| | | (Σx ∈ dom(ps). Σy ∈ dom(qs) × dom(rs). when (x * y) =b u. ps[x] * (qs ** rs)[y]) = 0
| | |
1 BY (RWW "lookup_omral_times_a" 0 ...a)
| |
| | (Σx ∈ dom(ps).
| | | Σy ∈ dom(qs) × dom(rs).
| | | when (x * y) =b u.
| | | ps[x] * (Σx1 ∈ dom(qs). Σy1 ∈ dom(rs). when (x1 * y1) =b y. qs[x1] * rs[y1]))
| | | = 0
| | |
1 BY % float up sigmas and whens %
| | | (RWW "rng_times_mssum_l
| | | | rng_times_mssum_r
| | | | rng_mssum_when_swap<
| | | | rng_times_when_l
| | | | rng_times_when_r" 0 ...a)
| | |
| | | (Σx ∈ dom(ps).
| | | | Σy ∈ dom(qs) × dom(rs).
| | | | Σx1 ∈ dom(qs).
| | | | Σy1 ∈ dom(rs). when (x * y) =b u. when (x1 * y1) =b y. ps[x] * (qs[x1] * rs[y1]))
| | | | = 0
| | |
1 BY (RWN 2 (HereDnC (PolyC "rng_mssum_swap rng_when_swap"))) 0 ...a)

```

```

| |
| |  $\vdash (\sum x \in \text{dom}(\text{ps}).$ 
| |    $\sum x1 \in \text{dom}(\text{qs}).$ 
| |      $\sum y1 \in \text{dom}(\text{rs}).$ 
| |        $\sum y \in \text{dom}(\text{qs}) \times \text{dom}(\text{rs}).$ 
| |         when  $(x1 * y1) =_b y.$  when  $(x * y) =_b u.$   $\text{ps}[x] * (\text{qs}[x1] * \text{rs}[y1]))$ 
| |   = 0
| |
1 BY (Unfold 'oset_of_ocmon' 0 ...a)
| |
| |  $\vdash (\sum x \in \text{dom}(\text{ps}).$ 
| |    $\sum x1 \in \text{dom}(\text{qs}).$ 
| |      $\sum y1 \in \text{dom}(\text{rs}).$ 
| |        $\sum y \in \text{dom}(\text{qs}) \times \text{dom}(\text{rs}).$ 
| |         when  $(x1 * y1) =_b y.$  when  $(x * y) =_b u.$   $\text{ps}[x] * (\text{qs}[x1] * \text{rs}[y1]))$ 
| |   = 0
| |
1 BY (RWN 1 (LemmaC 'grp_eq_sym') 0
| |   THENM RWH dset_of_monC 0 ...a)
| |
| |  $\vdash (\sum x \in \text{dom}(\text{ps}).$ 
| |    $\sum x1 \in \text{dom}(\text{qs}).$ 
| |      $\sum y1 \in \text{dom}(\text{rs}).$ 
| |        $\sum y \in \text{dom}(\text{qs}) \times \text{dom}(\text{rs}).$ 
| |         when  $y =_b (x1 * y1).$  when  $(x * y) =_b u.$   $\text{ps}[x] * (\text{qs}[x1] * \text{rs}[y1]))$ 
| |   = 0
| |
1 BY (RWO "rng_fset_for_when_eq" 0 ...a)
| | \
| | 7. x: |(g↓set)|
| | 8. ↑(x ∈b dom(ps))
| | 9. x1: |(g↓set)|
| | 10. ↑(x1 ∈b dom(qs))
| | 11. x2: |(g↓set)|
| | 12. ↑(x2 ∈b dom(rs))
| |  $\vdash \uparrow(x1 * x2 \in_b \text{dom}(\text{qs}) \times \text{dom}(\text{rs}))$ 
| | |
1 2 BY (BLemma 'prod_in_mset_prod' ...)
| | \
| |  $\vdash (\sum x \in \text{dom}(\text{ps}).$ 
| |    $\sum x1 \in \text{dom}(\text{qs}). \sum y1 \in \text{dom}(\text{rs}).$  when  $(x * (x1 * y1)) =_b u.$   $\text{ps}[x] * (\text{qs}[x1] * \text{rs}[y1]))$ 
| |   = 0
| | |
| | INCOMPLETE
| | \
| |  $\vdash 0 = ((\text{ps} ** \text{qs}) ** \text{rs})[u]$ 
| | |
| | INCOMPLETE
*T omral_times_comm 52.7 sec.
 $\vdash \forall g:\text{OCMon}. \forall a:\text{CRng}. \text{Comm}(|\text{omral}(g;a)|; \lambda \text{ps}, \text{qs}. \text{ps} ** \text{qs})$ 
|
BY Force '5' (Eval ''comm'' 0)
| THEN RenameBVars ['x','ps'; 'y','qs'] 0
| THEN (RepD ...a)
|
1. g: OCMon
2. a: CRng

```

```

3. ps: |omral(g;a)|
4. qs: |omral(g;a)|
⊢ ps ** qs = qs ** ps
|
BY (BLemma 'omral_lookups_same_a'
| THENM D 0 ...a)
|
5. u: |g|
⊢ (ps ** qs)[u] = (qs ** ps)[u]
|
BY (RWH (LemmaC 'lookup_omral_times') 0
| THENM Fold 'rng_mssum' 0 ...a)
|
⊢ (∑x ∈ dom(ps). ∑y ∈ dom(qs). when (x * y) =b u. ps[x] * qs[y])
| = (∑x ∈ dom(qs). ∑y ∈ dom(ps). when (x * y) =b u. qs[x] * ps[y])
|
BY % make bvars same in each equand %
| RenameBVars ['x','y';'y','x'] 0
|
|
BY % normalize %
(RWN 1 (LemmaC 'rng_mssum_swap') 0
THENM RW CRngNormC 0
THENM RW AbMonNormC 0 ...)
*T omral_times_comm_a 0.3 sec.
⊢ ∀g:OCMon. ∀a:CRng. ∀ps,qs:|omral(g;a)|. ps ** qs = qs ** ps
|
BY AssertLemma 'omral_times_comm' []
THEN Force '5' (Eval ''comm'' 1)
THEN Trivial
*T omral_bilinear 155.1 sec.
⊢ ∀g:OCMon. ∀a:CRng. BiLinear(|omral(g;a)|;λps,qs.ps ++ qs;λps,qs.ps ** qs)
|
BY (RepD THENM BLemma 'bilinear_comm_elim'
| THENM Force '5' (Reduce 0) ...a)
| \
| 1. g: OCMon
| 2. a: CRng
| ⊢ Comm(|omral(g;a)|;λps,qs.ps ** qs)
| |
1 BY (BLemma 'omral_times_comm' ...)
\
1. g: OCMon
2. a: CRng
⊢ ∀a1,x,y:|omral(g;a)|. a1 ** (x ++ y) = (a1 ** x) ++ (a1 ** y)
|
BY (RenameBVars ['a1','ps';'x','qs';'y','rs'] 0
| THENM RepD
| THENM BLemma 'omral_lookups_same_a'
| THENM D 0 ...a)
|
3. ps: |omral(g;a)|
4. qs: |omral(g;a)|
5. rs: |omral(g;a)|
6. u: |g|
⊢ (ps ** (qs ++ rs))[u] = ((ps ** qs) ++ (ps ** rs))[u]
|

```

```

BY % distribute lookup over plus and times %
| (RWW "lookup_omral_plus lookup_omral_times
|   f:rng_mssum" 0 ...a)
|
|  $(\sum x \in \text{dom}(ps). \sum y \in \text{dom}(qs ++ rs). \text{when } (x * y) =_b u. ps[x] * (qs[y] +a rs[y]))$ 
| =  $(\sum x \in \text{dom}(ps). \sum y \in \text{dom}(qs). \text{when } (x * y) =_b u. ps[x] * qs[y])$ 
|   +a  $(\sum x \in \text{dom}(ps). \sum y \in \text{dom}(rs). \text{when } (x * y) =_b u. ps[x] * rs[y])$ 
|
BY % equalize domains of summation %
|
| (RWNs [2;4;6] (LemmaWithC ['q', $\lceil \text{dom}(qs) \cup \text{dom}(rs) \rceil$ ]
|   'mset_for_dom_shift') 0
|   THENM Fold 'rng_mssum' 0 ...a)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| ⊢ ↑(dom(qs ++ rs) ⊆b dom(qs) ∪ dom(rs))
| |
1 BY (BLemma 'omral_plus_dom' ...)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| 9. x1: |(g↓oset)|
| 10. ↑(x1 ∈b (dom(qs) ∪ dom(rs)) - dom(qs ++ rs))
| ⊢ when (x * x1) =b u. ps[x] * (qs[x1] +a rs[x1]) = e
| |
1 BY % Have to backtrack a little with lookup distributing %
| |
| | (Reduce 0
| |   THENM RWH (RevLemmaC 'lookup_omral_plus') 0
| |   THENM RWN 2 (LemmaC 'lookup_omral_eq_zero') 0 ...a)
| |\
| | ⊢ ¬↑(x1 ∈b dom(qs ++ rs))
| | |
1 2 BY (RWH (LemmaC 'mset_mem_diff') 10
| |   THENM RW bool_to_propC 10 ...)
| | \
| | ⊢ when (x * x1) =b u. ps[x] * 0 = 0
| | |
1 BY (RW RngNormC 0
|   THENM RWH (LemmaC 'rng_when_of_zero') 0 ...)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| ⊢ ↑(dom(qs) ⊆b dom(qs) ∪ dom(rs))
| |
1 BY (BLemma 'mem_bsubset'
|   THENM RepD
|   THENM RWH (LemmaC 'fset_mem_union') 0
|   THENM RW bool_to_propC 0
|   THENM Sel 1 (D 0) ...)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| 9. x1: |(g↓oset)|
| 10. ↑(x1 ∈b (dom(qs) ∪ dom(rs)) - dom(qs))
| ⊢ when (x * x1) =b u. ps[x] * qs[x1] = e

```

```

| |
1 BY (RWH (LemmaC 'mset_mem_diff') 10
|   THENM RW bool_to_propC 10
|   THENM Reduce 0
|   THENM RWN 2 (LemmaC 'lookup_omral_eq_zero') 0
|   THENM RW RngNormC 0
|   THENM RWH (LemmaC 'rng_when_of_zero') 0 ...)
| \
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| ⊢ ↑(dom(rs) ⊆b dom(qs) ∪ dom(rs))
| |
1 BY (BLemma 'mem_bsubset'
|   THENM RepD
|   THENM RWH (LemmaC 'fset_mem_union') 0
|   THENM RW bool_to_propC 0
|   THENM Sel 2 (D 0) ...)
| \
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| 9. x1: |(g↓oset)|
| 10. ↑(x1 ∈b (dom(qs) ∪ dom(rs)) - dom(rs))
| ⊢ when (x * x1) =b u. ps[x] * rs[x1] = e
| |
1 BY (RWH (LemmaC 'mset_mem_diff') 10
|   THENM RW bool_to_propC 10
|   THENM Reduce 0
|   THENM RWN 2 (LemmaC 'lookup_omral_eq_zero') 0
|   THENM RW RngNormC 0
|   THENM RWH (LemmaC 'rng_when_of_zero') 0 ...)
| \
| ⊢ (∑x ∈ dom(ps). ∑y ∈ dom(qs) ∪ dom(rs). when (x * y) =b u. ps[x] * (qs[y] +a rs[y]))
|   = (∑x ∈ dom(ps). ∑y ∈ dom(qs) ∪ dom(rs). when (x * y) =b u. ps[x] * qs[y])
|     +a (∑x ∈ dom(ps). ∑y ∈ dom(qs) ∪ dom(rs). when (x * y) =b u. ps[x] * rs[y])
|
| BY % Distribute xa, Sigma and when over + %
|   (RWW "rng_times_over_plus.1
|     rng_mssum_of_plus
|     rng_when_thru_plus" 0 ...)
*T omral_bilinear_a 0.8 sec.
⊢ ∀g:OCMon. ∀a:CRng. ∀ps,qs,rs:|omral(g;a)|.
|   ps ** (qs ++ rs) = (ps ** qs) ++ (ps ** rs) ∧ (qs ++ rs) ** ps = (qs ** ps) ++ (rs ** ps)
|
BY AssertLemma 'omral_bilinear' []
  THEN Force '5' (Eval ''bilinear'' 1)
  THEN Trivial
*C omral_one_act_com
  =====
  OMRAL ONE AND ACTION
  =====
*D omral_one_df
  Parends ::Prec(preop):: 11<g:g:L>,<r:r:L>== omral_one{<g>; <r>}
  11== omral_one{<g>; <r>}
*A omral_one 11 == inj(e,1)
*T omral_one_wf 1.1 sec.
⊢ ∀g:OCMon. ∀r:CRng. 11 ∈ |omral(g;r)|
|

```

```

BY (Unfold 'omral_one' 0 ...)
*T omral_dom_one 12.4 sec.
├  $\forall g:OCMon. \forall r:CRng. \neg(0 = 1) \Rightarrow \text{dom}(11) = \text{mset\_inj}\{g \downarrow \text{oset}\}(e)$ 
|
BY (RepD THENM RWW "u:omral_one omral_dom_inj" 0 ...a)
|
1. g: OCMon
2. r: CRng
3.  $\neg(0 = 1)$ 
├ if 1 =b 0 then 0{g↓oset} else mset_inj{g↓oset}(e) fi = mset_inj{g↓oset}(e)
|
BY (SplitOnConclITE ...a)
|\
| 4. 1 = 0
| ─ 0{g↓oset} = mset_inj{g↓oset}(e)
| |
1 BY (RelRST ...)
\
4.  $\neg(1 = 0)$ 
├ mset_inj{g↓oset}(e) = mset_inj{g↓oset}(e)
|
BY Auto
*D omral_action_df
    Parens ::Prec(inop)::
      <v:v:L> ..<g:g:L>, <r:r:L> <ps:ps:L>
      == omral_action{<g>; <r>; <v>; <ps>}
    Parens ::Prec(inop)::
      <v:v:L> .. <ps:ps:L>
      == omral_action{<g>; <r>; <v>; <ps>}
*A omral_action          v .. ps == <e,v>* ps
*T omral_action_wf 2.0 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall v:|r|. \forall ps:|\text{omral}(g;r)|. v .. ps \in |\text{omral}(g;r)|$ 
|
BY (Unfold 'omral_action' 0 ...)
*T comb_for_omral_action_wf 1.6 sec.
├  $(\lambda g,r,v,ps,z.v .. ps) \in g:OCMon$ 
|
|                                     → r:CRng
|                                     → v:|r|
|                                     → ps:|\text{omral}(g;r)|
|                                     → ↓True
|                                     → |\text{omral}(g;r)|
|
BY ProveOpCombTyping 'omral_action_wf'
*C omral_dom_action_com
    Nice simple example of monotonicity
    reasoning here.
*T omral_dom_action 24.0 sec.
├  $\forall g:OCMon. \forall r:CRng. \forall v:|r|. \forall ps:|\text{omral}(g;r)|. \uparrow(\text{dom}(v .. ps) \subseteq_b \text{dom}(ps))$ 
|
BY (RepD ...a)
|
1. g: OCMon
2. r: CRng
3. v: |r|
4. ps: |\text{omral}(g;r)|
├  $\uparrow(\text{dom}(v .. ps) \subseteq_b \text{dom}(ps))$ 
|

```

```

BY % A nice example of monotonic reasoning %
|
| (Unfold 'omral_action' 0
|   THENM RWH (LemmaC 'omral_dom_scale') 0 ...a)
|
|  $\vdash \uparrow(\text{fs-map}(\lambda k'.k' * e, \text{dom}(\text{ps})) \subseteq_b \text{dom}(\text{ps}))$ 
|
BY (BLemma 'mem_bsubset' THENM RepD ...a)
|
5. x: |(g↓oset)|
6.  $\uparrow(x \in_b \text{fs-map}(\lambda k'.k' * e, \text{dom}(\text{ps})))$ 
|  $\vdash \uparrow(x \in_b \text{dom}(\text{ps}))$ 
|
BY (RWW "u:fset_map fset_of_mset_mem" 6 ...a)
|
6.  $\uparrow(x \in_b \text{mmap}\{g\downarrow\text{oset}, g\downarrow\text{oset}\}(\lambda k'.k' * e; \text{dom}(\text{ps})))$ 
|
BY (RWH (AssertC  $\uparrow(\lambda k'.k' * e) = \text{Id}\{|(g\downarrow\text{oset})|\}$ ) 6
|   THENM RWW "mset_map_id" 6 ...)
|
|  $\vdash (\lambda k'.k' * e) = \text{Id}\{|(g\downarrow\text{oset})|\}$ 
|
BY (Ext THENM Reduce 0 THENM RW MonNormC 0 ...)
*T lookup_omral_action 12.9 sec.
|  $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall k:|g|. \forall v:|r|. \forall ps:|\text{omral}(g;r)|. (v \cdot ps)[k] = v * ps[k]$ 
|
BY (RepD THENM Unfold 'omral_action' 0 ...a)
|
1. g: OCMon
2. r: CRng
3. k: |g|
4. v: |r|
5. ps: |omral(g;r)|
|  $\vdash (<e,v>* ps)[k] = v * ps[k]$ 
|
BY (ReplaceWithEqv
|   MonNormC
|    $\uparrow(<e,v>* ps)[e * k] = v * ps[k]$ 
|   0 ...a)
|
|  $\vdash (<e,v>* ps)[e * k] = v * ps[k]$ 
|
BY (RWW "lookup_omral_scale_a" 0 ...)
*C omral_alg_com
=====
ASSEMBLY OF OMRAL FREE MONOID ALGEBRA
=====
*T omral_times_ident_r 2.1 sec.
|  $\vdash \forall g:\text{OCMon}. \forall r:\text{CRng}. \forall ps:|\text{omral}(g;r)|. ps ** 11 = ps$ 
|
BY (RepD
|   THENM InstLemma 'omral_times_comm'  $\uparrow[g]; \uparrow[r]$  ...a)
|
1. g: OCMon
2. r: CRng
3. ps: |omral(g;r)|
4. Comm(|omral(g;r)|;  $\lambda ps, qs. ps ** qs$ )

```



```

┆ ps ** 11 = ps
|
BY (Force '5' (Eval 'comm' (-1))
    THENM RW (HypC (-1)) 0
    THENM BLemma 'omral_times_ident_1' ...a)
*T omral_times_ident_1 92.8 sec.
┆  $\forall g:OCMon. \forall r:CRng. \forall ps:|omral(g;r)|. 11 ** ps = ps$ 
|
BY (RepD ...a)
|
1. g: OCMon
2. r: CRng
3. ps: |omral(g;r)|
┆ 11 ** ps = ps
|
BY Unfold 'omral_one' 0
|
┆ inj(e,1) ** ps = ps
|
BY (BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
4. u: |g|
┆ (inj(e,1) ** ps)[u] = ps[u]
|
BY (RWW "lookup_omral_times omral_dom_inj" 0 ...a)
|
┆ msFor{r↓+gp} x ∈ if 1 =b 0 then 0{g↓oset} else mset_inj{g↓oset}(e) fi
|   msFor{r↓+gp} y ∈ dom(ps). when (x * y) =b u. inj(e,1)[x] * ps[y]
|   = ps[u]
|
BY (SplitOnConclITE ...a)
| \
| 5. 1 = 0
|   ┆ msFor{r↓+gp} x ∈ 0{g↓oset}
|     ┆ msFor{r↓+gp} y ∈ dom(ps). when (x * y) =b u. inj(e,1)[x] * ps[y]
|     ┆ = ps[u]
|     |
|     1 BY Reduce 0
|     |
|     ┆ 0 = ps[u]
|     |
|     1 BY (InvertRel 0
|         THENM BLemma 'ring_triv' ...)
|     \
|     5.  $\neg(1 = 0)$ 
|     ┆ msFor{r↓+gp} x ∈ mset_inj{g↓oset}(e)
|       ┆ msFor{r↓+gp} y ∈ dom(ps). when (x * y) =b u. inj(e,1)[x] * ps[y]
|       ┆ = ps[u]
|       |
|       BY (RWW "mset_for_mset_inj" 0
|           THENM RW MonNormC 0 ...a)
|       |
|       ┆ msFor{r↓+gp} y ∈ dom(ps). when y =b u. inj(e,1)[e] * ps[y] = ps[u]
|       |
|       BY (RWW "lookup_omral_inj" 0 ...)
|       |
|       ┆ msFor{r↓+gp} y ∈ dom(ps). when y =b u. (when e =b e. 1) * ps[y] = ps[u]

```

```

|
BY (RWN 2 (LemmaC 'mon_when_true') 0
|   THENM RW RngNormC 0 ...a)
|\
| 6. x: |(g↓oset)|
| 7. ↑(x ∈b dom(ps))
| ⊢ ↑(e =b e)
| |
1 BY (RW bool_to_propC 0 ...)
\
| ⊢ msFor{r↓+gp} y ∈ dom(ps). when y =b u. ps[y] = ps[u]
|
BY (Decide [↑(u ∈b dom(ps))] ...a)
|\
| 6. ↑(u ∈b dom(ps))
| |
1 BY (RWH oset_of_ocmonC 0
|   THENM RWW "fset_for_when_eq" 0 ...)
\
| 6. ¬↑(u ∈b dom(ps))
|
BY (RWW "mset_for_when_none" 0 ...a)
|\
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps))
| ⊢ ¬↑(x =b u)
| |
1 BY (D 0 THENM RW bool_to_propC (-1)
|   THENM RWW "-1" 8 ...)
\
| ⊢ e = ps[u]
|
BY (Reduce 0
|   THENM RWW "lookup_omral_eq_zero" 0 ...)
*T omral_action_one 8.3 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀ps:|omral(g;r)|. 1 .. ps = ps
|
BY (RepD THENM BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. ps: |omral(g;r)|
4. u: |g|
⊢ (1 .. ps)[u] = ps[u]
|
BY (RWW "lookup_omral_action" 0
|   THENM RW RngNormC 0 ...)
*T omral_action_times 11.5 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀v,w:|r|. ∀ps:|omral(g;r)|. (v * w) .. ps = v .. (w .. ps)
|
BY (RepD THENM BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. v: |r|
4. w: |r|
5. ps: |omral(g;r)|

```

```

6. u: |g|
⊢ ((v * w) .. ps)[u] = (v .. (w .. ps))[u]
|
BY (RWW "lookup_omral_action" 0
    THENM RW RngNormC 0 ...)
*T omral_action_times_r1 135.2 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀v:|r|. ∀ps,qs:|omral(g;r)|. v .. (ps ** qs) = (v .. ps) ** qs
|
BY (RepD THENM BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. v: |r|
4. ps: |omral(g;r)|
5. qs: |omral(g;r)|
6. u: |g|
⊢ (v .. (ps ** qs))[u] = ((v .. ps) ** qs)[u]
|
BY (RWW "lookup_omral_action lookup_omral_times" 0
    | THENM Fold 'rng_mssum' 0 ...a)
|
⊢ v * (Σx ∈ dom(ps). Σy ∈ dom(qs). when (x * y) =b u. ps[x] * qs[y])
| = (Σx ∈ dom(v .. ps). Σy ∈ dom(qs). when (x * y) =b u. (v * ps[x]) * qs[y])
|
BY (Unfold 'rng_mssum' 0
    | THENM RWH (LemmaC 'omral_dom_action') 0
    | THENM Fold 'rng_mssum' 0 ...a)
| \
| 7. x: |(g↓oset)|
| 8. ↑(x ∈b dom(ps) - dom(v .. ps))
| ⊢ msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b u. (v * ps[x]) * qs[y] = e
| |
1 BY % Fold back action in concl to make proof easy %
| | (RWN 2 (RevLemmaC 'lookup_omral_action') 0
| | THENM RWW "mset_mem_diff" 8
| | THENM RW bool_to_propC 8
| | THENM RepD ...a)
| |
| 8. ↑(x ∈b dom(ps))
| 9. ¬↑(x ∈b dom(v .. ps))
| ⊢ msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b u. (v .. ps)[x] * qs[y] = e
| |
1 BY (RWN 1 (LemmaC 'lookup_omral_eq_zero') 0
| | THENM Reduce 0 ...a)
| |
| ⊢ msFor{r↓+gp} y ∈ dom(qs). when (x * y) =b u. 0 * qs[y] = 0
| |
1 BY (RW RngNormC 0
| THENM RWH (LemmaC 'rng_when_of_zero') 0
| THENM RWH add_grp_of_rngC 0
| THENM RWH (LemmaC 'mset_for_of_id') 0
| THENM Reduce 0 ...)
| \
| ⊢ v * (Σx ∈ dom(ps). Σy ∈ dom(qs). when (x * y) =b u. ps[x] * qs[y])
| = (Σx ∈ dom(ps). Σy ∈ dom(qs). when (x * y) =b u. (v * ps[x]) * qs[y])
|
BY (RWW "rng_times_mssum_1 rng_times_when_1" 0

```

```

    THENM RW RngNormC 0 ...)
*T omral_action_times_r2 5.0 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀v:|r|. ∀ps,qs:|omral(g;r)|. v .. (ps ** qs) = ps ** (v .. qs)
|
BY (RepD
    THENM RWH (LemmaC 'omral_times_comm_a' 0
    THENM RWO "omral_action_times_r1<" 0 ...)
*T omral_action_plus_l 13.5 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀v,w:|r|. ∀ps:|omral(g;r)|. (v +r w) .. ps = (v .. ps) ++ (w .. ps)
|
BY (RepD THENM BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. v: |r|
4. w: |r|
5. ps: |omral(g;r)|
6. u: |g|
⊢ ((v +r w) .. ps)[u] = ((v .. ps) ++ (w .. ps))[u]
|
BY (RWW "lookup_omral_action lookup_omral_plus" 0
    THENM RW RngNormC 0 ...)
*T omral_action_plus_r 17.0 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀v:|r|. ∀ps,qs:|omral(g;r)|. v .. (ps ++ qs) = (v .. ps) ++ (v .. qs)
|
BY (RepD THENM BLemma 'omral_lookups_same_a' THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. v: |r|
4. ps: |omral(g;r)|
5. qs: |omral(g;r)|
6. u: |g|
⊢ (v .. (ps ++ qs))[u] = ((v .. ps) ++ (v .. qs))[u]
|
BY (RWW "lookup_omral_action lookup_omral_plus" 0
    THENM RW RngNormC 0 ...)
*T omral_action_inj 6.9 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k:|g|. ∀v,v':|r|. v .. inj(k,v') = inj(k,v * v')
|
BY (RepD THENM BLemma 'omral_lookups_same_a'
| THENM D 0
| THENM RWW "lookup_omral_action lookup_omral_inj" 0 ...a)
|
1. g: OCMon
2. r: CRng
3. k: |g|
4. v: |r|
5. v': |r|
6. u: |g|
⊢ v * (when k =b u. v') = when k =b u. v * v'
|
BY (Fold 'rng_when' 0
    THENM RWW "rng_times_when_l" 0 ...)
*D omral_alg_df omral_alg(<g:g:*>;<r:r:*>)== omral_alg{<g>;<r>}
*A omral_alg omral_alg(g;r) ==
    <|omral(g;r)|

```

```

    , =b
    , λx,y.tt
    , λx,y.x ++ y
    , 00g,r
    , λx.--x
    , λx,y.x ** y
    , 11
    , λx,y.(inr . )
    , λa,x.a .. x>
*T omral_alg_wf 22.5 sec.
⊢ ∀g:OCMon. ∀r:CRng. omral_alg(g;r) ∈ AlgebraSig(|r|)
|
BY (Unfolds ‘‘omral_alg algebra_sig‘‘ 0 ...)
*T omral_alg_wf2 60.0 sec.
⊢ ∀g:OCMon. ∀r:CRng. omral_alg(g;r) ∈ r-CAgebra
|
BY (RepD THENM RepeatM MemTypeCD
| THENM Force ‘5‘ (Reduce 0) ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| ⊢ omral_alg(g;r) ∈ AlgebraSig(|r|)
| |
1 BY Auto
|\
| 1. g: OCMon
| 2. r: CRng
| ⊢ IsGroup(|omral_alg(g;r)|;+omral_alg(g;r);0omral_alg(g;r);-omral_alg(g;r))
| |
1 BY (Assert [oal_grp(g↓oset;r↓+gp) ∈ Group{i}]
| | THENM AddAllProperties (-1)
| | THENM All (\i.Force ‘5‘ (Eval ‘‘oal_grp omral_alg‘‘ i)) ...a)
| |
| 3. <|oal(g↓oset;r↓+gp)|, =b, λx,y.x ≤b y, λx,y.x ++ y, 00, λx.--x> ∈ Group{i}
| 4. IsMonoid(|oal(g↓oset;r↓+gp)|;λx,y.x ++ y;00) ∧ IsEqFun(|oal(g↓oset;r↓+gp)|;=b)
| 5. Inverse(|oal(g↓oset;r↓+gp)|;λx,y.x ++ y;00;λx.--x)
| ⊢ IsGroup(|omral(g;r)|;λx,y.x ++ y;00g,r;λx.--x)
| |
1 BY Unfolds ‘‘omralist omral_plus omral_zero omral_minus‘‘ 0
| THEN (AGenRepD ["compound"] ...)
|\
| 1. g: OCMon
| 2. r: CRng
| ⊢ Comm(|omral_alg(g;r)|;+omral_alg(g;r))
| |
1 BY % Need finer grading of strengths to avoid the
| ugliness of the first 3 steps here %
|
| (ARepD ["basic"]
| THENM RWH AbRedexC 0
| THENM Force ‘2‘ (Reduce 0)
| THENM BLemma ‘omral_plus_comm‘ ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| ⊢ IsAction(|r|;*;1;|omral_alg(g;r)|;·omral_alg(g;r))
| |

```

```

1 BY (AGenRepD ["compound";"basic"] ...a)
| | \
| | 3. a: |r|
| | 4. b: |r|
| | 5. u: |omral_alg(g;r)|
| | ⊢ (a * b) ·omral_alg(g;r) u = a ·omral_alg(g;r) (b ·omral_alg(g;r) u)
| | |
1 2 BY % Can't keep reducing or would go too far... %
| | | RWH AbRedexC 0 THENM Force '5' (Reduce 0)
| | |
| | | ⊢ (a * b) .. u = a .. (b .. u)
| | |
1 2 BY (BLemma 'omral_action_times' ...)
| | \
| | 3. u: |omral_alg(g;r)|
| | ⊢ 1 ·omral_alg(g;r) u = u
| | |
1 BY RWH AbRedexC 0 THENM Force '5' (Reduce 0)
| | |
| | | ⊢ 1 .. u = u
| | |
1 BY (BLemma 'omral_action_one' ...)
| | \
| | 1. g: OCMon
| | 2. r: CRng
| | ⊢ IsBilinear(|r|;|omral_alg(g;r)|;|omral_alg(g;r)|;+r;+omral_alg(g;r);+omral_alg(g;r);·omral_
| | |
| | | alg(g;r))
| | |
1 BY Unfolds 'bilinear_p' 0
| | THEN RWH AbRedexC 0 THEN Force '5' (Reduce 0)
| | THEN (Backchain 'omral_action_plus_l omral_action_plus_r' ...a)
| | \
| | 1. g: OCMon
| | 2. r: CRng
| | ⊢ IsEqFun(|omral_alg(g;r)|;omral_alg(g;r).eq)
| | |
1 BY RWH AbRedexC 0
| | THENM (Assert [|omral(g;r) ∈ DSet|
| | THENM AddProperties (-1) ...)
| | \
| | 1. g: OCMon
| | 2. r: CRng
| | ⊢ IsMonoid(|omral_alg(g;r)|;xomral_alg(g;r);1omral_alg(g;r))
| | |
1 BY RWH AbRedexC 0
| | THENM D 0
| | \
| | ⊢ Assoc(|omral(g;r)|;λx,y.x ** y)
| | |
1 2 BY (BLemma 'omral_times_assoc' ...)
| | \
| | ⊢ Ident(|omral(g;r)|;λx,y.x ** y;11)
| | |
1 BY (AGenRepD ["basic"])
| | THENM Force '5' (Reduce 0)
| | THENM Backchain

```

```

|      'omral_times_ident_l omral_times_ident_r' ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| ⊢ BiLinear(|omral_alg(g;r)|;+omral_alg(g;r);xomral_alg(g;r))
| |
1 BY RWH AbRedexC 0
|   THENM (BLemma 'omral_bilinear' ...a)
|\
| 1. g: OCMon
| 2. r: CRng
| 3. a: |r|
| ⊢ Dist1op2opLR(|omral_alg(g;r)|;·omral_alg(g;r) a;xomral_alg(g;r))
| |
1 BY RWH AbRedexC 0
|   THENM Force '5' (Eval 'dist_1op_2op_lr' 0)
|   THENM (Backchain 'omral_action_times_r1 omral_action_times_r2' ...a)
\
  1. g: OCMon
  2. r: CRng
  ⊢ Comm(|omral_alg(g;r)|;xomral_alg(g;r))
  |
  BY (ARepD ["basic"]
      THENM RWH AbRedexC 0
      THENM Force '2' (Reduce 0)
      THENM BLemma 'omral_times_comm_a' ...a)
*T omral_inj_mon_op 101.8 sec.
⊢ ∀g:OCMon. ∀r:CRng. ∀k,k':|g|. inj(k * k',1) = inj(k,1) ** inj(k',1)
|
BY (RepD THENM BLemma 'omral_lookups_same_a'
|   THENM D 0 ...a)
|
1. g: OCMon
2. r: CRng
3. k: |g|
4. k': |g|
5. u: |g|
⊢ inj(k * k',1)[u] = (inj(k,1) ** inj(k',1))[u]
|
BY (RWW "lookup_omral_times lookup_omral_inj" 0
|   THENM Folds 'rng_when rng_mssum' 0 ...a)
|
⊢ when (k * k') =b u. 1
| = (∑x ∈ dom(inj(k,1)).
|   ∑y ∈ dom(inj(k',1)). when (x * y) =b u. (when k =b x. 1) * (when k' =b y. 1))
|
BY (RWH (LemmaC 'omral_dom_inj') 0
|   THENM SplitOnConclITE ...a)
|\
| 6. 1 = 0
| ⊢ when (k * k') =b u. 1
| | = (∑x@0 ∈ 0{g↓oset}.
| |   ∑y ∈ 0{g↓oset}. when (x@0 * y) =b u. (when k =b x@0. 1) * (when k' =b y. 1))
| |
1 BY Unfold 'rng_mssum' 0
| | THEN Reduce 0
| |

```

```

| ⊢ when (k * k') =b u. 1 = 0
| |
1 BY (RWW "6 rng_when_of_zero" 0 ...)
\
6. ¬(1 = 0)
| ⊢ when (k * k') =b u. 1
|   = (∑x ∈ mset_inj{g↓oset}(k).
|     ∑y ∈ mset_inj{g↓oset}(k'). when (x * y) =b u. (when k =b x. 1) * (when k' =b y. 1))
|
BY (RWW "mset_for_mset_inj" 0 ...a)
|
| ⊢ when (k * k') =b u. 1 = when (k * k') =b u. (when k =b k. 1) * (when k' =b k'. 1)
|
BY (Unfold 'rng_when' 0
|   THEN RNWs [3;4] (LemmaC 'mon_when_true') 0 ...a)
|\
| | ⊢ ↑(k =b k)
| |
1 BY (RW bool_to_propC 0 ...)
|\
| | ⊢ ↑(k' =b k')
| |
1 BY (RW bool_to_propC 0 ...)
\
| ⊢ when (k * k') =b u. 1 = when (k * k') =b u. 1 * 1
|
|   BY (RW RngNormC 0 ...)
*D omral_alg_umap_df
      alg_umap{<g:mon:*>,<a:rng:*>}(<n:alg:*>,<f:mon->alg:*>)
      == omral_alg_umap{<g>; <a>; <n>; <f>}
      alg_umap(<n:alg:*>,<f:mon->alg:*>)== omral_alg_umap{<g>; <a>; <n>; <f>}
*A omral_alg_umap      alg_umap(n,f) == λps:|omral(g;a)|. ∑k ∈ dom(ps). ps[k] ·n (f k)
*T omral_alg_umap_wf 16.0 sec.
| ⊢ ∀g:OCMon. ∀a:CRng. ∀n:a-Algebra. ∀f:|g| → |n|. alg_umap(n,f) ∈ |omral(g;a)| → |n|
|
BY (Unfold 'omral_alg_umap' 0 ...)
*T omral_alg_umap_is_hom 496.0 sec.
| ⊢ ∀g:OCMon. ∀a:CRng. ∀n:a-Algebra. ∀f:MonHom(g,n↓rg↓xmn).
|   IsAlgHom{a,omral_alg(g;a),n}(alg_umap(n,f))
|
BY (AGenRepD ["compound";"basic"] ...a)
|\
| 1. g: OCMon
| 2. a: CRng
| 3. n: a-Algebra
| 4. f: MonHom(g,n↓rg↓xmn)
| 5. a1: |omral_alg(g;a)|
| 6. a2: |omral_alg(g;a)|
| ⊢ alg_umap(n,f) (a1 +omral_alg(g;a) a2) = (alg_umap(n,f) a1) +n (alg_umap(n,f) a2)
| |
1 BY All (RW (HigherC AbRedexC))
| | THENM Force '5' (Eval 'omral_alg_umap' 0)
| |
| 5. a1: |omral(g;a)|
| 6. a2: |omral(g;a)|
| ⊢ (∑k ∈ dom(a1 ++ a2). (a1 ++ a2)[k] ·n (f k))
| | = (∑k ∈ dom(a1). a1[k] ·n (f k)) +n (∑k ∈ dom(a2). a2[k] ·n (f k))

```



```

| |
1 BY % First equalize summation domains %
| | (RWH (LemmaWithC ['q', 'dom(a1) ∪ dom(a2)']
| | 'rng_mssum_dom_shift') 0 ...a)
| | \
| | ⊢ ↑(dom(a1 ++ a2) ⊆b dom(a1) ∪ dom(a2))
| | |
1 2 BY (BLemma 'omral_plus_dom' ...)
| | \
| | 7. x: |(g↓oset)|
| | 8. ↑(x ∈b (dom(a1) ∪ dom(a2)) - dom(a1 ++ a2))
| | ⊢ (a1 ++ a2)[x] ·n (f x) = 0
| | |
1 2 BY (RWW "lookup_omral_eq_zero.2" 0 ...a)
| | | \
| | | ⊢ ¬↑(x ∈b dom(a1 ++ a2))
| | | |
1 2 3 BY (RWW "fset_mem_union mset_mem_diff" 8
| | | THENM RW bool_to_propC 8
| | | THENM ProveProp ...)
| | | \
| | | ⊢ 0 ·n (f x) = 0
| | | |
1 2 BY (RWO "module_act_zero_1" 0
| | | THENM Reduce 0 ...)
| | | \
| | | ⊢ ↑(dom(a1) ⊆b dom(a1) ∪ dom(a2))
| | | |
1 2 BY (RWW "mem_bsubset fset_mem_union" 0
| | | THENM RepD
| | | THENM RW bool_to_propC 0
| | | THENM ProveProp ...a)
| | | \
| | | 7. x: |(g↓oset)|
| | | 8. ↑(x ∈b (dom(a1) ∪ dom(a2)) - dom(a1))
| | | ⊢ a1[x] ·n (f x) = 0
| | | |
1 2 BY (RWW "lookup_omral_eq_zero module_act_zero_1" 0
| | | THENM Reduce 0 ...)
| | | |
| | | ⊢ ¬↑(x ∈b dom(a1))
| | | |
1 2 BY (RWW "fset_mem_union mset_mem_diff" 8
| | | THENM RW bool_to_propC 8
| | | THENM ProveProp ...)
| | | \
| | | ⊢ ↑(dom(a2) ⊆b dom(a1) ∪ dom(a2))
| | | |
1 2 BY (RWW "mem_bsubset fset_mem_union" 0
| | | THENM RepD
| | | THENM RW bool_to_propC 0
| | | THENM ProveProp ...a)
| | | \
| | | 7. x: |(g↓oset)|
| | | 8. ↑(x ∈b (dom(a1) ∪ dom(a2)) - dom(a2))
| | | ⊢ a2[x] ·n (f x) = 0
| | | |

```

```

1 2 BY (RWW "lookup_omral_eq_zero module_act_zero_1" 0
| | | THENM Reduce 0 ...)
| | |
| | |  $\neg \uparrow (x \in_b \text{dom}(a2))$ 
| | |
1 2 BY (RWW "fset_mem_union mset_mem_diff" 8
| | THENM RW bool_to_propC 8
| | THENM ProveProp ...)
| \
|  $\vdash (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). (a1 ++ a2)[k] \cdot n (f k))$ 
|  $= (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). a1[k] \cdot n (f k)) + n (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). a2[k] \cdot n (f k))$ 
| |
1 BY (RWW "lookup_omral_plus
| | module_act_plus.1" 0 ...a)
| |
|  $\vdash (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). (a1[k] \cdot n (f k)) + n (a2[k] \cdot n (f k)))$ 
|  $= (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). a1[k] \cdot n (f k)) + n (\sum k \in \text{dom}(a1) \cup \text{dom}(a2). a2[k] \cdot n (f k))$ 
| |
1 BY RWH rng_of_algC 0
| THENM (RWW "rng_mssum_of_plus<" 0 ...)
|\
| 1. g: OCMon
| 2. a: CRng
| 3. n: a-Algebra
| 4. f: MonHom(g,n|rg|xmn)
| 5. u: |a|
| 6. a@0: |omral_alg(g;a)|
|  $\vdash \text{alg\_umap}(n,f) (\cdot \text{omral\_alg}(g;a) u a@0) = \cdot n u (\text{alg\_umap}(n,f) a@0)$ 
| |
1 BY RenameVar 'a1' 6
| | THENM All (RW (HigherC AbRedexC))
| | THENM Force '5' (Eval 'omral_alg_umap' 0)
| |
| 6. a1: |omral(g;a)|
|  $\vdash (\sum k \in \text{dom}(u \cdot a1). (u \cdot a1)[k] \cdot n (f k)) = \cdot n u (\sum k \in \text{dom}(a1). a1[k] \cdot n (f k))$ 
| |
1 BY % Equalize summation domains %
| | (Unfold 'rng_mssum' 0
| | THENM RWO "omral_dom_action" 0
| | THENM Fold 'rng_mssum' 0 ...a)
| | \
| | 7. x: |(g|oset)|
| | 8.  $\uparrow (x \in_b \text{dom}(a1) - \text{dom}(u \cdot a1))$ 
| |  $\vdash (u \cdot a1)[x] \cdot n (f x) = e$ 
| | |
1 2 BY Reduce 0
| | |
| | |  $(u \cdot a1)[x] \cdot n (f x) = 0n$ 
| | |
1 2 BY (RWW "lookup_omral_eq_zero module_act_zero_1" 0 ...)
| | |
| | |  $\neg \uparrow (x \in_b \text{dom}(u \cdot a1))$ 
| | |
1 2 BY (RWW "fset_mem_union mset_mem_diff" 8
| | THENM RW bool_to_propC 8
| | THENM ProveProp ...)
| \

```

```

|   ⊢ (∑k ∈ dom(a1). (u · a1)[k] ·n (f k)) = ·n u (∑k ∈ dom(a1). a1[k] ·n (f k))
|   |
1  BY (RWH (LemmaC 'lookup_omral_action' 0 ...a)
|   |
|   ⊢ (∑k ∈ dom(a1). (u * a1[k]) ·n (f k)) = ·n u (∑k ∈ dom(a1). a1[k] ·n (f k))
|   |
1  BY % n.act is grp hom... Type matching can't figure out n binding:
|   |   because matching requires fill in of forgetful functor. %
|   |
|   |   (RWH (LemmaWithC ['n', [n↓rg↓+gp]] 'dist_hom_over_mset_for' 0
|   |     THENM Fold 'rng_mssum' 0 ...a)
|   | \
|   | ⊢ ·n u ∈ MonHom(n↓rg↓+gp, n↓rg↓+gp)
|   | |
1  2 BY (Fold 'grp_of_module' 0
|   |   THENM MemTypeCD
|   |   THEN IfLabL ['set predicate',
|   |     BLemma 'module_act_is_grp_hom' ] ...)
|   | \
|   | ⊢ (∑k ∈ dom(a1). (u * a1[k]) ·n (f k)) = (∑k ∈ dom(a1). ·n u (a1[k] ·n (f k)))
|   | |
1  BY (RWW "module_action_p.1" 0 ...a)
|   | |
|   | ⊢ (∑k ∈ dom(a1). u ·n (a1[k] ·n (f k))) = (∑k ∈ dom(a1). ·n u (a1[k] ·n (f k)))
|   | |
1  BY (Unfold 'infix_ap' 0 ...)
| \
| 1. g: OCMon
| 2. a: CRng
| 3. n: a-Algebra
| 4. f: MonHom(g, n↓rg↓xmn)
| 5. a1: |omral_alg(g;a)|
| 6. a2: |omral_alg(g;a)|
| ⊢ alg_umap(n,f) (a1 xomral_alg(g;a) a2) = (alg_umap(n,f) a1) xn (alg_umap(n,f) a2)
| |
1 BY All (RW (HigherC AbRedexC))
| | THEN Force '5' (Eval 'omral_alg_umap' 0)
| | THEN RWH rng_to_mod_mssumC 0
| |
| 5. a1: |omral(g;a)|
| 6. a2: |omral(g;a)|
| ⊢ (∑n k ∈ dom(a1 ** a2). (a1 ** a2)[k] ·n (f k))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| |
1 BY % Suitably widen summation domain %
| | (Unfold 'mod_mssum' 0
| |   THENM RWO "omral_times_dom" 0
| |   THENM Fold 'mod_mssum' 0 ...a)
| | \
| | 7. x: |(g↓oset)|
| | 8. ↑(x ∈b (dom(a1) × dom(a2)) - dom(a1 ** a2))
| | ⊢ (a1 ** a2)[x] ·n (f x) = e
| | |
1 2 BY (RWW "lookup_omral_eq_zero" 0 ...a)
| | | \
| | | ⊢ ¬↑(x ∈b dom(a1 ** a2))
| | | |

```

```

1 2 3 BY (RWW "fset_mem_union mset_mem_diff" 8
| | | THENM RW bool_to_propC 8
| | | THENM ProveProp ...)
| | \
| | ⊢ 0 ·n (f x) = e
| | |
1 2 BY (Reduce 0
| | THENM RWW "module_act_zero_1" 0 ...)
| | \
| | ⊢ (∑n k ∈ dom(a1) × dom(a2). (a1 ** a2)[k] ·n (f k))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| | |
1 BY (RWW "lookup_omral_times" 0
| | THENM Fold 'rng_mssum' 0 ...a)
| | |
| | ⊢ (∑n k ∈ dom(a1) × dom(a2)
| | (∑x ∈ dom(a1). ∑y ∈ dom(a2). when (x * y) =b k. a1[x] * a2[y]) ·n (f k))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| | |
1 BY % Pull 1st ∑n k and when together.
| | Would be nice to automate this with some metric-guided rewrites.
| | Didn't Bundy write a CADE paper on this a while ago?
| | %
| | (RWD (PolyC "mod_action_mssum_r mod_action_when_r") 0 ...a)
| | |
| | ⊢ (∑n k ∈ dom(a1) × dom(a2)
| | ∑n x ∈ dom(a1). ∑n y ∈ dom(a2). when (x * y) =b k. (a1[x] * a2[y]) ·n (f k))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| | |
1 BY (RWD (LemmaC 'mod_mssum_swap') 0
| | THENM RWH (LemmaC 'grp_eq_sym') 0 ...a)
| | |
| | ⊢ (∑n x ∈ dom(a1)
| | ∑n y ∈ dom(a2). ∑n k ∈ dom(a1) × dom(a2). when k =b (x * y). (a1[x] * a2[y]) ·n (f k))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| | |
1 BY (RWN 3 (UnfoldTopC 'mod_mssum') 0
| | THENM RWH dset_of_monC 0
| | THENM Unfold 'oset_of_ocmon' 0
| | THENM RWH (LemmaC 'fset_for_when_eq') 0 ...a)
| | |
| | \
| | 7. x: |(g↓set)|
| | 8. ↑(x ∈b dom(a1))
| | 9. x1: |(g↓set)|
| | 10. ↑(x1 ∈b dom(a2))
| | ⊢ ↑(x * x1 ∈b dom(a1) × dom(a2))
| | |
1 2 BY (BLemma 'prod_in_mset_prod' ...)
| | \
| | ⊢ (∑n x ∈ dom(a1). ∑n y ∈ dom(a2). (a1[x] * a2[y]) ·n (f (x * y)))
| | = (∑n k ∈ dom(a1). a1[k] ·n (f k)) xn (∑n k ∈ dom(a2). a2[k] ·n (f k))
| | |
1 BY (RenameBVars' ['k', 'x'; 'k', 'y'] 0
| | THENM RWH (LemmaC 'mod_times_mssum_r') 0
| | THENM RWH (LemmaC 'mod_times_mssum_l') 0 ...a)
| | |
| | ⊢ (∑n x ∈ dom(a1). ∑n y ∈ dom(a2). (a1[x] * a2[y]) ·n (f (x * y)))

```

```

|      | = (∑n x ∈ dom(a1). ∑n y ∈ dom(a2). (a1[x] ·n (f x)) xn (a2[y] ·n (f y)))
|      |
1      | BY (RWH (LemmaC 'monoid_hom_op') 0
|      |   THENM Reduce 0 ...a)
|      |
|      | ⊢ (∑n x ∈ dom(a1). ∑n y ∈ dom(a2). (a1[x] * a2[y]) ·n (xn (f x) (f y)))
|      | = (∑n x ∈ dom(a1). ∑n y ∈ dom(a2). (a1[x] ·n (f x)) xn (a2[y] ·n (f y)))
|      |
1      | BY (RWW "algebra_act_times_lr" 0 ...)
|
| \
| 1. g: OCMon
| 2. a: CRng
| 3. n: a-Algebra
| 4. f: MonHom(g,n↓rg↓xmn)
| ⊢ alg_umap(n,f) 1omral_alg(g;a) = 1n
|
| BY All (RW (HigherC AbRedexC))
| | THENM Force '5' (Eval 'omral_alg_umap' 0)
| |
| ⊢ (∑k ∈ dom(11). 11[k] ·n (f k)) = 1n
|
| BY Unfold 'omral_one' 0
|
| ⊢ (∑k ∈ dom(inj(e,1)). inj(e,1)[k] ·n (f k)) = 1n
|
| BY (RWH (LemmaC 'omral_dom_inj') 0
| | THENM SplitOnConclITE ...a)
| \
| | 5. 1 = 0
| | ⊢ (∑k ∈ 0{g↓oset}. inj(e,1)[k] ·n (f k)) = 1n
| |
| 1 BY Eval 'rng_mssum' 0
| |
| | ⊢ 0n = 1n
| |
| 1 BY (InvertRel 0 THENM BLemma 'module_over_triv_rng' ...)
| \
| 5. ¬(1 = 0)
| ⊢ (∑k ∈ mset_inj{g↓oset}(e). inj(e,1)[k] ·n (f k)) = 1n
|
| BY (Unfold 'rng_mssum' 0
| | THENM RWH (LemmaC 'mset_for_mset_inj') 0 ...a)
| |
| ⊢ inj(e,1)[e] ·n (f e) = 1n
|
| BY (RWH (LemmaC 'lookup_omral_inj') 0 ...a)
|
| ⊢ (when e =b e. 1) ·n (f e) = 1n
|
| BY (RWH (LemmaC 'mon_when_true') 0 ...a)
| \
| | ⊢ ↑(e =b e)
| |
| 1 BY (RW bool_to_propC 0 ...)
| \
| ⊢ 1 ·n (f e) = 1n
|

```

```

      BY (RWW "module_action_p.2 monoid_hom_id" 0
          THENM Reduce 0 ...)
*T omral_alg_umap_tri_comm 48.8 sec.
⊢ ∀g:OCMon. ∀a:CRng. ∀n:a-Algebra. ∀f:|g| → |n|. alg_umap(n,f) o (λk.inj(k,1)) = f
|
BY (RepD THENM New ['k1'] Ext
  | THENM Eval ``omral_alg_umap`` 0 ...a)
|
1. g: OCMon
2. a: CRng
3. n: a-Algebra
4. f: |g| → |n|
5. k1: |g|
⊢ (Σk ∈ dom(inj(k1,1)). inj(k1,1)[k] ·n (f k)) = f k1
|
BY (RWO "omral_dom_inj" 0 THENM SplitOnConclITE
  | THENM Eval ``rng_mssum`` 0 ...a)
| \
| 6. 1 = 0
| ⊢ 0n = f k1
| |
1 BY (InvertRel 0 THENM BLemma 'module_over_triv_rng' ...)
| \
| 6. ¬(1 = 0)
| ⊢ msFor{n↓rg↓+gp} k ∈ mset_inj{g↓oset}(k1). inj(k1,1)[k] ·n (f k) = f k1
|
BY (RWW "mset_for_mset_inj" 0 ...a)
|
⊢ inj(k1,1)[k1] ·n (f k1) = f k1
|
BY (RWW "lookup_omral_inj mon_when_true" 0 ...a)
| \
| ⊢ ↑(k1 =b k1)
| |
1 BY (RW bool_to_propC 0 ...)
| \
| ⊢ 1 ·n (f k1) = f k1
|
  BY (RWW "module_action_p.2" 0 ...)
*T omral_alg_umap_unique 148.2 sec.
⊢ ∀g:OCMon. ∀a:CRng. ∀n:a-Algebra. ∀f:|g| → |n|. ∀f':a-AlgebraHom(omral_alg(g;a);n).
  | f' o (λk:|g|. inj(k,1)) = f ⇒ f' = alg_umap(n,f)
|
BY (RepD THENM New ['ps'] Ext
  | THENM Eval ``omral_alg_umap`` 0 ...a)
|
1. g: OCMon
2. a: CRng
3. n: a-Algebra
4. f: |g| → |n|
5. f': a-AlgebraHom(omral_alg(g;a);n)
6. f' o (λk:|g|. inj(k,1)) = f
7. ps: |omral(g;a)|
⊢ f' ps = (Σk ∈ dom(ps). ps[k] ·n (f k))
|
BY (RWO "<" 0 THENM Reduce 0 ...a)
|

```

```

⊢ f' ps = (∑k ∈ dom(ps). ps[k] ·n (f' inj(k,1)))
|
BY (Unfold 'rng_mssum' 0
| THENM Fold 'grp_of_module' 0
| THENM RWO "module_hom_action<" 0 ...a)
|
⊢ f' ps = msFor{n|grp} k ∈ dom(ps). f' (·omral_alg(g;a) ps[k] inj(k,1))
|
BY (RWH (RevLemmaWithC ['m', 'omral_alg(g;a)↓grp'] 'dist_hom_over_mset_for') 0 ...a)
| \
| ⊢ f' ∈ MonHom(omral_alg(g;a)↓grp, n|grp)
| |
1 BY (AddAllProperties 5
| | THENM MemTypeCD
| | THEN IfLabL ['set predicate',
| | BLemma 'module_hom_is_grp_hom']
| | THENM AGenRepD ["compound"; "basic"]
| | THENM HypBackchain ...a)
| | \
| | 6. ∀a1,a2:|omral_alg(g;a)|. f' (a1 +omral_alg(g;a) a2) = (f' a1) +n (f' a2)
| | 7. ∀u:|a|. fun_thru_top(|omral_alg(g;a)|; |n|; ·omral_alg(g;a) u; ·n u; f')
| | 8. ∀a1,a2:|omral_alg(g;a)|. f' (a1 xomral_alg(g;a) a2) = (f' a1) xn (f' a2)
| | 9. f' 1omral_alg(g;a) = 1n
| | 10. f' o (λk:|g|. inj(k,1)) = f
| | 11. ps: |omral(g;a)| | | |
| | ⊢ f' ∈ |(omral_alg(g;a)↓grp)| → |(n|grp)|
| | |
1 2 BY % Inclusion should be fixed to get this %
| | (RWH AbRedexC 0 ...)
| | \
| | 6. ∀a1,a2:|omral_alg(g;a)|. f' (a1 +omral_alg(g;a) a2) = (f' a1) +n (f' a2)
| | 7. ∀u:|a|. fun_thru_top(|omral_alg(g;a)|; |n|; ·omral_alg(g;a) u; ·n u; f')
| | 8. ∀a1,a2:|omral_alg(g;a)|. f' (a1 xomral_alg(g;a) a2) = (f' a1) xn (f' a2)
| | 9. f' 1omral_alg(g;a) = 1n
| | 10. f' o (λk:|g|. inj(k,1)) = f
| | 11. ps: |omral(g;a)|
| | 12. u: |a|
| | 13. a@0: |omral_alg(g;a)|
| | ⊢ f' (·omral_alg(g;a) u a@0) = ·n u (f' a@0)
| | |
1 BY (Unfold 'fun_thru_top' 7
| THENM HypBackchain ...)
| \
| ⊢ f' ps = f' (msFor{omral_alg(g;a)↓grp} k ∈ dom(ps). ·omral_alg(g;a) ps[k] inj(k,1))
|
BY (EqCD ...)
|
⊢ ps = msFor{omral_alg(g;a)↓grp} k ∈ dom(ps). ·omral_alg(g;a) ps[k] inj(k,1)
|
BY RWH AbRedexC 0
| THENM Force '5' (Reduce 0)
|
⊢ ps = msFor{omral_alg(g;a)↓grp} k ∈ dom(ps). ps[k] .. inj(k,1)
|
BY (RWH "omral_action_inj" 0 ...)
|
⊢ ps = msFor{omral_alg(g;a)↓grp} k ∈ dom(ps). inj(k, ps[k] * 1)

```

```

|
BY (RW RngNormC 0
    THENM RWW "omral_fact_a<" 0 ...)
*D omral_fma_df          omral_fma(<g:g:*>;<a:a:*>)== omral_fma{<g>; <a>}
*A omral_fma    omral_fma(g;a) == <omral_alg(g;a), λk.inj(k,1), λn,f.alg_umap(n,f)>
*T omral_fma_wf  22.4 sec.
⊢ ∀g:OCMon. ∀a:CRng. omral_fma(g;a) ∈ FMASig(g;a)
|
BY (Unfolds 'omral_fma fma_sig' 0 ...)
*T omral_fma_wf2  53.0 sec.
⊢ ∀g:OCMon. ∀a:CRng. omral_fma(g;a) ∈ FMonAlg(g;a)
|
BY (RepD THENM MemTypeCD ...)
| \
| 1. g: OCMon
| 2. a: CRng
| ⊢ IsMonHom{g,omral_fma(g;a).alg|rg|xmn}(omral_fma(g;a).inj)
| |
1 BY (AGenRepD ["compound";"basic"]
| |   THENM Force '5' (Reduce 0) ...a)
| | \
| | 3. a1: |g|
| | 4. a2: |g|
| | ⊢ inj(a1 * a2,1) = inj(a1,1) ** inj(a2,1)
| | |
1 2 BY (BLemma 'omral_inj_mon_op' ...)
| | \
| | ⊢ inj(e,1) = 11
| | |
1   BY (Fold 'omral_one' 0 ...)
| \
| 1. g: OCMon
| 2. a: CRng
| 3. n: a-Algebra
| 4. f: MonHom(g,n|rg|xmn)
| ⊢ omral_fma(g;a).umap n f = !f':|omral_fma(g;a).alg| → |n|
|                               IsAlgHom{a,omral_fma(g;a).alg,n}(f')
|                               ∧ f' o omral_fma(g;a).inj = f
|
BY (Unfold 'uni_sat' 0 THEN GenRepD
|   THENM All (\i.Forced '5' (Reduce i)) ...a)
| \
| ⊢ IsAlgHom{a,omral_alg(g;a),n}(alg_umap(n,f))
| |
1 BY (BLemma 'omral_alg_umap_is_hom' ...)
| \
| ⊢ alg_umap(n,f) o (λk.inj(k,1)) = f
| |
1 BY (BLemma 'omral_alg_umap_tri_comm' ...)
| \
| 5. a': |omral(g;a)| → |n|
| 6. IsAlgHom{a,omral_alg(g;a),n}(a')
| 7. a' o (λk.inj(k,1)) = f
| ⊢ a' = alg_umap(n,f)
|
BY (BLemma 'omral_alg_umap_unique' ...)
|

```



```

⊢ a' ∈ a-AlgebraHom(omral_alg(g;a);n)
|
BY % Peculiarities of definitions make this a mess. %
| (ARepD ["compound";"basic"]
| THENM RepeatM MemTypeCD
| THEN IfLabL ['set predicate',AGenRepD ["compound";"basic"]
| THENM HypBackchain] ...)
|
6. ∀a1,a2:|omral_alg(g;a)|. a' (a1 +omral_alg(g;a) a2) = (a' a1) +n (a' a2)
7. ∀u:|a|. fun_thru_1op(|omral_alg(g;a)|;|n|;·omral_alg(g;a) u;·n u;a')
8. ∀a1,a2:|omral_alg(g;a)|. a' (a1 xomral_alg(g;a) a2) = (a' a1) xn (a' a2)
9. a' 1omral_alg(g;a) = 1n
10. a' o (λk.inj(k,1)) = f
11. u: |a|
12. a@0: |omral_alg(g;a)|
⊢ a' (·omral_alg(g;a) u a@0) = ·n u (a' a@0)
|
BY (Unfold 'fun_thru_1op' 7
THENM HypBackchain ...)
*C polynom_3_end *****

```