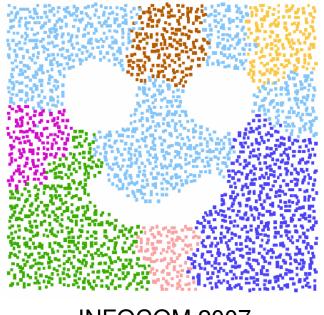
Shape Segmentation and Applications in Sensor Networks

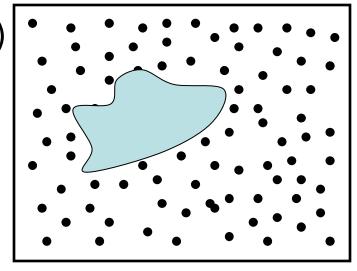
Xianjin Zhu Rik Sarkar Jie Gao





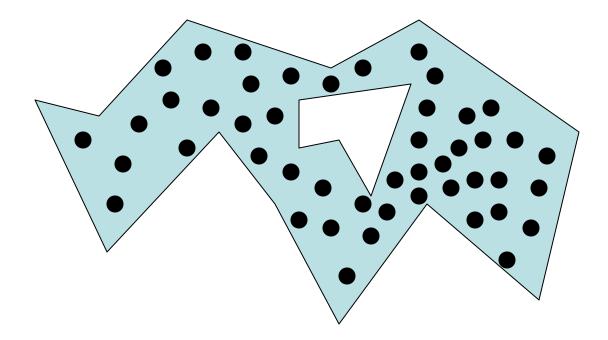
Motivation

- Common assumption: sensors are deployed uniformly randomly inside a simple region (e.g., square).
- In practice, can be complex.
 - Obstacles (lakes, buildings)
 - Terrain variation
 - Degradation over time



Sensor Distribution in Practice

 Nodes are distributed in a geometric region with possible complex shape, with holes.



With holes or a complex shape...

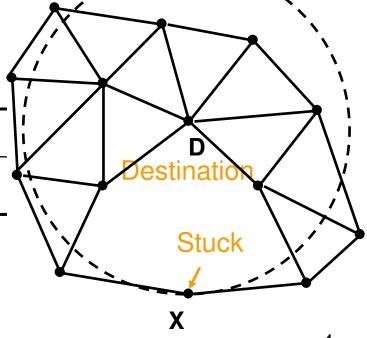
Some protocols may fail:

Greedy forwarding: packets are

greedily forward to the neighbor

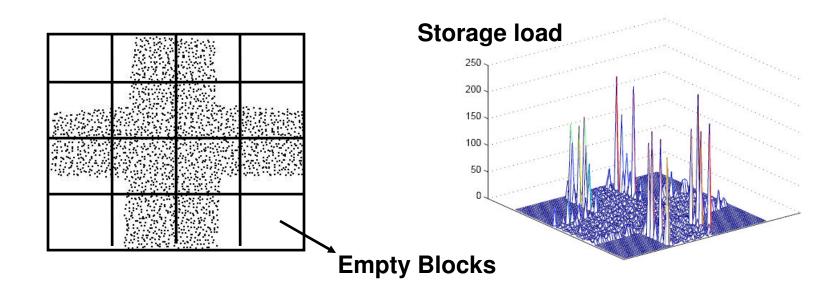
closest to the destination

Dense uniform	Sparse, non-uniform
Works well	May get stuck

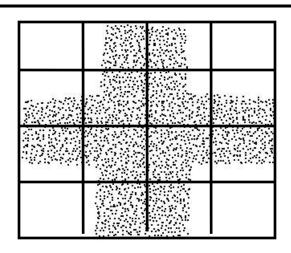


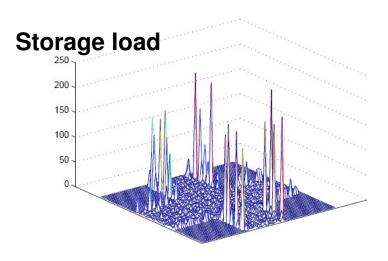
With holes or a complex shape...

- Some protocols have degraded performance
 - Quad-tree type data storage hierarchy
 - Data is hashed uniformly to the quads

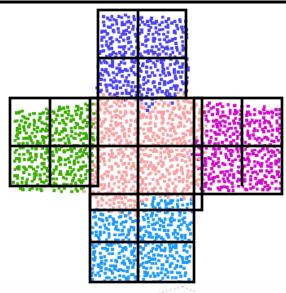


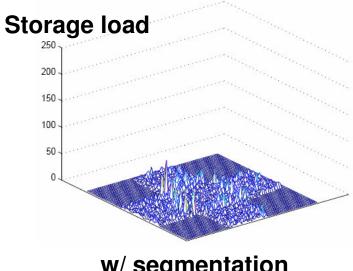
Quad-Tree Type Hierarchy





w/o segmentation



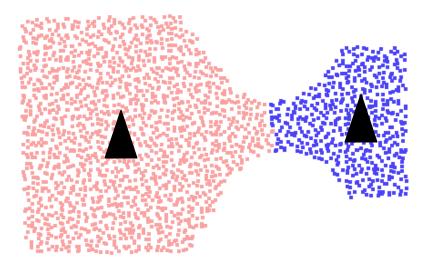


w/ segmentation

Lesson Learned

- Global geometric features affect many aspects of sensor networks.
 - Affect system performance.
 - Affect network design.

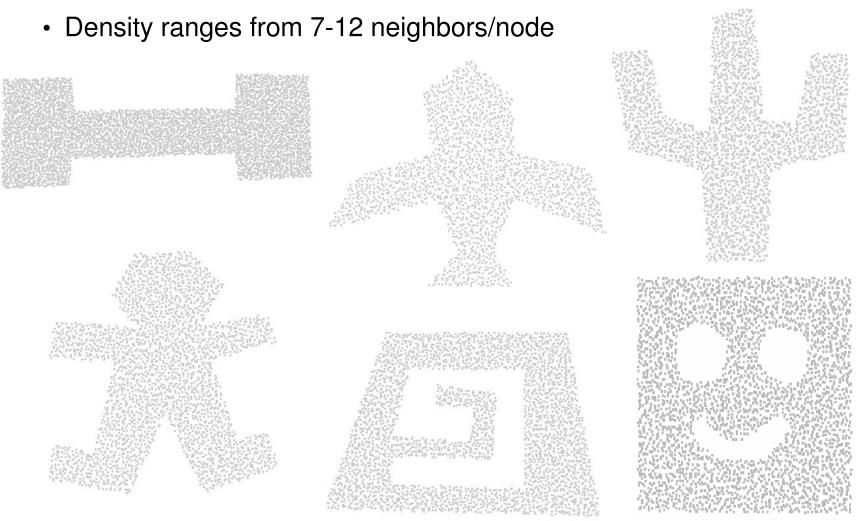
Place base stations and avoid traffic bottleneck



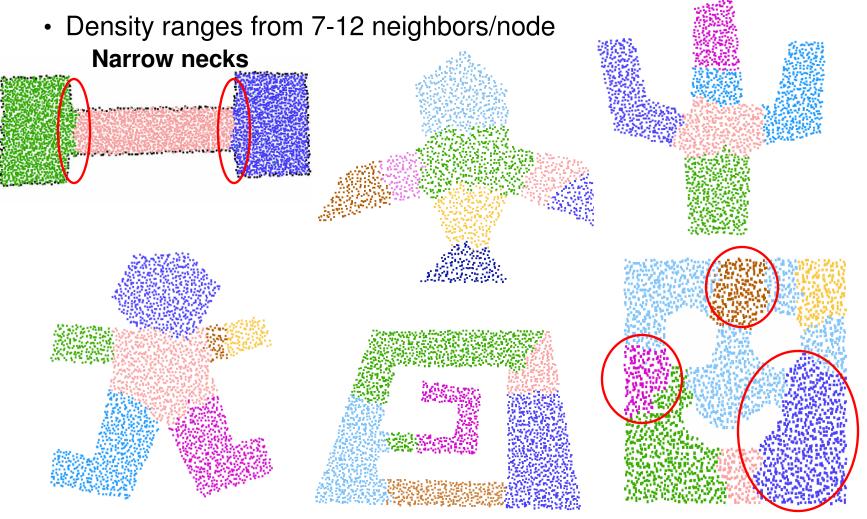
How to Handle Complex Shape?

- Previous work
 - Build problem specific virtual coordinate system (e.g., for routing)
 - Redevelop every algorithm on virtual coordinate system
- Our approach: shape segmentation
 - A unified approach to handle complex geometry
 - Make existing protocols reusable

Sensor Field with Arbitrary Shape



Simulation Results on Segmentation



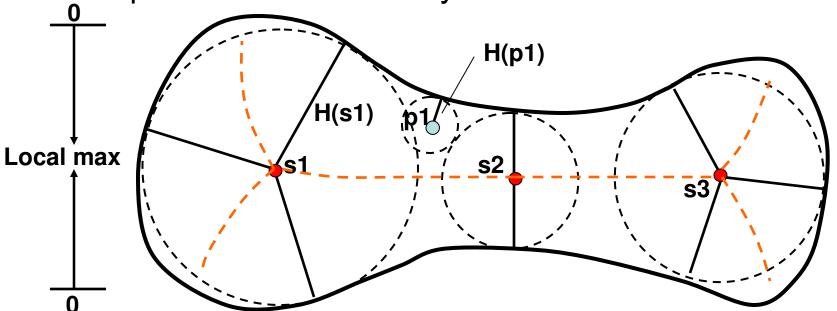
Our Approach: Shape Segmentation

- Segment the irregular field into "nice" pieces.
 - Each piece has no holes, and has a relatively nice shape
- Apply existing algorithms inside each piece.
 - Existing protocols are reusable
- Integrate the pieces together with a problem-dependent structure.

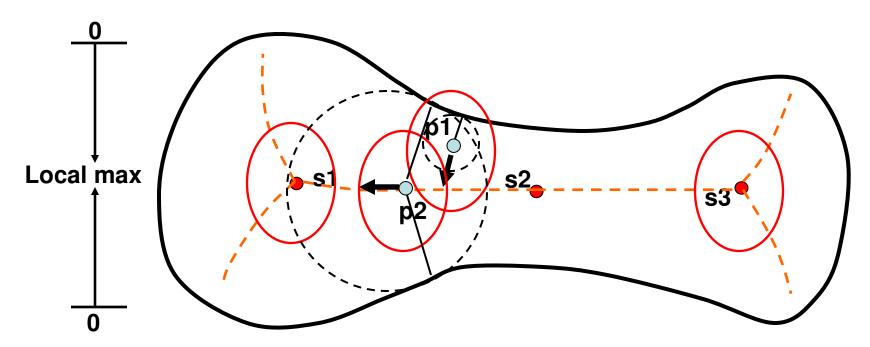
The rest of the talk ...

- Segmentation algorithm
- Implementation issues

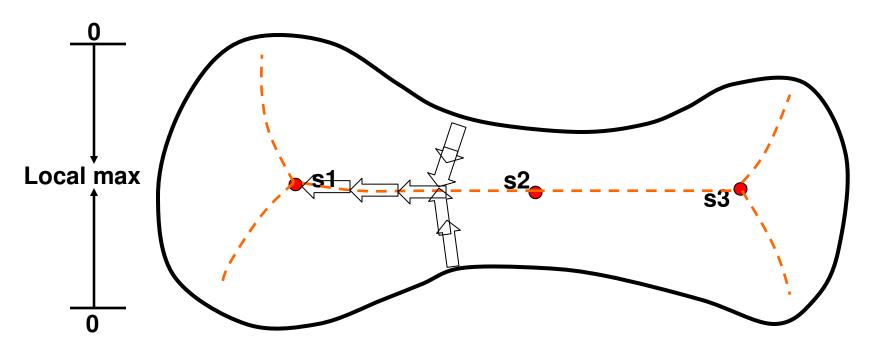
- Flow Complex in continuous domain
 - Distance function $h(x)=min\{||x-p||^2: p on boundary\}$
 - Medial axis: a set of points with at least two closest points on the boundary



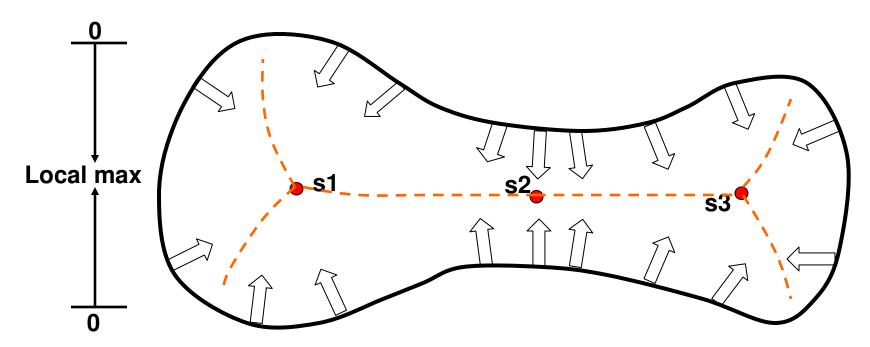
- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



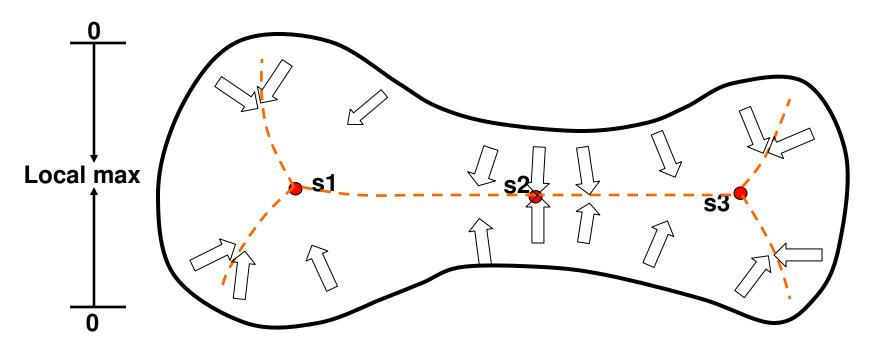
- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



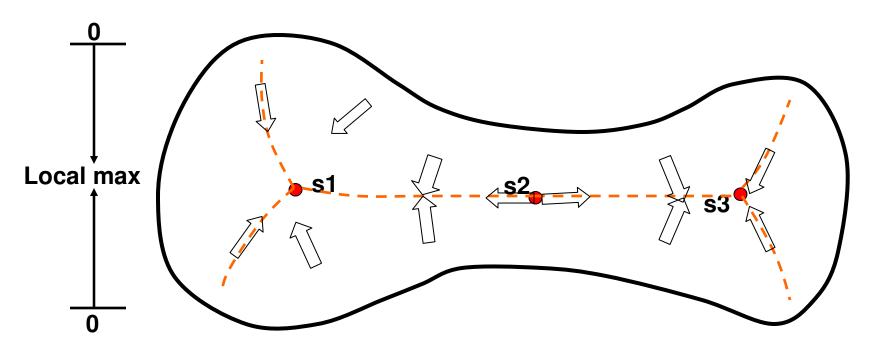
- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



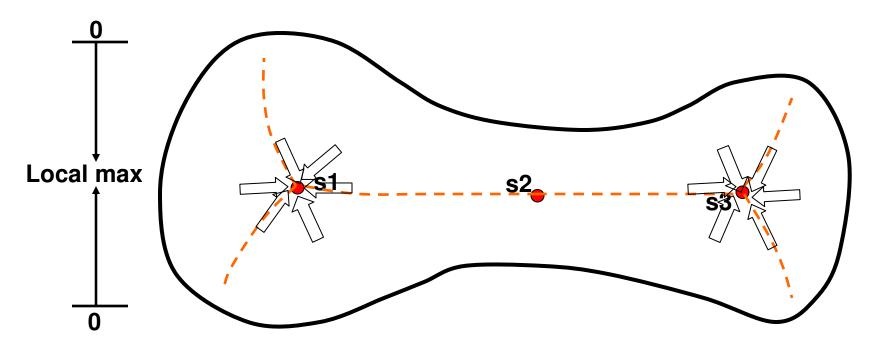
- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)



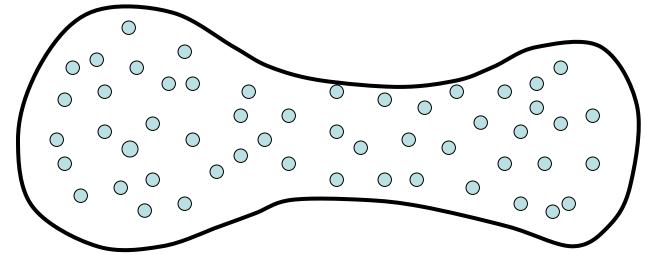
- Flow Complex in Continuous domain
 - Flow direction: the direction that h(x) increases fastest
 - Sinks: local maximum, no flow direction (s1 & s3 here)

- Segments: set of points flow to the same sink

Naturally partition along narrow necks

Implementation Challenges

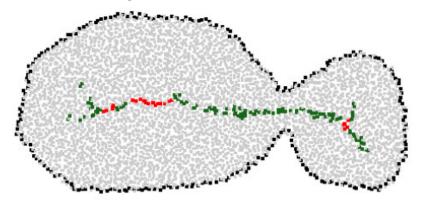
- No global view, no centralized authority
- No location, only connectivity information
 - Distances are approximated by hop count
- Robust to inaccuracy, packet loss, etc.
- Goal: a distributed and robust segmentation algorithm.



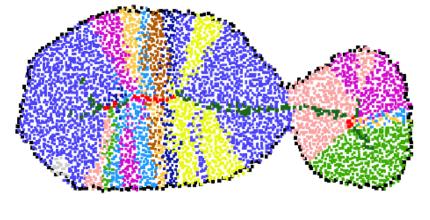
Algorithm Outline

1. Compute the medial axis

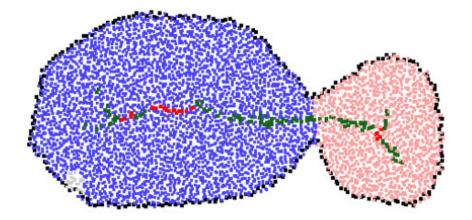
2. Compute the flow

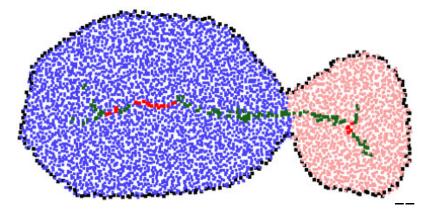


3. Merge nearby sinks



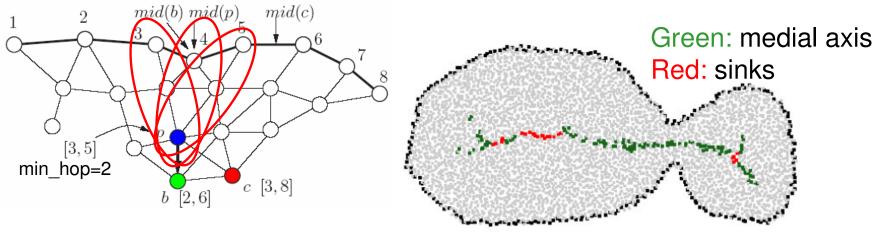
4. Final clean-up





Step 1: Compute the medial axis

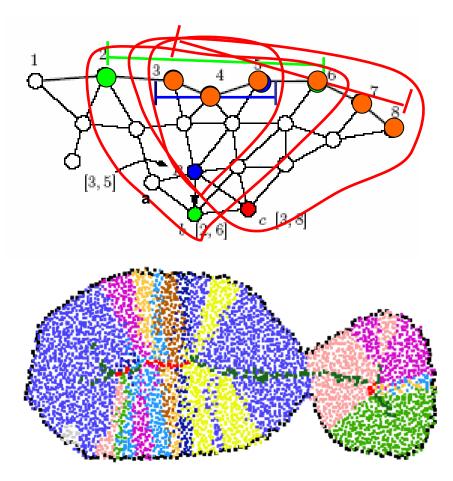
- Boundary nodes flood inward simultaneously.
- Nodes record: minimum hop count & closest intervals on the boundary
- Medial axis: more than two closest intervals



Reference: Boundary Detection [Wang, Gao, Mitchell, MobiCom'06]

Step2: Compute the flow

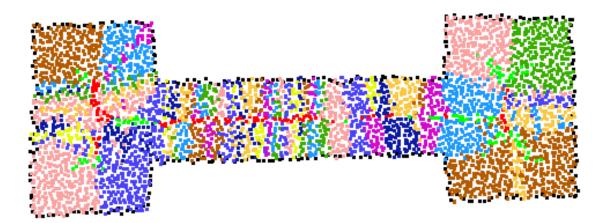
- Flow direction: a pointer to a neighbor with a higher hop count from the same boundary
 - Prefer neighbor with the most symmetric interval
- Sinks must be on the medial axis.
- Network is organized into forests, sinks are roots
- Nodes are classified into segments by their sinks.



Too many segments!

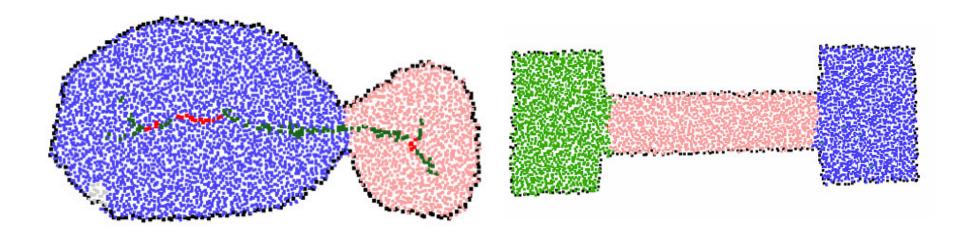
Example of Heavy Fragmentation

 Fragmentation problem becomes severe with parallel boundaries.



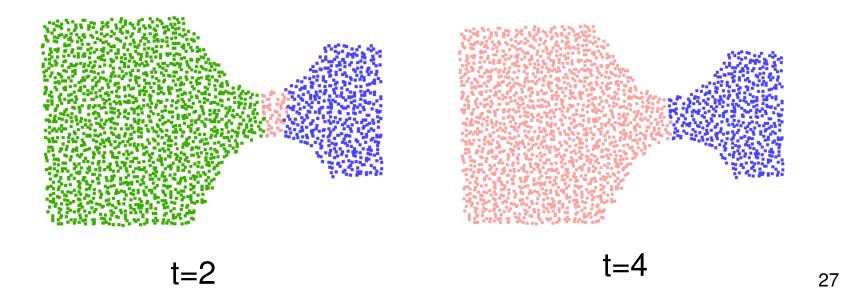
Step3: Merge nearby sinks

 Nearby sinks with similar hop count to the boundaries are merged (together with their segments).



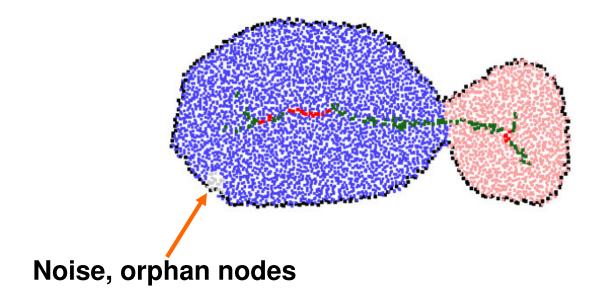
Step3: Merge nearby sinks

- Nearby sinks with similar hop count to the boundaries are merged (together with their segments).
 - Segmentation granularity: |H_{max}-H_{min}|< t

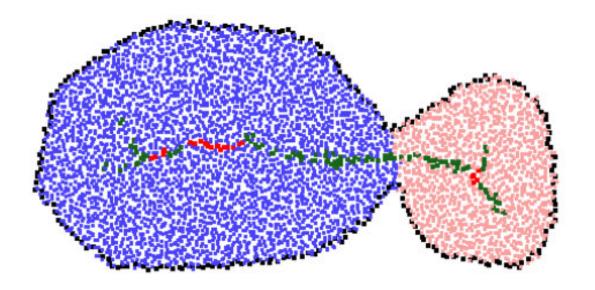


Step4: Final clean-up

- Merge orphan nodes with nearby segments
- Orphan nodes: local maximum and nodes that flow into them

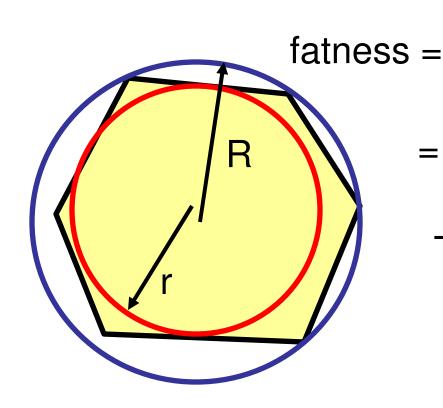


Final Result



Properties of Segmentation

- A few "fat" segments
- Further merging only hurts fatness



max inscribing ball radius

min enclosing ball radius

The bigger the fatter.

Conclusion

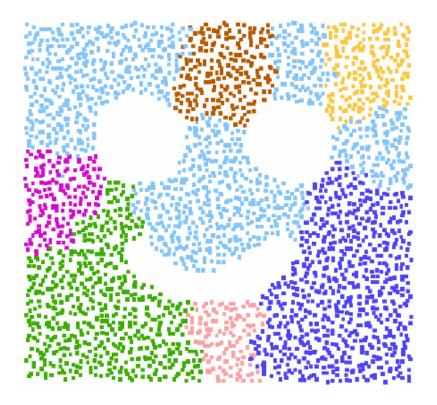
 A unified approach handling complex shape in sensor networks.

 A good example to extract high-level geometry from connectivity information.

Network self-organizes by local operations.

Thank you!

• Questions?



Email: {xjzhu, rik, jgao}@cs.sunysb.edu