



Hardware just-in-time compilation Kimberley Stonehouse, Tom Spink and Björn Franke UK Systems, 18th April 2024

What is just-in-time (JIT) compilation?

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 - Reduces resource usage and inventory cost

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What are the benefits of just-in-time?

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- Performance
 - Accelerating Java, Python, C#, PHP, JavaScript, Ruby, WebAssembly...

What challenges does just-in-time face?

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The primary constraint on just-in-time compilers is speed. The program must not pause during execution.

How is latency mitigated?

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- Tiered compilation
 - Each tier delivers higher quality code, taking longer to do so
 - Low latency early tiers conceal background optimisations

Introducing hardware

- Hardware acceleration
 - Adding hardware accelerated instructions can improve the efficiency of code generation by an average of 15% [1, 2]

Carbon, M.A. et al., 2013. Hardware acceleration for just-in-time compilation on heterogeneous embedded systems. IEEE ASAP'13.
Carbon, M.A. et al., 2014. Hardware acceleration of red-black tree management and application to just-in-time compilation. Journal of Signal Processing Systems.

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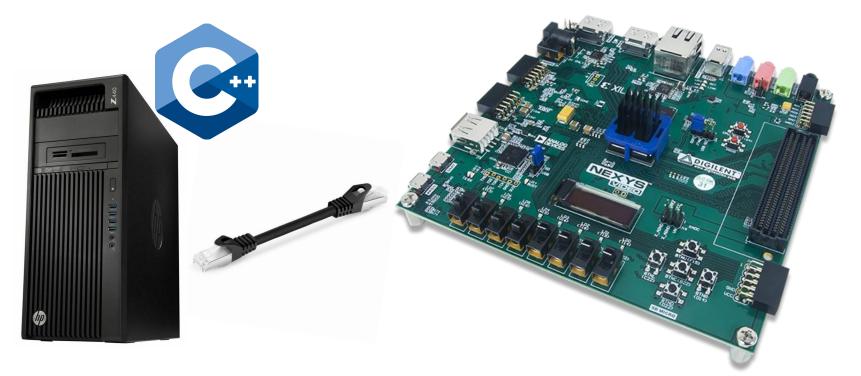
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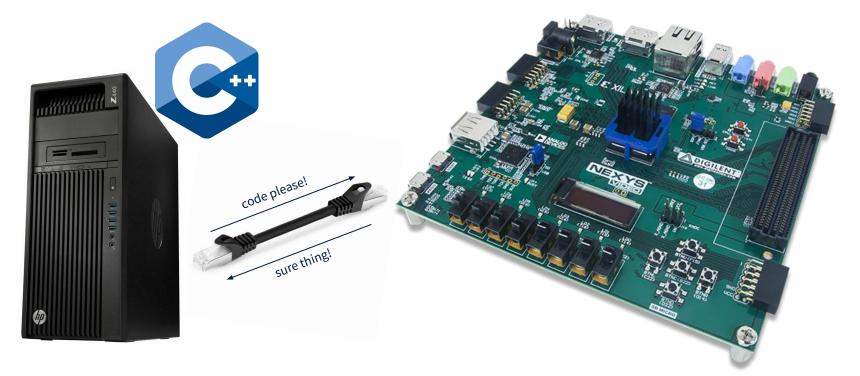
Can we combine tiered compilation with hardware acceleration?

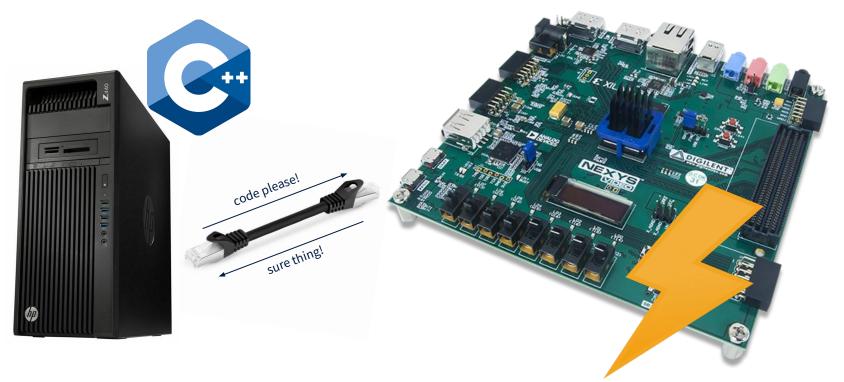
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What about a first tier entirely in hardware?











Can we use high-level synthesis (HLS)?

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 - Synthesise onto a programmable logic device



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 - No double pointers, function pointers or general pointer casting
 - No library calls or standard template library containers

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- Demonstrated **benefits** in multiple contexts
 - Lightweight compilation in constrained environments [1]

[1] Coffin, E. et al., 2020. MicroJIT: a case for templated just-in-time compilation in constrained environments. CASCON'20.
[2] Kaur, H. et al., 2023. Performance Evaluation of Template-based JIT Compilation in OpenJ9. CASCON'23.
[3] Xu, H. et al., 2021. Copy-and-patch compilation: a fast compilation algorithm for high-level languages and bytecode. OOPSLA'21.

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 - Lightweight compilation in constrained environments [1]
 - Reduced start up latency for larger virtual machines [2, 3]

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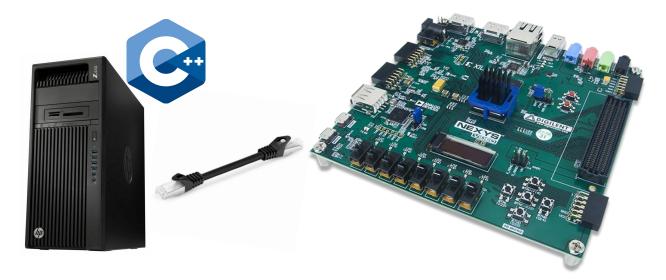
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- Translates register-based intermediate representation to native x86
- Transfers code back and forth via physical Ethernet connection



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- Measure system **performance**
 - Choose appropriate benchmarks for evaluation

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- Hardware just-in-time is a promising space
 - Hardware acceleration has shown positive results
 - Research into an entirely hardware first tier is ongoing
 - But there are still many interesting problems left to solve