<Links> Web programming without tiers

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Team Links



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Coauthors

Cooper, Topor & Yallop

Buneman & Wong

Hejlsberg & Meijer

Hasoya & Pierce

Boag, Chamberlin, Fernandéz, Florescue, Robie, & Siméon

Graunke, Findler, Krishnamurthi & Felleisen

Quenniec

Graham

Armstrong, Virding, Williams, & Wikström

Eich

A Grand Challenge

Design a programming language with a sound basis in theory that becomes the leader in its domain.

A Jolly Good Challenge

Design a programming language with a sound basis in theory that becomes the leader in its domain.

Wadler's theorem of language adoption

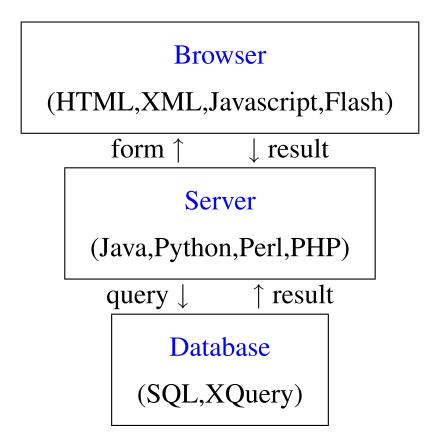
A programming language will be adopted if and only if it permits its users to do something that cannot be done in any other way.

Wadler's theorem of language adoption

A programming language will be adopted if and only if it permits its users to

boldly go where no progamming language has gone before.

Three-tier model



Links builds on successes of functional programming

- Databases Kleisli, LINQ
 Compile Links into SQL, XQuery
- XML Xduce, XQuery
 XML support with regular expression types
- Continuations PLT Scheme, Yahoo stores
 Continuations for web dialogue
- Distribution Erlang, JoCaml
 Reliablity as in Erlang/OTP
- Javascript AJAX
 Compile Links into Javascript

Part I

Kleisli:

Comprehensions for Queries

Comprehensions

Monads and Comprehensions

```
(1) [t \mid ()] = unit t

(2) [t \mid x \leftarrow u] = map(\lambda x. t) u

(3) [t \mid (p,q)] = join[[t \mid q] \mid p]

(1') unit x = [x]

(2') map f xs = [f x \mid x \leftarrow xs]

(3') join xss = [x \mid xs \leftarrow xss, x \leftarrow xs]
```

Monad laws and Comprehension laws

```
(I) 	 join \cdot unit = id
(II) 	 join \cdot map unit = id
(III) 	 join \cdot join = join \cdot map join
(I') 	 [t | (), q] = [t | q]
(II') 	 [t | q, ()] = [t | q]
(III') 	 [t | (p, q), r] = [t | p, (q, r)]
```

Relational Data

BOOKS

title	isbn	year
What Can You Do With a Shoe?	0613733266	1997
Where the Wild Things Are	0060254920	1963

AUTHORS

author	isbn
Beatrice Schenk de Regniers	0613733266
Maurice Sendak	0613733266
Maurice Sendak	0060254920

Relational Query

SQL

```
select b.title, a.author
from BOOKS b, AUTHORS a
where b.isbn = a.isbn
and b.year < 2000</pre>
```

Kleisli (Buneman & Wong & others)

```
{ (title: b.title, author: a.author) |
  \b <--- BOOKS, \a <--- AUTHORS,
  b.isbn = a.isbn, b.year < 2000 }</pre>
```

LINQ (Hejlsberg & Meijer & others)

```
from b in BOOKS
from a in AUTHORS
where b.isbn == a.isbn && b.year < 2000
select new { title=b.title, author=a.author };</pre>
```

An odd relational Query

odd(y)

Kleisli

```
{ (title: b.title, author: a.author) |
    \b <--- BOOKS, \a <--- AUTHORS,
    b.isbn = b.isbn, b.year < 2000, odd(b.year) }

Optimized Kleisli

{ (title: t, author: a) |
    (title: \t, year: \y, author: \a)
    <--- { (title: b.title, author: a.author, year: b.year) |
        \t <--- BOOKS, \a <--- AUTHORS,
        t.isbn = a.isbn, t.year < 2000 }</pre>
```

Kleisli for bioinformatics

```
localblast-blastp (#name: "scop-blast", #db: "scopseq");
localblast-blastp (#name: "pat-blast", #db: "patseq");
scop-add "scop";
setindex-access (#name:"sid2seq", #file: "scopseq",
                 #key: "#sid");
{ (#sf: (#desc: xinfo.#desc.#sf, #hit:x.#accession,
         #pscore:x.#pscore),
  #bridge: (#hit: s, #patent: p.#title, #pscore: p.#pscore))
\mid x \mid x \mid process SEQ using scop-blast, x.#pscore <= PSCORE,
 \xinfo <- process <#sidinfo: x.#accession> using scop,
  \s <- process <#numsid: xinfo.#type.#sf> using scop,
  \y <- process <#key: s> using sid2seq,
  \p <- process y.#seq using pat-blast, p.#pscore <= PSCORE };</pre>
```

Kleisli was first to perform "twelve impossible queries" identified by DoE Workshop for Human Genome Project

XML Data

```
<books>
  <book>
   <title>Where the Wild Things Are</title>
   <author>Maurice Sendak</author>
   <isbn>0060254920</isbn>
   <year>1963
  </book>
 <book>
   <title>What Can You Do With a Shoe?</title>
    <author>Beatrice Schenk de Regniers</author>
    <author>Maurice Sendak</author>
   <isbn>0613733266</isbn>
   <year>1997
  </book>
</books>
```

XML Query

XQuery

```
for $b from input()/books/book
      $a from $b/author
  where $b/year < 2000
  return
    <book>{ $b/title, $a }</book>
Kleisli
  { (title: b.title, author: a) |
    \b <--- BOOKS, \a <-- t.authors,
    b.year < 2000 }
LINQ
  from b in BOOKS
  from a in b.authors
  where b.isbn == a.isbn && b.year < 2000
  select (b.title, a.author)
```

Links

- General-purpose, compiles to SQL or XQuery differs for Kleisli
- No syntactic distinction, one expression may query multiple sources differs from LINQ
- Other related work

Mnesia in Erlang (Mattsson, Nilsson & Wikstrom) Natural Expert (Hutchison, Neuhaus, Schmidt-Schauss & Hall)

Part II

Xduce:

Regular expression types for XML

XML data

Hosoya and Pierce

Xduce types

```
type Addrbook = addrbook[Person*]
type Person = person[Name, Email*, Tel?]
type Name = name[String]
type Email = email[String]
type Tel = tel[String]

type TelBook = telbook[TelPerson*]
type TelPerson = person[Name, Tel]
```

XML Schema

```
<xs:element name="addrbook">
  <xs:complexType>
    <xs:sequence>
      <xs:element name="Person"</pre>
        minOccurs="0" maxOccurs="unbounded"/>
    </xs:sequence>
  </xs:complexType>
</xs:element>
<xs:element name="addrbook">
  <xs:complexType>
    <xs:sequence>
      <xs:element name="name" type="xs:string"/>
      <xs:element name="email" type="xs:string"/>
        minOccurs="0" maxOccurs="unbounded"/>
      <xs:element name="tel" type="xs:string"/>
        minOccurs="0"/>
    </xs:sequence>
  </xs:complexType>
</xs:element>
```

Xduce transformation

```
fun telbook(doc : Addrbook) : TelBook =
 match doc with
    addrbook[val persons as Person*] ->
      telbook[telpersons(persons)]
fun telpersons (val ps as Person*) : TelPerson* =
 match ps with
   person[name[val n as String], Email*,
           tel[val t as String]],
   val rest as Person*
      -> person[name[n], tel[t]],
         telpersons (rest)
  | person[name[val n as String], Email*],
   val rest as Person*
      -> telpersons(rest)
  | ()
      -> ()
```

XQuery transformation



XQuery 1.0: An XML Query Language

W3C Working Draft 16 August 2002

This version:

http://www.w3.org/TR/2002/WD-xquery-20020816/

Latest version:

http://www.w3.org/TR/xquery/

Previous versions:

http://www.w3.org/TR/2002/WD-xquery-20020430/

http://www.w3.org/TR/2001/WD-xquery-20011220/

http://www.w3.org/TR/2001/WD-xquery-20010607/

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XQuery 1.0 and XPath 2.0 Formal Semantics

W3C Working Draft 16 August 2002

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Previous versions:

http://www.w3.org/TR/2002/WD-query-semantics-20020326/

http://www.w3.org/TR/2001/WD-query-semantics-20010607/

http://www.w3.org/TR/2001/WD-query-algebra-20010215/

http://www.w3.org/TR/2000/WD-query-algebra-20001204/

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3.3.2 Matches

Notation

The judgment

Value matches Type

holds when the given value matches the given type.

Semantics

This judgment is specified by the following rules.

The empty sequence matches the empty sequence type.

If two values match two types, then their sequence matches the corresponding sequence type.

Links

• Also use regular expressions for lists

$$A, A?, A+, A*$$

A? corresponds to maybe type of Haskell

• XML syntax that works for cut-and-paste

```
<greet>Hello, \{x\}< vs. <greet>'Hello, ', x</>
```

• Other related work

Xtatic (Gapayev, Levin & Pierce)

Cduce (Benzaken, Castagna & Frisch)

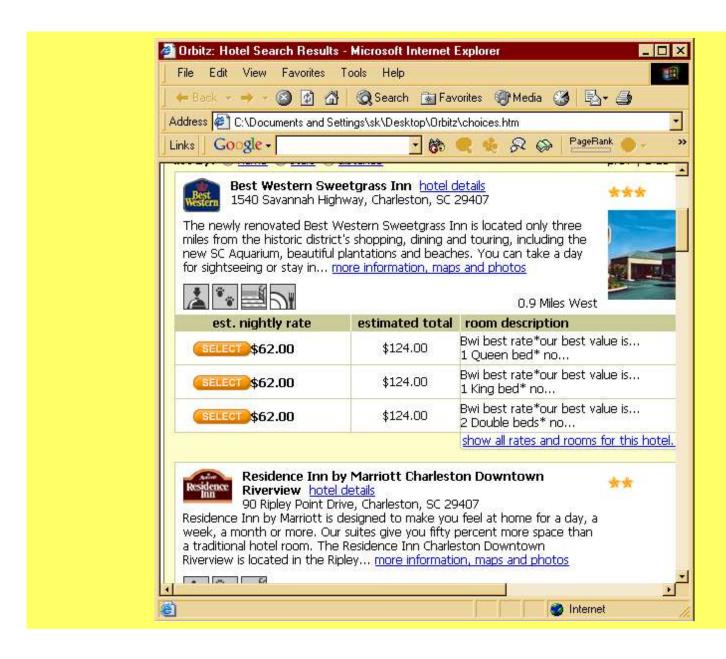
Jwig (Schwartbach & Møller)

Part III

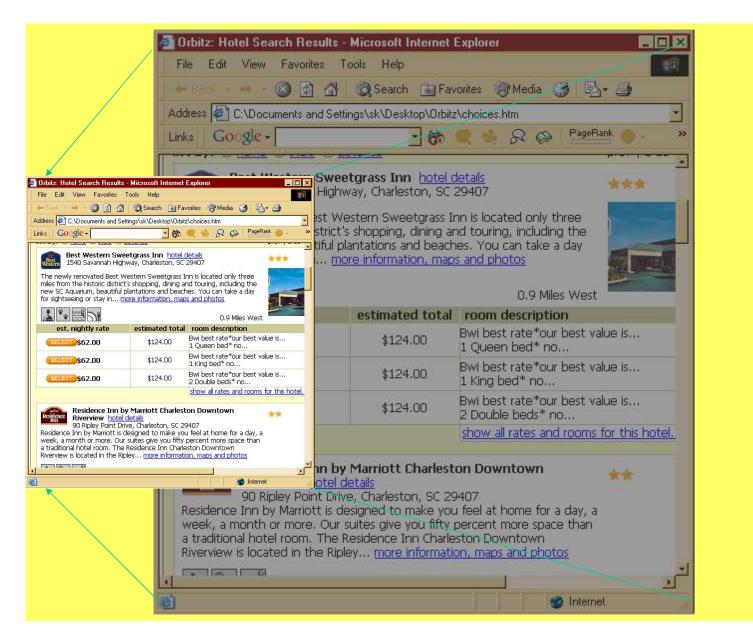
PLT Scheme: Continuations for the Web



Graunke, Findler, Krishnamurthi, Felleisen slides courtesy of Shriram Krishnamurthi



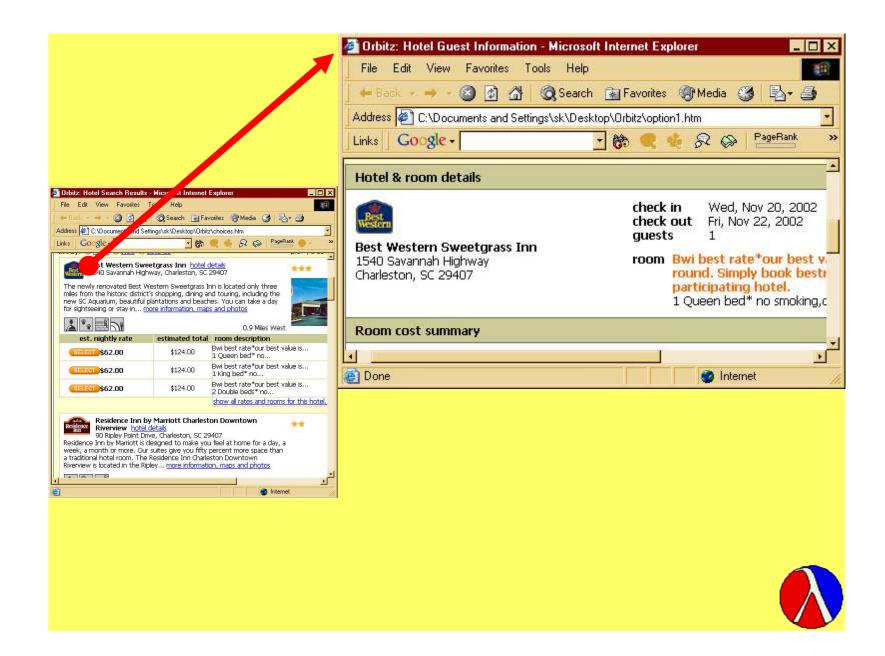


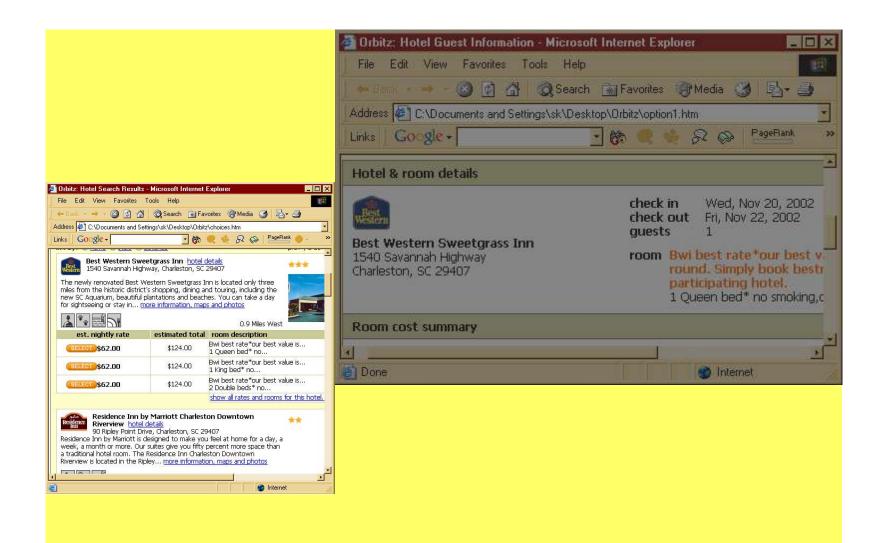




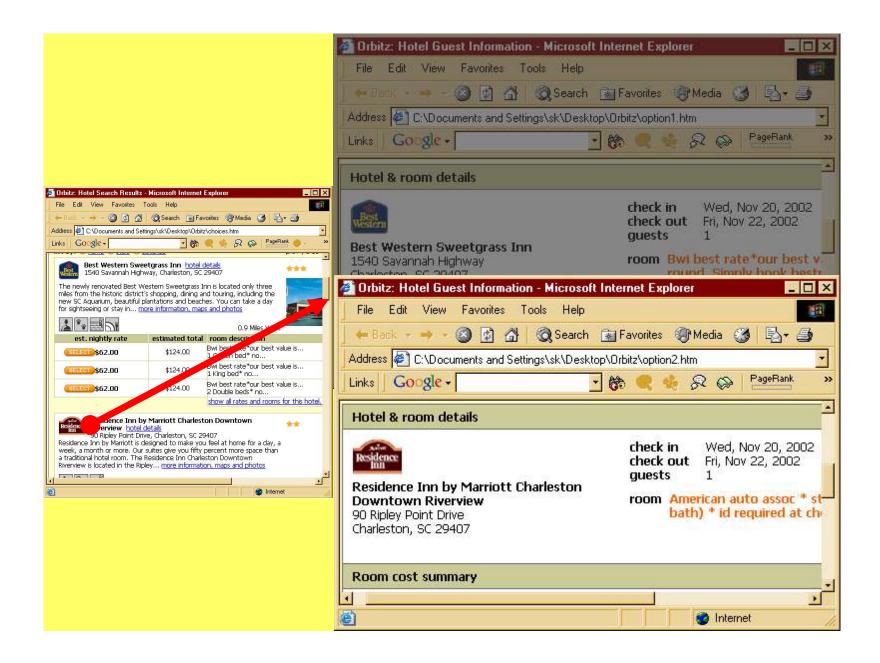


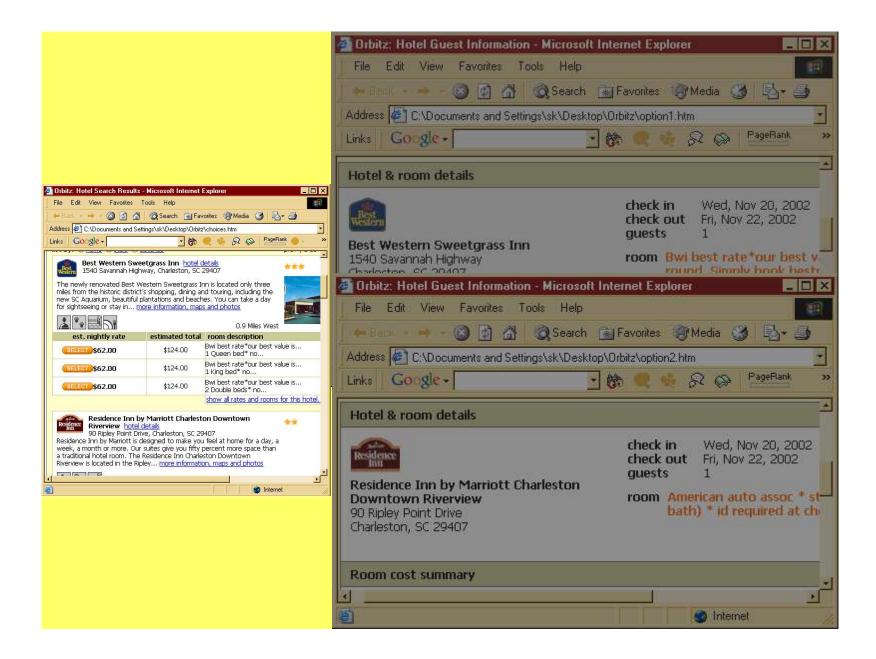


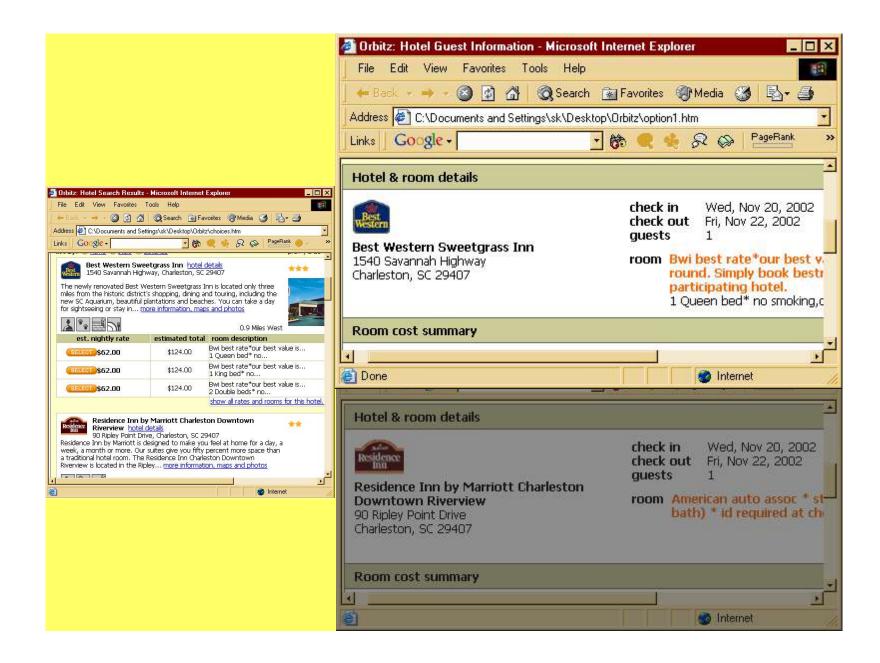


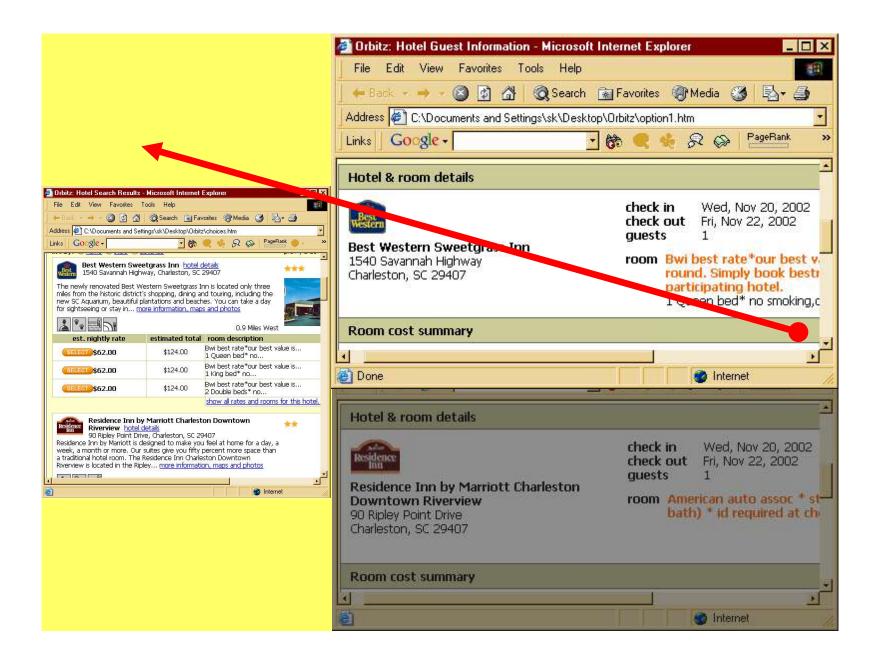


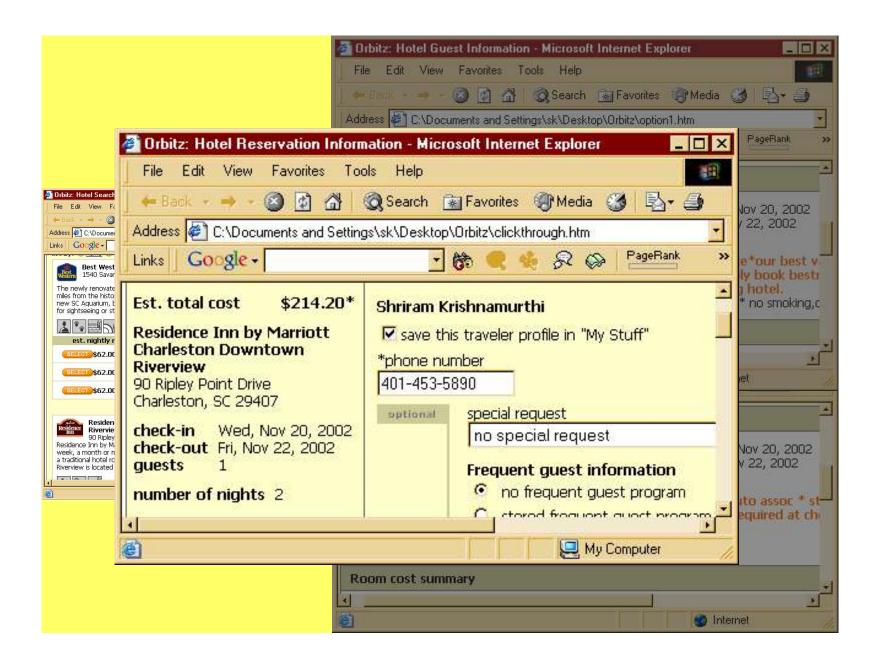


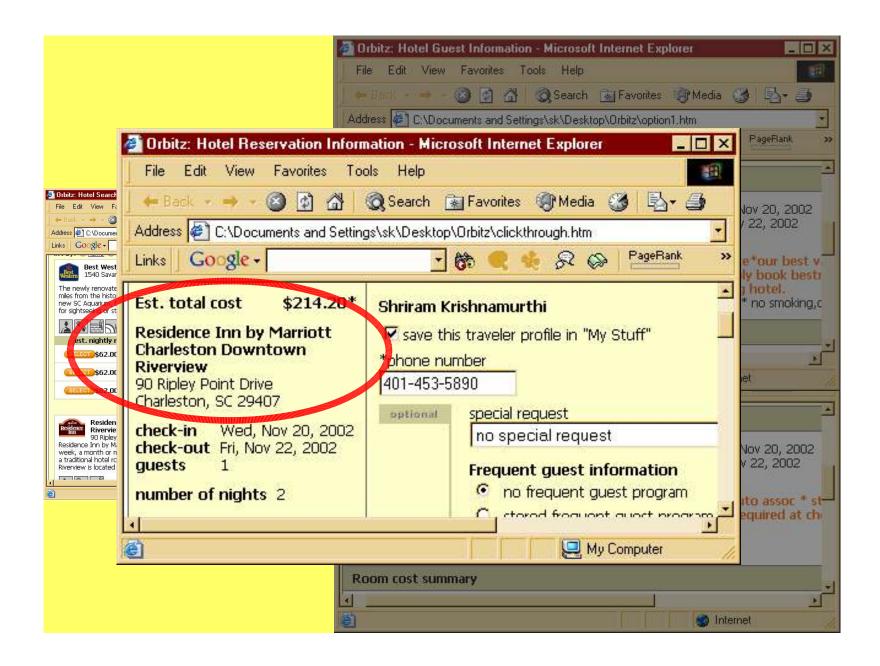






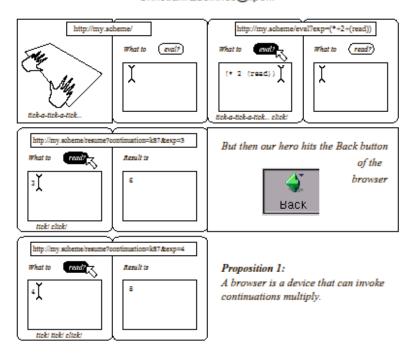




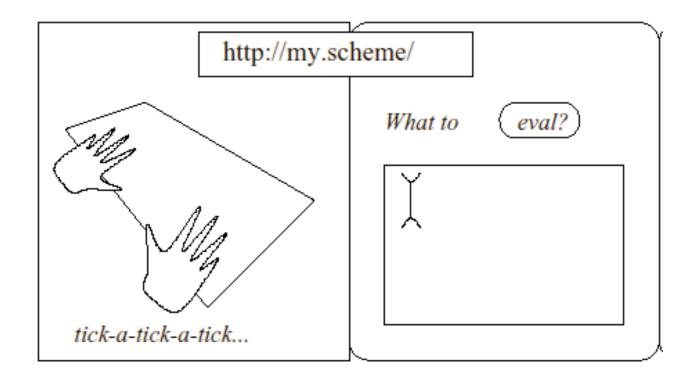


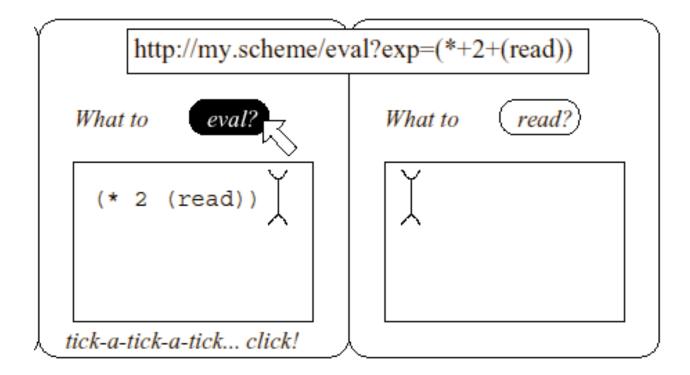
The Influence of Browsers on Evaluators or, Continuations to Program Web Servers [150] \$\frac{1}{2} \text{ [150]} \text{ [150]}\$

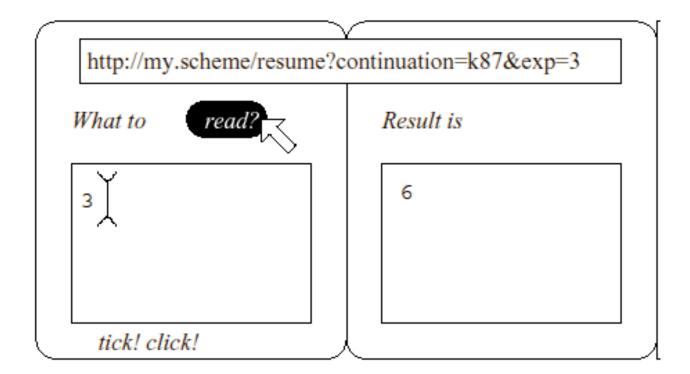
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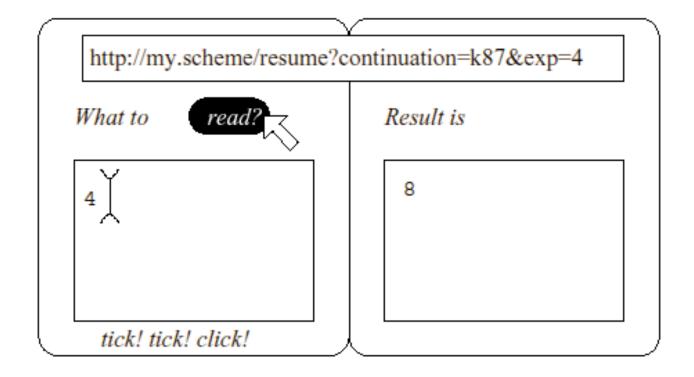
Christian Quenniec (ICFP 2000)











also used by Paul Graham for Yahoo stores

Links

- What are the right scopes? session scope vs. global scope key question for web applications!
- "Scalable" state stored in client, not in server
- Mechanisms
 hooks to choose implementation technique
 URL parameters vs. hidden fields
 cryptographically protect state kept in client
- Other related work
 Mawl (Ramming, Atkins, Ball, Bruns, Cox)
 WASH (Thiemann)

Part IV

Erlang:

Communication via values

Erlang: An area server

```
start() ->
    register (area_server,
         spawn(fun() \rightarrow loop(0) end)).
loop(Tot) ->
    receive
         {Pid, {square, X}} ->
             Pid! X*X,
             loop (Tot + X * X);
         {Pid, {rectangle, [X,Y]}} ->
             Pid! X*Y,
             loop (Tot + X \star Y);
         {Pid, areas} ->
             Pid! Tot,
             loop (Tot)
    end.
```

Armstrong, Virding, Williams & Wikström

Erlang: Generic server

```
start(Name, Data, Fun) ->
    register (Name,
        spawn(fun() -> loop(Data, Fun) end)).
Rpc(Name, Query) ->
    Tag = ref(),
    Name ! {query, self(), Tag, Query},
    receive
        {Tag, Reply} -> Reply
    end.
loop(Data, Fun) ->
    receive
        {query, Pid, Taq, Query} ->
            {Reply, Data1} = Fun(Query, Data),
            Pid! {Tag, Reply},
            loop(Data1, Fun)
    end.
```

Erlang: Instantiating the Generic Server

Erlang: Instantiating a Replicated Server

Links

• Regular expressions for session types

```
session AreaServer =
  (square[Int] | rectangle[Int,Int] | areas[]) *
session InitializedAreaServer =
  initialize[],
  (square[Int] | rectangle[Int,Int] | areas[]) *
```

• User interface as a process

```
session DOM =
  (root[HTML] | todolist[li[String]*])*
```

• Resumption-passing style

Server processes hibernate on client

• Other related work

JoCaml (Fournet & Gonthier)

Timber (Carlsson, Nordlander & Kieburtz)

Part V

Javascript:

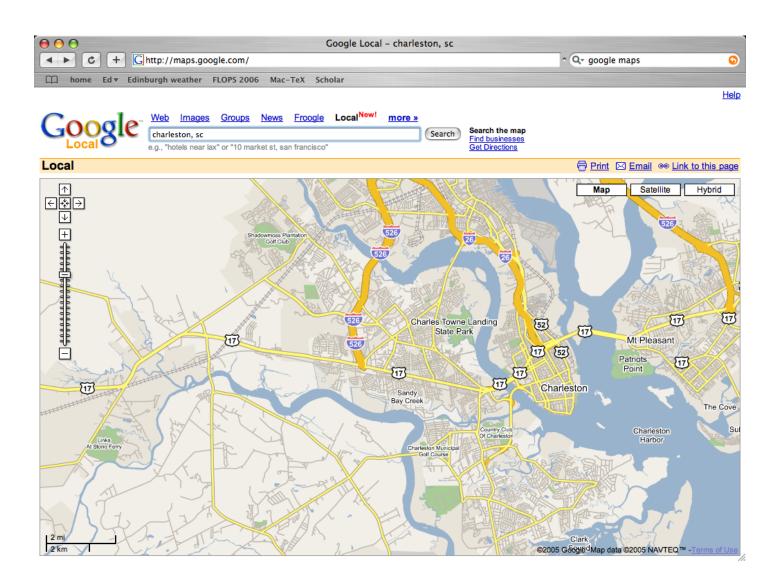
The world's most-widely deployed functional language

Javascript is a functional language

```
Array.prototype.reduce=function(templateFunction) {
  var l=this.length;
  var s='';
  for (var i=0;i<1;i++) s+=templateFunction(this[i]);</pre>
  return s;
function wrap(tag) {
  var stag='<'+tag+'>';
  var etag='</'+tag.replace(/s.*/,'')+'>';
  return function(x) {
    return stag+x+etag;
document.write(
  ' <TABLE><TR>' +
    arr.reduce(wrap('TD class="small"'))+
  '</TR></TABLE>'
);
```

Eich

AJAX

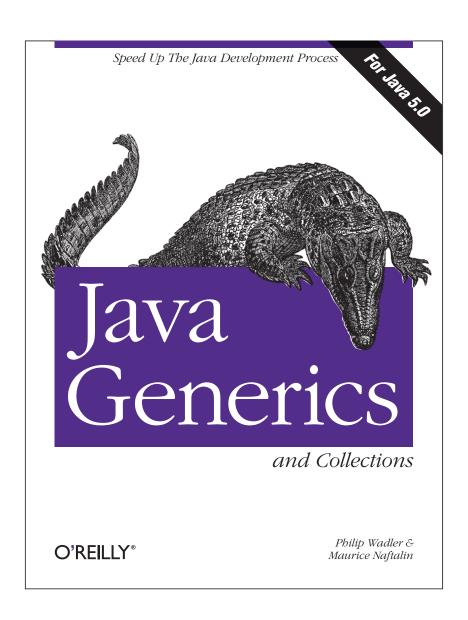


Part VI

Java:

Anothre widely deployed functional language

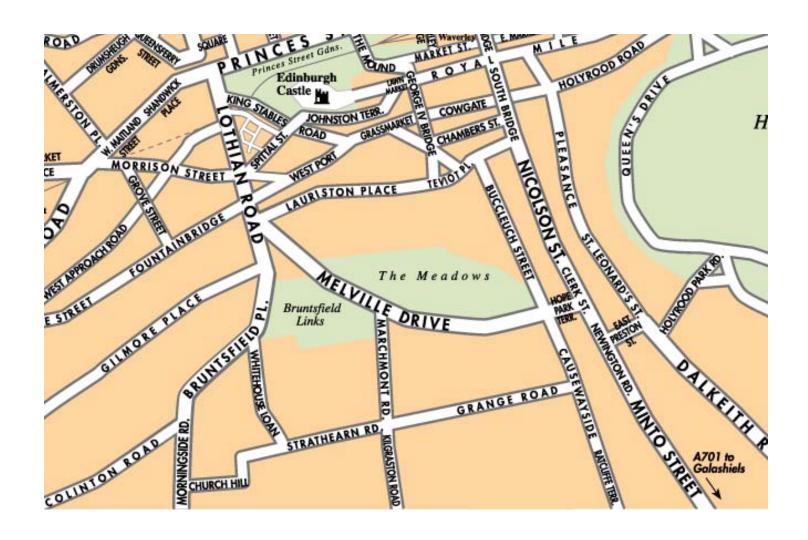
Java generics



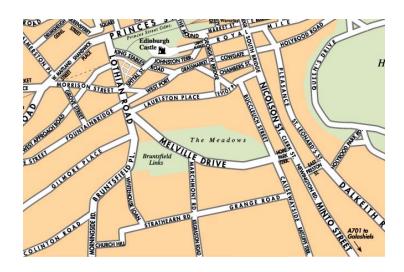
Part VII

Links

Hope and Links



Hope and Links



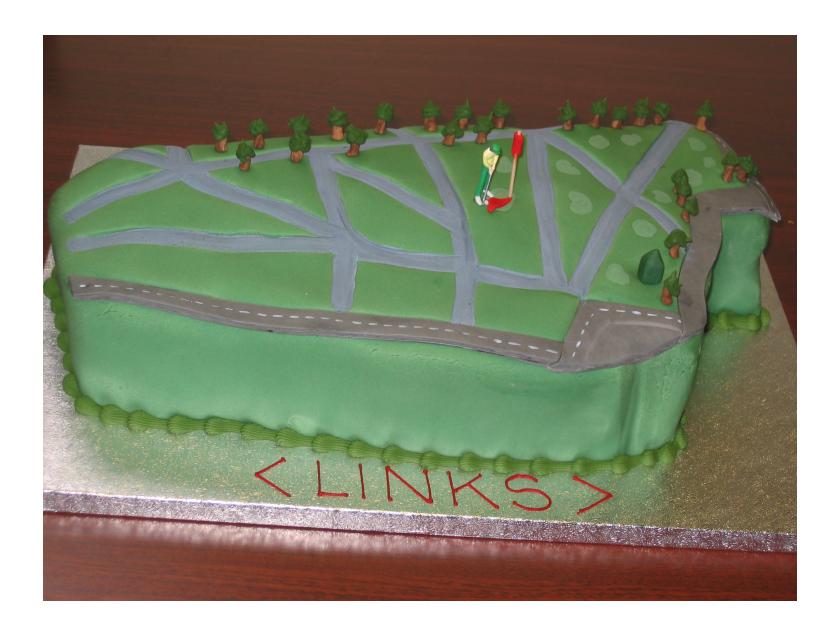
Links

(Bruntsfield Links)

Wadler et al (2005)

Hope
(Hope Park Square)
Burstall, MacQueen,
Sannella (1980)

Links meeting, 6 April 2005

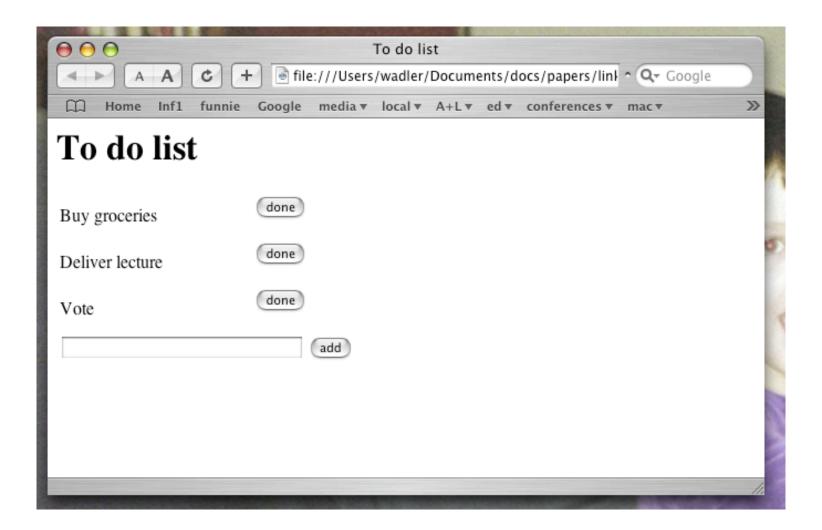


Links meeting, 6 April 2005



Part VIII

A Links program state in client



```
main() { todo([]) }
todo(items) {
  <html><body>
   <h1>Items to do</h1>
   {
     for item in items return
       {item}
         < t.d >
           <form l:action="{todo(items\\[item])}">
             <input type="submit" value="done"/>
           </form>
         }
   <form l:action="{todo(items++[new])}">
     <input l:name="{new}" type="text" size="40">
     <input type="submit" value="add"/>
   </form>
 </body></html>
```

Part IX

A Links program state in server

```
table TODO of (name : String, item: String)
lookup(n) { [ i | (name:n,item:i) <- TODO ] }</pre>
add(n,i) {
  insert into TODO values (name:n, item:i);
  todo(name)
remove(n,i) {
  remove from TODO values (name:n, item:i);
  todo (name)
main() {
  <html><body>
    <h1>Login</h1>
    <form l:action="todo(name)">
      <input l:name="{name}" type="text" size="40">
      <input type="submit" value="login"/>
    </form>
  </body></html>
```

```
todo(name) {
 let items = lookup(name) in
 <html><body>
    <h1>Items to do</h1>
   {
     for item in items return
       \langle t.r \rangle
         {item}
         <form l:action="{remove(name, item)}">
             <input type="submit" value="done"/>
           </form>
         }
   <form l:action="{add(name, new)}">
     <input l:name="{new}" type="text" size="40">
     <input type="submit" value="add"/>
   </form>
 </body></html>
```

Part X

Conclusions

Other ideas

- Multimethods
 Integrate functional and OO styles
- Type classes
 Semantics should determine types, not types determine semantics
- Lists and dictionaries as data structures
 Regular expression matching for lists down with cons!
- Testing and validity Contracts, Quickcheck

Antinomies (technical)

- Pure or effectful?Effects with effect types
- Lazy vs. strict?Strict with support for lazy closures
- Type inference vs. subtyping?
 Give up on type inference
- How to import imperative libraries?
 Hide them in a process

Antinomies (social)

- Is it really research?
 Hard for academics to build real systems
- Is it too much research?
 Haskell and SML built on strong concensus
- How do we build a community?
 We can only succeed if we do it together

</Links>